

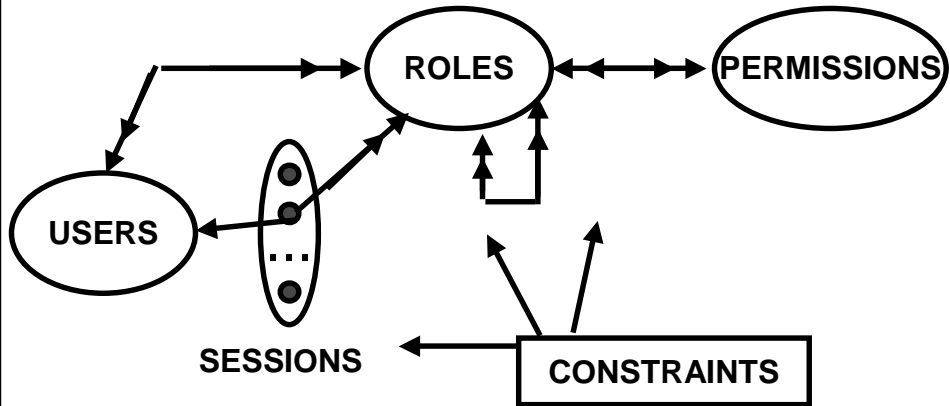
The *RCL2000* Language for Specifying Role-Based Authorization Constraints

Gail-Joon Ahn

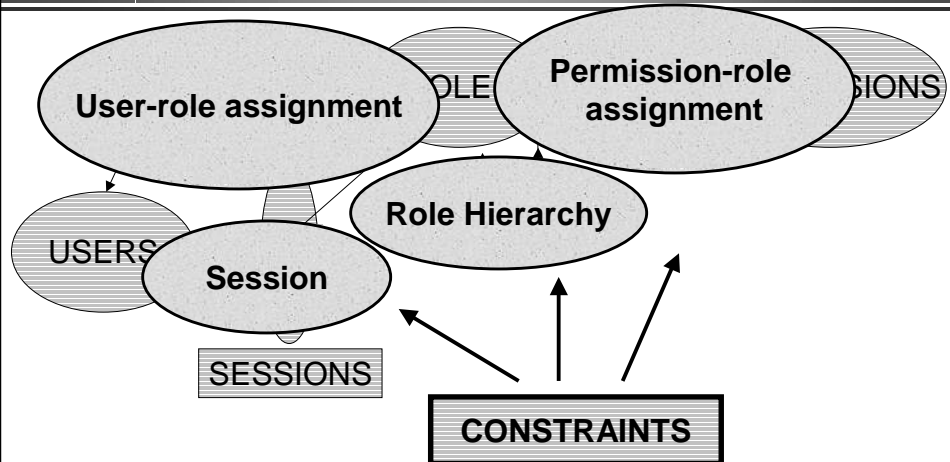
ABSTRACT

- ◆ **This presentation includes**
 - **The first formal (and intuitive) language for role-based authorization constraints**
 - **A formal semantics for this language**
 - **Demonstration of the expressive power of the language**
 - **Characterization of role-based constraints into prohibition and obligation constraints**

RBAC96



RBAC96



SEPARATION OF DUTY (1)

- ◆ **SOD is fundamental technique for preventing fraud and errors**

- ◆ **Related Work**
 - **Enumerate several forms of SOD**
 - **Little work on specifying SOD in a comprehensive way**

SEPARATION OF DUTY (2)



**PURCHASING
MANAGER**

**ACCOUNTING PAYABLE
MANAGER**

PROHIBITION

- ◆ **Separation of Duty constraints**

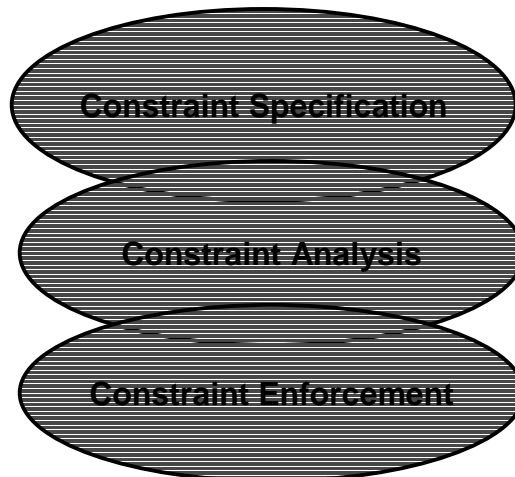
OBLIGATION

- ◆ **Every faculty member must be assigned to at least one departmental committee**

RESEARCH PLAN

- ◆ **Need to specify these constraints**
 - **Language**
- ◆ **Show the meaning of expression**
 - **Formal semantics**
- ◆ **Expressive power of the language**
 - **Well-known constraints and simulations**
- ◆ **Analysis of the work**
 - **Characterization**

