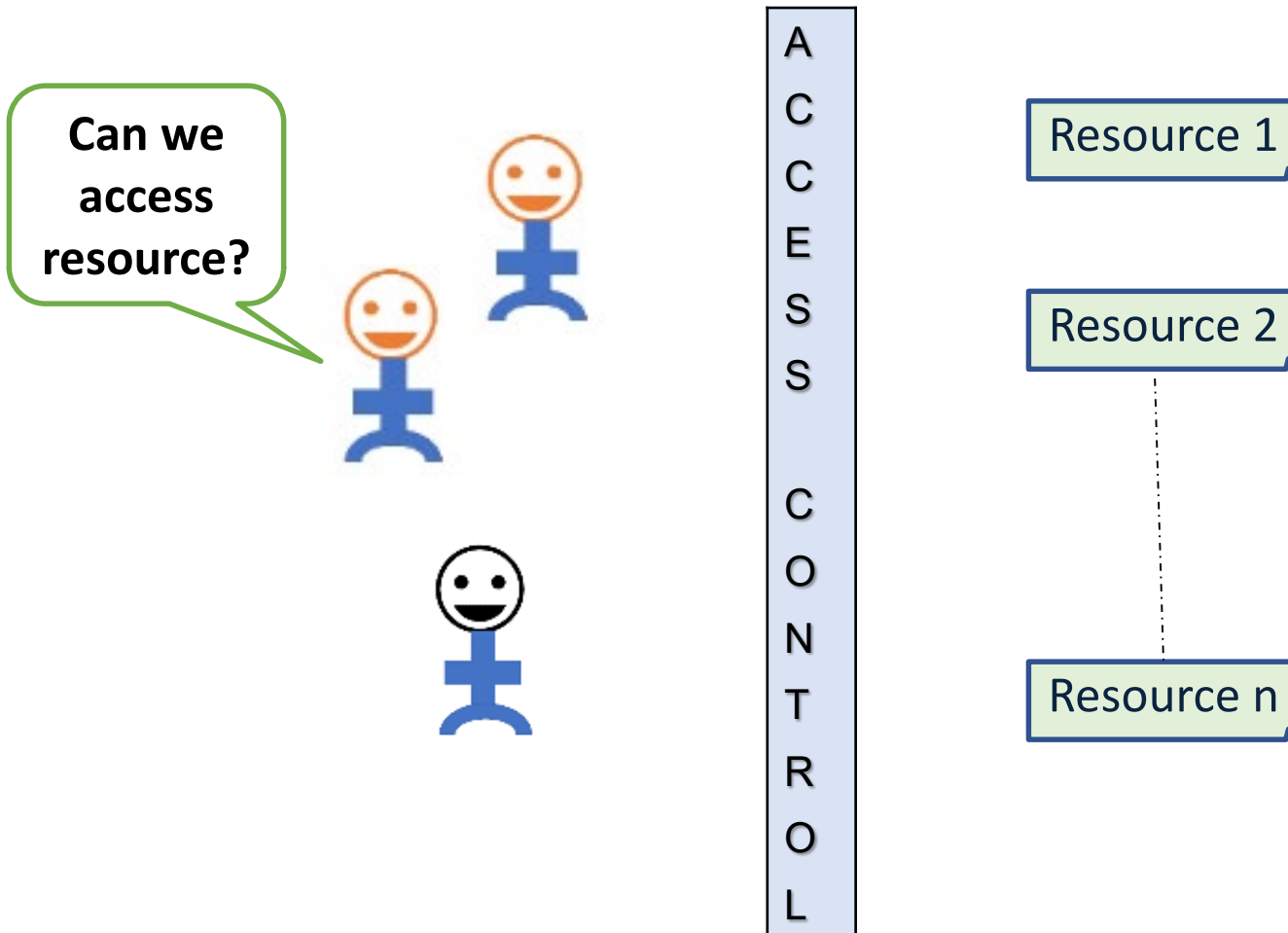


Formal Analysis of ReBAC Policy Mining Feasibility

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Legitimate users get legitimate access only
i.e., Role-Based Access Control (RBAC), Attribute-Based Access Control (ABAC)