BIG PICTURE



WHO IS THE USER

- ◆ **Security Researcher**
- ◆ **Security Policy Designer**
- ◆ **Security Architect**

RCL 2000

- ◆ **RCL 2000 (Role-based Constraints Language 2000)**
- ◆ **Specification Language**
 - **to formally express constraints in role-based systems**
- ◆ **Most components are built upon RBAC96**

BASIC ELEMENT (from RBAC96)

- ◆ **U** : a set of users
- ◆ **R** : a set of roles
 - $RH \subseteq R \times R$: role hierarchy
- ◆ **OBJ** : a set of objects
- ◆ **OP** : a set of operations
- ◆ **P = OP × OBJ** : a set of permissions
- ◆ **S** : a set of sessions

BASIC ELEMENT (from RBAC96)

- ◆ **UA** : a many-to-many user-to-role assignment relation
- ◆ **PA** : a many-to-many permissions-to-role assignment relation

SYSTEM FUNCTIONS (from RBAC96)

- ◆ **user** : $R \rightarrow 2^U$
- ◆ **roles** : $U \cup P \cup S \rightarrow 2^R$
- ◆ **sessions** : $U \rightarrow 2^S$
- ◆ **permissions** : $R \rightarrow 2^P$
- ◆ **operations** : $R \times \text{OBJ} \rightarrow 2^{\text{OP}}$
- ◆ **object** : $P \rightarrow 2^{\text{OBJ}}$

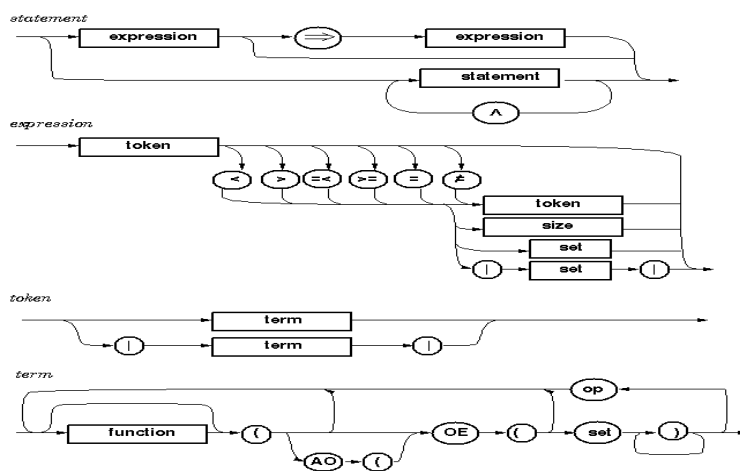
BASIC ELEMENT (beyond RBAC96)

- ◆ **CR** : all conflicting role sets
- ◆ **CU** : all conflicting user sets
- ◆ **CP** : all conflicting permission sets

NON-DETERMINISTIC FUNCTIONS (beyond RBAC96)

- ◆ introduced by Chen and Sandhu (1995)
- ◆ oneelement (OE)
 - $\text{oneelement}(X) = x_i$, where $x_i \in X$
- ◆ allother (AO)
 - $\text{allother}(X) = X - \{\text{OE}(X)\}$
 $= X - \{x_i\}$
 - should occur along with OE function

SYNTAX



EXAMPLES OF CONSTRAINT EXPRESSION

Conflicting roles cannot have common users

- $|\text{roles}(\text{OE}(\text{U})) \cap \text{OE}(\text{CR})| \leq 1$

Conflicting users cannot have common roles

- $\text{roles}(\text{OE}(\text{OE}(\text{CU}))) \cap \text{roles}(\text{AO}(\text{OE}(\text{CU}))) = \emptyset$

Users cannot activate two conflicting roles

- $|\text{roles}(\text{sessions}(\text{OE}(\text{U}))) \cap \text{OE}(\text{CR})| \leq 1$

Users cannot activate two conflicting roles in a single session

- $|\text{roles}(\text{OE}(\text{sessions}(\text{OE}(\text{U})))) \cap \text{OE}(\text{CR})| \leq 1$

FORMAL SEMANTICS

◆ Reduction Algorithm

- to convert a constraint expression to a restricted form of first order predicate logic (RFOPL)

◆ Construction Algorithm

- to construct a constraint expression from RFOPL

REDUCTION ALGORITHM

$$OE(OE(CR)) \in \text{roles}(OE(U)) \Rightarrow AO(OE(CR)) \cap \text{roles}(OE(U)) = \emptyset$$