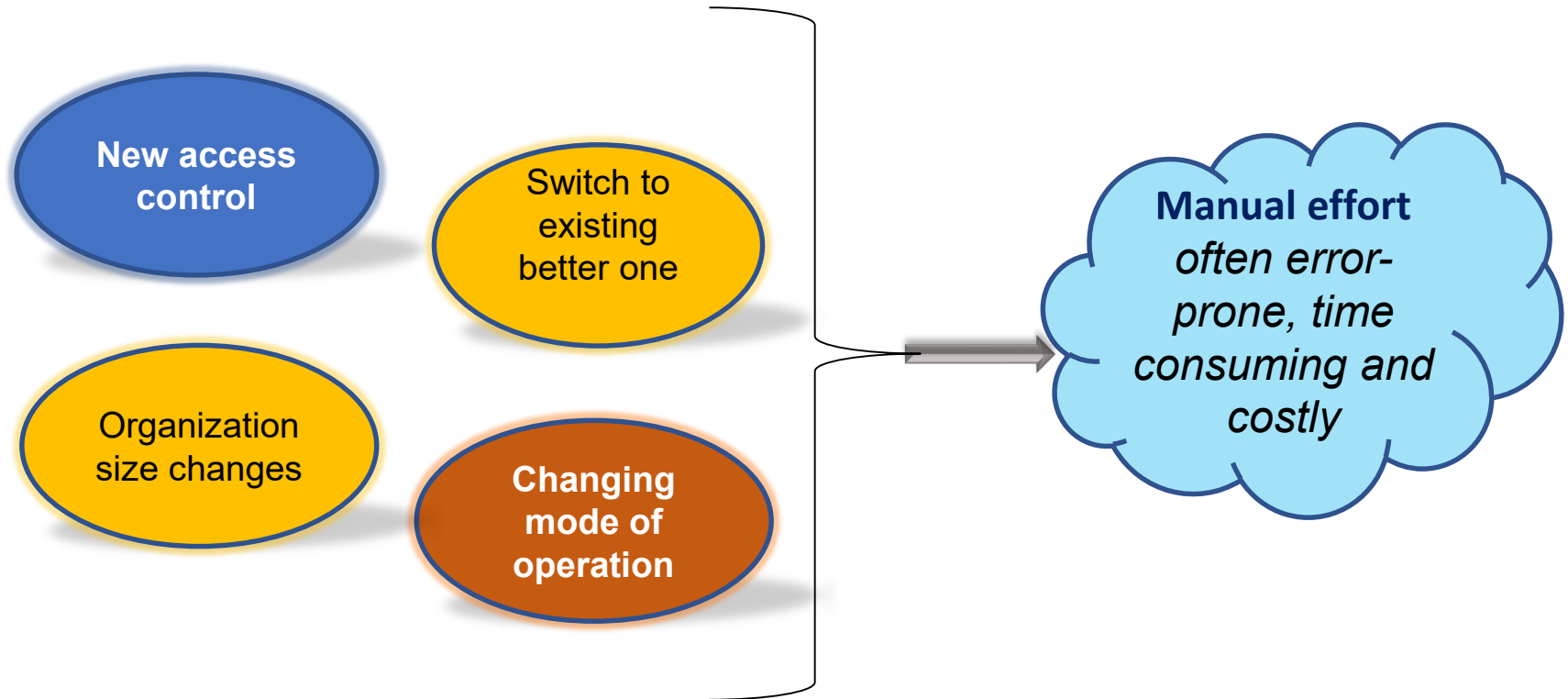
❖ *ReBAC* \equiv *Relationship-Based Access Control*

- ReBAC expresses authorization in terms of various direct and indirect relationships amongst entities, most commonly between users
- Access conditions are usually based on type, depth, or strength of relationships

❖ Assumption

- Relationship Graph (RG) where users(node) are connected(edge) by social relationships(edge label). Each edge in the RG is labeled with a relation type
- Only user-to-user relationships are considered

❖ **Problem:** migration from an existing access control model to another one



Is automation possible?

Access Control List / Log / RBAC +
Supporting attribute data



ABAC policy mining

Access Control List + Supporting
Relationship data



ReBAC policy mining



Given an access control system +
Supporting data



Another access control
model



*The feasibility analysis of the ReBAC policy mining problem studies whether the migration process from a given authorization set to ReBAC policy is feasible or not under the set of **imposed criteria**:*

- ❖ Relationship Graph (RG) is given
- ❖ ReBAC rule structure is given
- ❖ Use of entity ID is not allowed
 - Existing literature allows ID
- ❖ Equivalent set of ReBAC rules are required

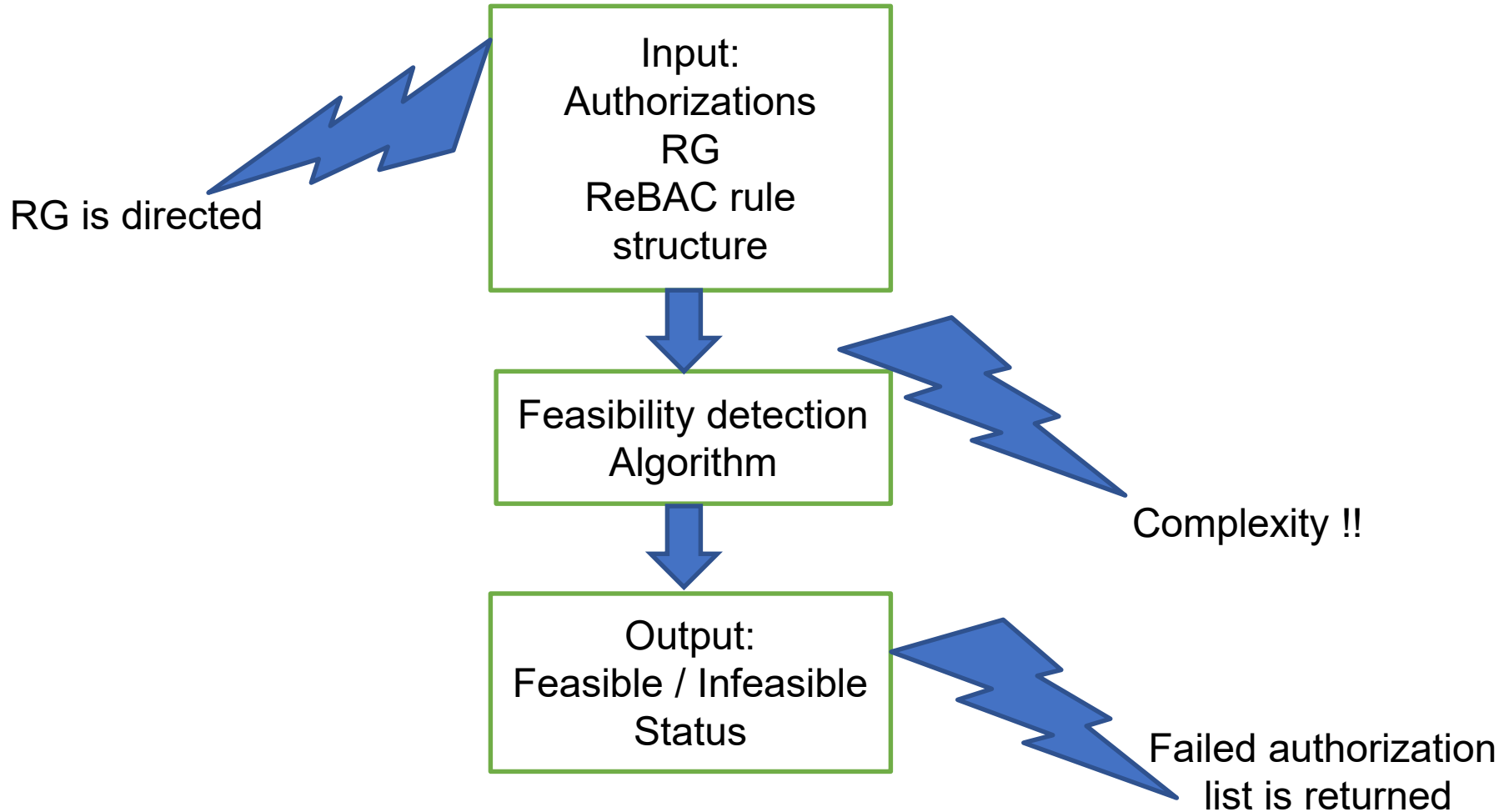
- ❖ Solution is guaranteed even if inconsistency arises
 - Infeasibility problem