1. $OE(OE(CR)) \in \text{roles}(OE(U)) \Rightarrow (OE(CR) - \{OE(OE(CR))\}) \cap \text{roles}(OE(U)) = \emptyset$
2. $\forall cr \in CR : OE(cr) \in \text{roles}(OE(U)) \Rightarrow (cr - \{OE(cr)\}) \cap \text{roles}(OE(U)) = \emptyset$
3. $\forall cr \in CR, \forall r \in cr : r \in \text{roles}(OE(U)) \Rightarrow (cr - \{r\}) \cap \text{roles}(OE(U)) = \emptyset$
4. $\forall cr \in CR, \forall r \in cr, \forall u \in U : r \in \text{roles}(u) \Rightarrow (cr - \{r\}) \cap \text{roles}(u) = \emptyset$

RFOPL STRUCTURE

- ◆ **sequence part : predicate**
- ◆ $\forall r \in R, \forall u \in U : r \in \text{roles}(u)$
- ◆ $\forall x_2 \in x_1, \forall x_3 \in x_2, \forall x_4 \in x_3 : \text{predicate}$

CONSTRUCTION ALGORITHM

$$\forall cr \in CR, \forall r \in cr, \forall u \in U : r \in \text{roles}(u) \Rightarrow (cr - \{r\}) \cap \text{roles}(u) = \emptyset$$

$$1. \forall cr \in CR, \forall r \in cr : r \in \text{roles}(\text{OE}(U)) \Rightarrow (cr - \{r\}) \cap \text{roles}(\text{OE}(U)) = \emptyset$$

$$2. \forall cr \in CR : \text{OE}(cr) \in \text{roles}(\text{OE}(U)) \Rightarrow (cr - \{\text{OE}(cr)\}) \cap \text{roles}(\text{OE}(U)) = \emptyset$$

$$3. \text{OE}(\text{OE}(CR)) \in \text{roles}(\text{OE}(U)) \Rightarrow (\text{OE}(CR) - \{\text{OE}(\text{OE}(CR))\}) \cap \text{roles}(\text{OE}(U)) = \emptyset$$

$$4. \text{OE}(\text{OE}(CR)) \in \text{roles}(\text{OE}(U)) \Rightarrow \text{AO}(\text{OE}(CR)) \cap \text{roles}(\text{OE}(U)) = \emptyset$$

SOUNDNESS AND COMPLETENESS

◆ **Theorem 1** *Given RCL2000 expression α , α can be translated into RFOPL expression β . Also α can be reconstructed from β .*

$$\mathbf{C(R(\alpha)) = \alpha}$$

◆ **Theorem 2** *Given RFOPL expression β , β can be translated into RCL2000 expression α . Also β' which is logically equivalent to β can be reconstructed from α .*

$$\mathbf{R(C(\beta)) = \beta'}$$

SEPARATION OF DUTY CONSTRAINTS

- ◆ **Specification of SOD constraints identified by Simon and Zurko (1997) and formulated by Virgil et al (1998)**
- ◆ **Identify new SOD properties**
 - **Role-centric**
 - **User-centric**
 - **Permission-centric**

ROLE-CENTRIC SOD CONSTRAINT EXPRESSION

- ◆ **Static SOD**
 - : **Conflicting roles cannot have common users**
$$U = \{u_1, u_2, \dots, u_n\}, R = \{r_1, r_2, \dots, r_n\},$$
$$CR = \{cr_1, cr_2\} : cr_1 = \{r_1, r_2, r_3\}, cr_2 = \{r_a, r_b, r_c\}$$
 - $|\text{roles}(\text{OE}(U)) \cap \text{OE}(CR)| \leq 1$

PERMISSION-CENTRIC SOD CONSTRAINT EXPRESSION

- ◆ **SSOD-CP**
 - $|\text{permissions}(\text{roles}(\text{OE}(\text{U}))) \cap \text{OE}(\text{CP})| \leq 1$

- ◆ **Variations of SSOD-CP**
 - $\text{SSOD-CP} \wedge |\text{permissions}(\text{OE}(\text{R})) \cap \text{OE}(\text{CP})| \leq 1$

USER-CENTRIC SOD CONSTRAINT EXPRESSION

- ◆ **SSOD-CU (User-centric)**
 - $\text{SSOD-CR} \wedge |\text{user}(\text{OE}(\text{CR})) \cap \text{OE}(\text{CU})| \leq 1$