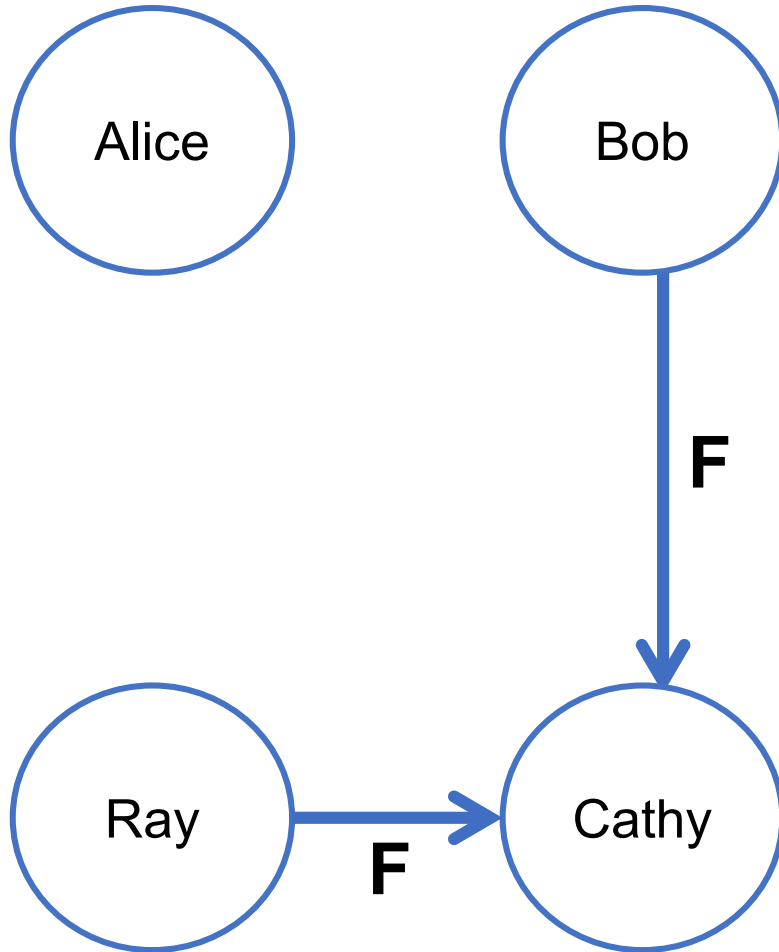
- ❖ Feasibility analysis on ReBAC policy mining for the first time
- ❖ Developing feasibility analysis algorithms for the given set of criteria with complexity analysis
 - Variety of ReBAC rule structures are considered
- ❖ In case of infeasibility, solution algorithms are presented to make it feasible under given criteria
 - Varieties available
- ❖ Demonstrate the generated algorithms with cases and show the effectiveness beyond complexity analysis
- ❖ Future scopes

$$\begin{aligned}
 Rule_{op} &::= Rule_{op} \vee Rule_{op} \mid pathRuleExpr \\
 pathRuleExpr &::= pathRuleExpr \wedge \\
 &\quad pathRuleExpr \mid pathLabelExpr \\
 pathLabelExpr &::= pathLabelExpr.pathLabelExpr \mid edgeLabel \\
 edgeLabel &::= \sigma, \sigma \in \Sigma
 \end{aligned}$$

- ❖ Evaluation of access request (a, b, op)
 - for each pathLabelExpr in $Rule_{op}$ substitute True if there exists a simple path p from a to b in RG with path label pathLabelExpr, otherwise substitute False
 - the resulting boolean expression evaluates true \rightarrow grant, deny otherwise

RREP(ReBAC Ruleset Existence Problem)-0





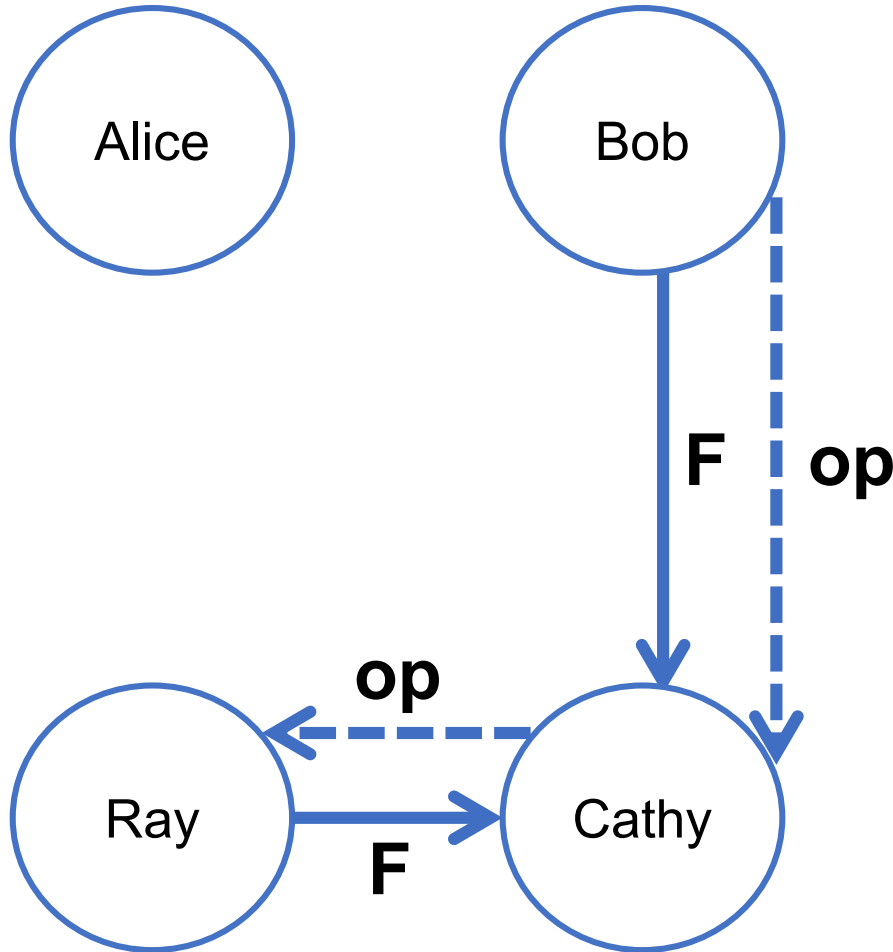
Feasible

(Bob, Cathy, op)
(Ray, Cathy, op)

Rule_{op} = F

Infeasible

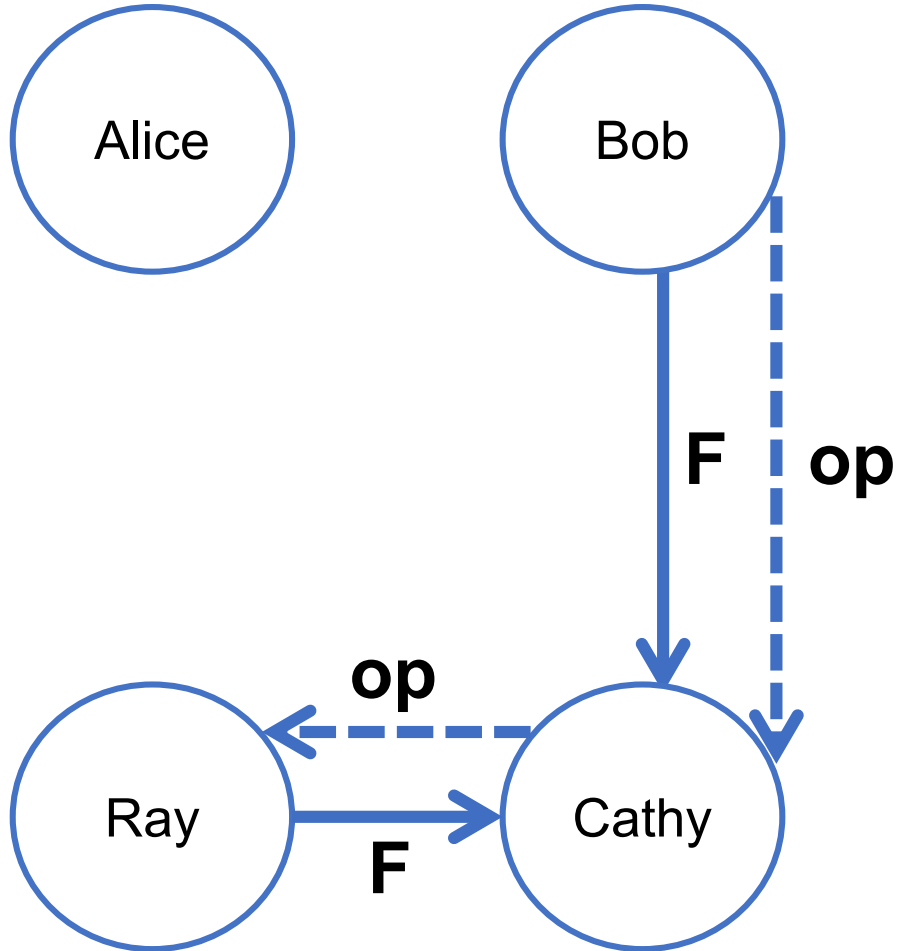
- i) (Bob, Cathy, op)
- ii) (Cathy, Ray, op)



Infeasible

- i) (Bob, Cathy, op)
- ii) (Cathy, Ray, op)

Rule_{op} = op



Infeasible

- i) (Bob, Cathy, op)
- ii) (Cathy, Ray, op)