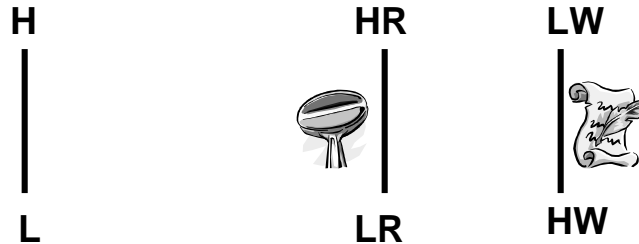
DYNAMIC SOD

- ◆ **User-based DSOD**
 - $|\text{roles}(\text{sessions}(\text{OE}(\text{U}))) \cap \text{OE}(\text{CR})| \leq 1$
- ◆ **User-based DSOD with CU**
 - $|\text{roles}(\text{sessions}(\text{OE}(\text{OE}(\text{CU})))) \cap \text{OE}(\text{CR})| \leq 1$
- ◆ **Session-based DSOD**
 - $|\text{roles}(\text{OE}(\text{sessions}(\text{OE}(\text{U})))) \cap \text{OE}(\text{CR})| \leq 1$
- ◆ **Session-based DSOD with CU**
 - $|\text{roles}(\text{OE}(\text{sessions}(\text{OE}(\text{OE}(\text{CU})))) \cap \text{OE}(\text{CR})| \leq 1$

CASE STUDIES

- ◆ **Lattice-based access control**
 - Ravi Sandhu (1993, 1996)
- ◆ **Chinese Wall policy**
 - Ravi Sandhu (1992)
- ◆ **Discretionary access control**
 - Sandhu and Munawer (1998)

LATTICE-BASED ACCESS CONTROL



- ◆ Subject s can write object o only if $\lambda(s) \leq \lambda(o)$
- ◆ Subject s can read object o only if $\lambda(o) \leq \lambda(s)$

Constraints on UA: *Each user is assigned to exactly two roles xR and LW*

LATTICE-BASED ACCESS CONTROL

- $AR = \{ar1, ar2\}$
 - $ar1 = \{HR, HW\}$, $ar2 = \{LR, LW\}$
- $ASR = \{asr1, asr2\}$
 - $asr1 = \{HR, LW\}$, $asr2 = \{LR, LW\}$
- ◆ **Constraint on UA:**
 - $roles(OE(U)) = OE(ASR)$
- ◆ **Constraint on sessions:**
 - $roles(OE(sessions(OE(U)))) = OE(AR)$

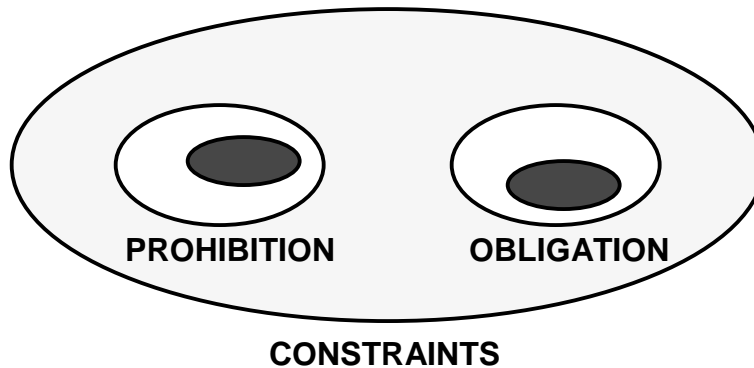
PROHIBITION CONSTRAINTS

- ◆ **Forbid the RBAC component from doing (or being) something which is not allowed to do (or be)**
 - **Separation of duty constraints**

OBLIGATION CONSTRAINTS

- ◆ **Force the RBAC component to do (or be) something**
 - **LBAC-RBAC, Chinese Wall-RBAC simulation**

CONSTRAINTS CHARACTERIZATION



SIMPLE PROHIBITION CONSTRAINTS

- ◆ **Type 1**
 - $|expr| \leq 1$
- ◆ **Type 2**
 - $expr = \phi$ or $|expr| = 0$
- ◆ **Type 3**
 - $|expr1| < |expr2|$

SIMPLE OBLIGATION CONSTRAINTS

- ◆ **Type 1**
 - $expr \neq 0$ or $|expr| > 0$
- ◆ **Type 2**
 - Set X = Set Y
- ◆ **Type 3**
 - obligation constraints \Rightarrow obligation constraints
- ◆ **Type 4**
 - $|expr| = 1$
 - $|expr| = 1 \equiv |expr| \leq 1 \wedge |expr| > 0$

CONTRIBUTIONS

- ◆ **Developed the first formal and intuitive language for role-based authorization constraints**
- ◆ **Provided a formal semantics for this language**
- ◆ **Demonstrated the expressive power of the language by**
 - specifying well-known separation of duty constraints
 - identifying new role-based SOD constraints
 - showing how to specify constraints identified in the simulations of other policies in RBAC
- ◆ **Characterized role-based constraints into prohibition and obligation constraints**

FUTURE WORK

- ◆ **Extension of RCL 2000**
 - Applying it the formalization of some realistic security policies
- ◆ **Implementation Issue**
 - Tool for checking syntax and semantic as well as visualization of specification
- ◆ **Enforcement of constraints**