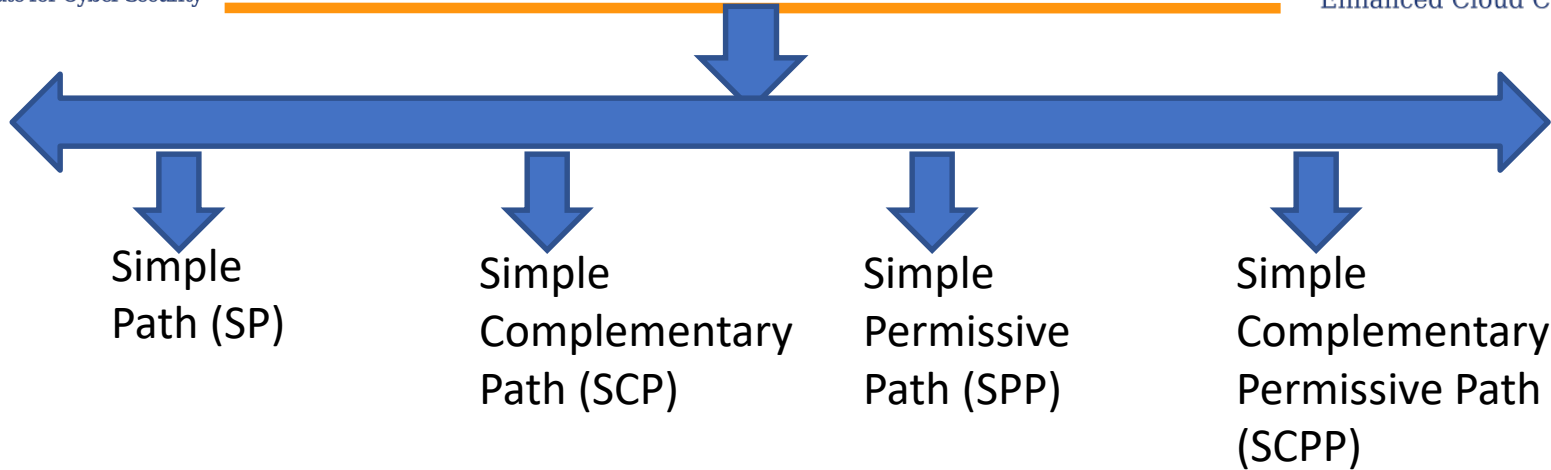
Rule_{op} = op

Simple

Operation \cap Relationship types = {}

Minimal edges not guaranteed

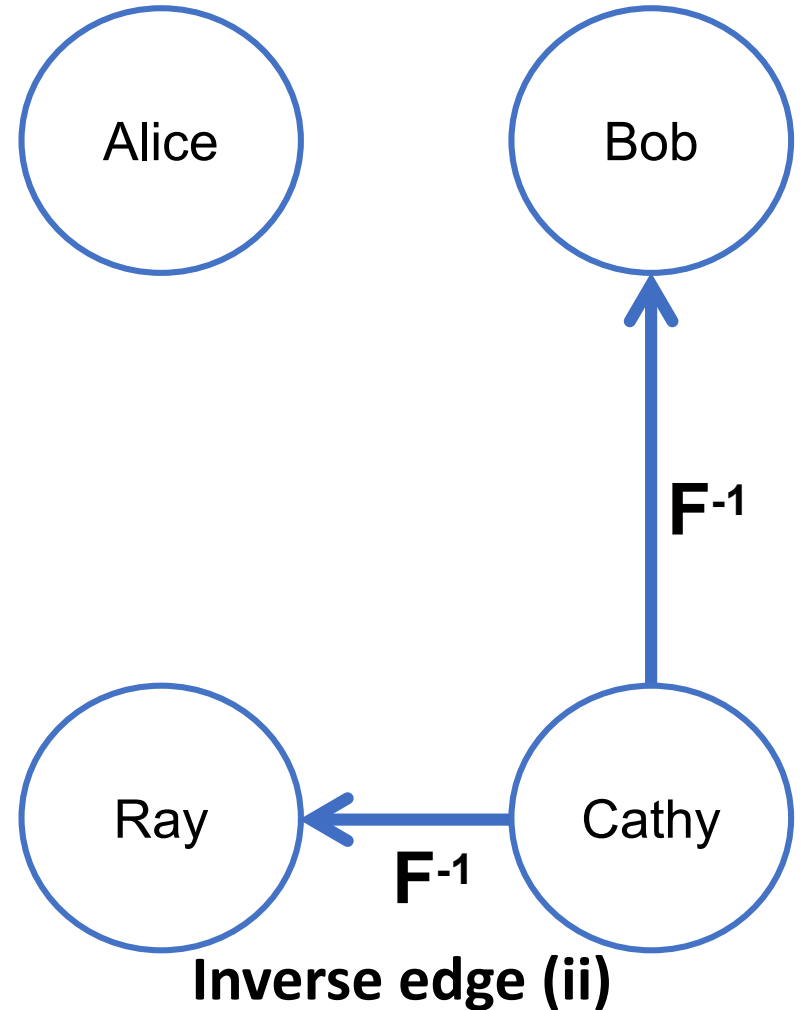
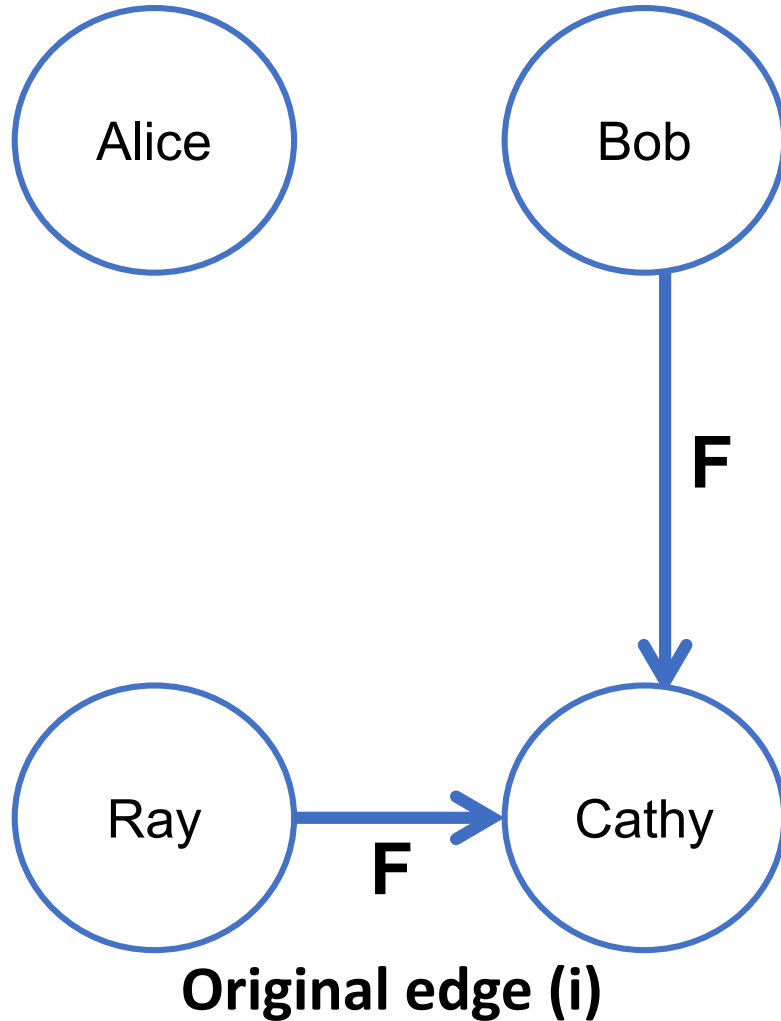
|Authorization| edges at worst!

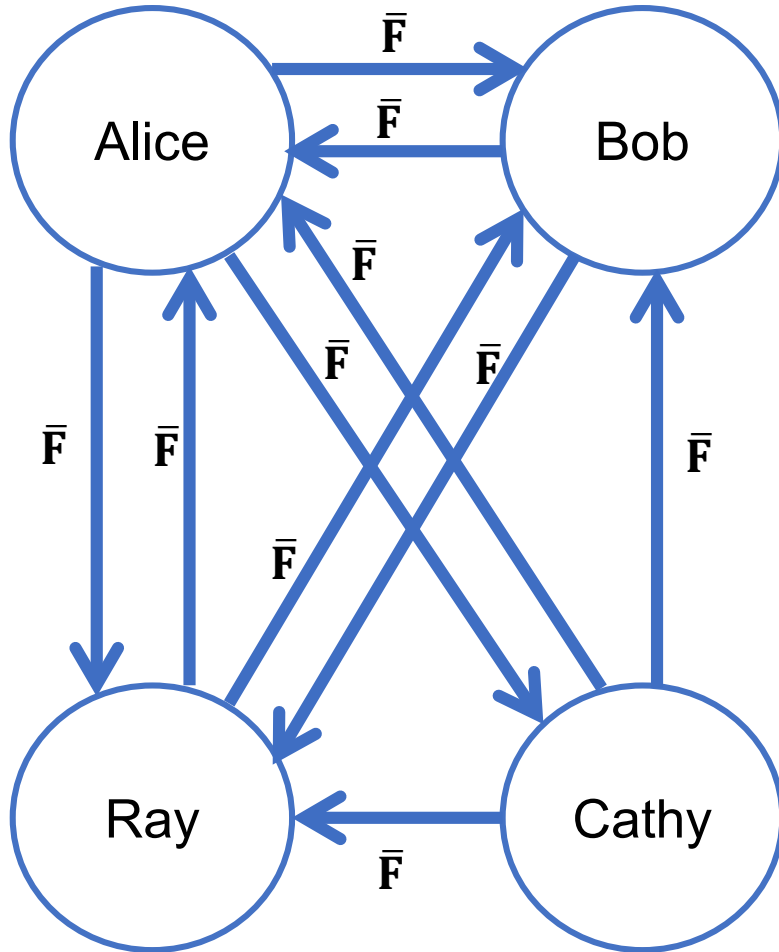


Characteristics	SCP	SPP	SCPP
$(a, b, \sigma) \rightarrow (a, b, \sigma) \in E, \sigma \in \Sigma$	✓	✓	✓
$(a, b, \bar{\sigma}) \rightarrow (a, b, \sigma) \notin E, \bar{\sigma} \in \bar{\Sigma}$	✓		✓
$(a, b, \sigma^{-1}) \rightarrow (b, a, \sigma) \in E, \sigma^{-1} \in \Sigma^{-1}$		✓	✓
$(a, b, \bar{\sigma}^{-1}) \rightarrow (b, a, \sigma) \notin E, \bar{\sigma}^{-1} \in \bar{\Sigma}^{-1}$			✓

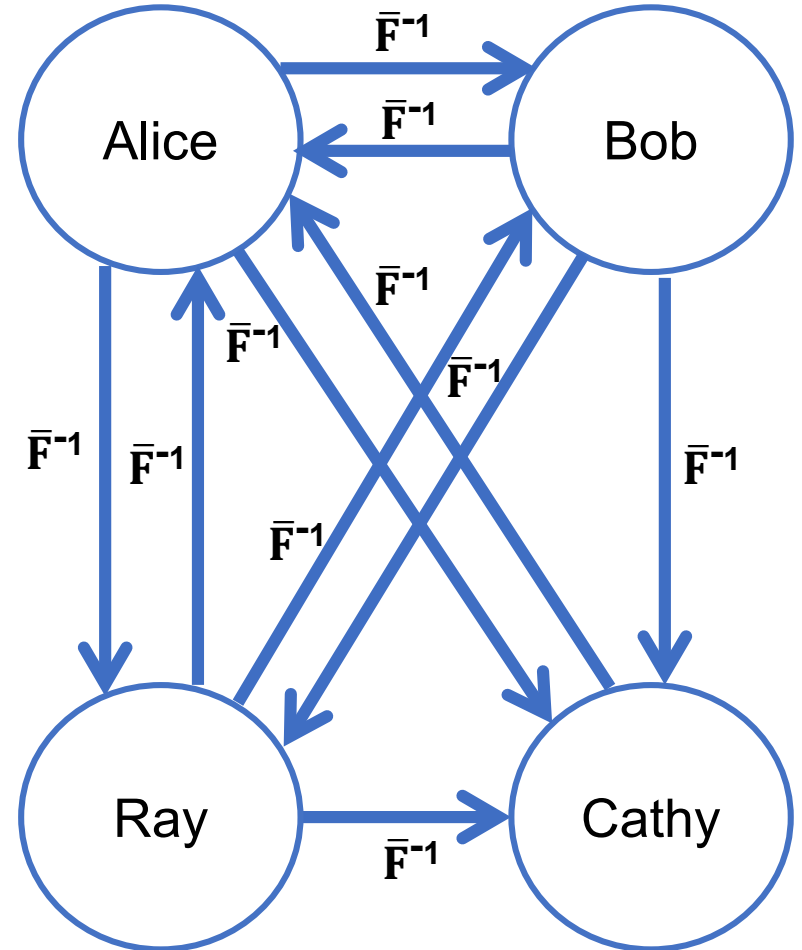
Table represents path variations with original, non-relationship, inverse and non-relationship inverse edges (row 1, 2, 3, and 4, respectively).

- a,b: users, E and Σ are the sets of edges and relationship type specifiers





Non-relationship edge (iii)



Non-relationship inverse edge (iv)

RREP-0 → SP (i)

RREP-1 → SCP (i + iii)

RREP-2 → SPP (i + ii)

RREP-3 → SCPP (i + ii + iii + iv)

Rule minimization techniques are described in the paper

- ❖ Complexity
- ❖ Inexact solution
- ❖ More path variations
- ❖ Cope up with changes in rule structures!
- ❖ Other infeasibility solutions
- ❖ Extend beyond user-user context

!! Just the beginning !!

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- ❖ Question/ Feedback