# **TOPIC**

# LATTICE-BASED ACCESS-CONTROL MODELS

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# LATTICE-BASED MODELS

- Denning's axioms
- Bell-LaPadula model (BLP)
- Biba model and its duality (or equivalence) to BLP
- Dynamic labels in BLP

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## **DENNING'S AXIOMS**

< SC,  $\rightarrow$ ,  $\oplus$  >

SC set of security classes

 $\rightarrow \subseteq$  SC X SC flow relation (i.e., can-flow)

⊕: SC X SC -> SC class-combining operator

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# **DENNING'S AXIOMS**

< SC,  $\rightarrow$ ,  $\oplus$  >

- 1 SC is finite
- $2 \rightarrow is a partial order on SC$
- 3 SC has a lower bound L such that L  $\rightarrow$  A for all A  $\in$  SC
- 4 ⊕ is a least upper bound (lub) operator on SC

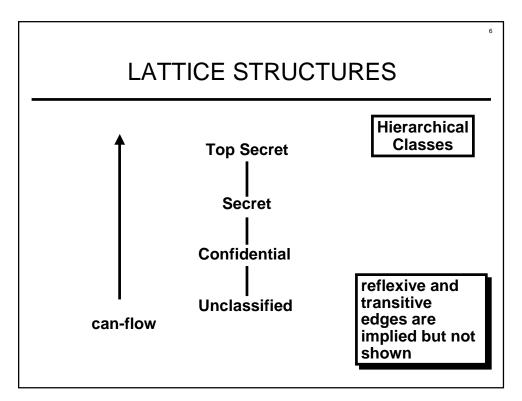
Justification for 1 and 2 is stronger than for 3 and 4. In practice we may therefore end up with a partially ordered set (poset) rather than a lattice.

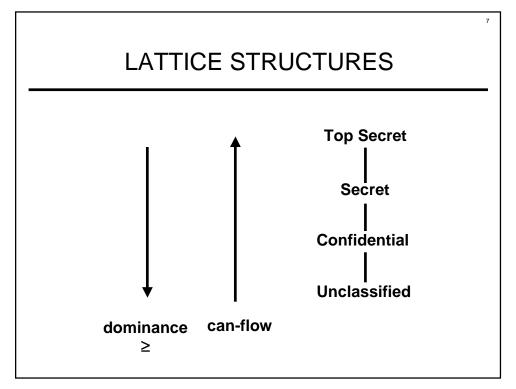
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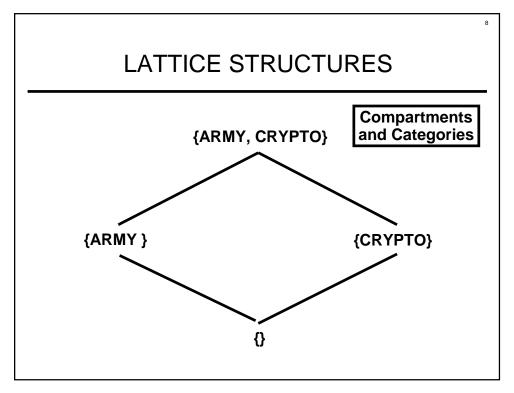
**DENNING'S AXIOMS IMPLY** 

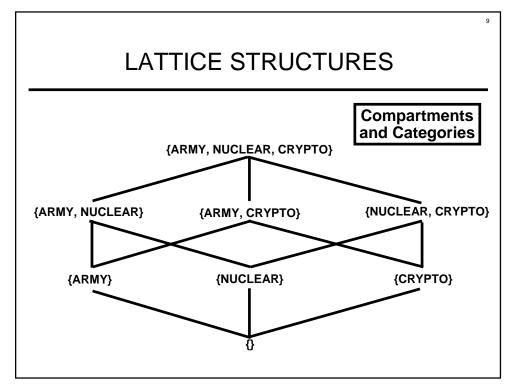
- SC is a universally bounded lattice
- there exists a Greatest Lower Bound (glb) operator ⊗ (also called meet)
- there exists a highest security class H

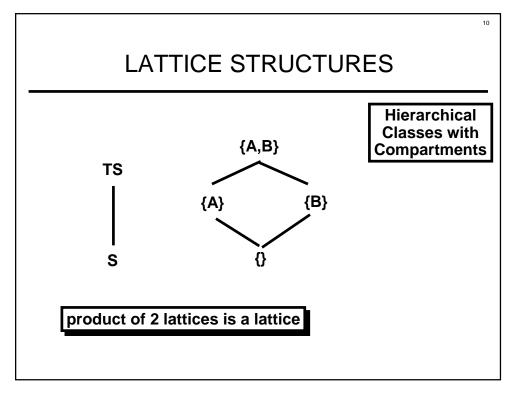
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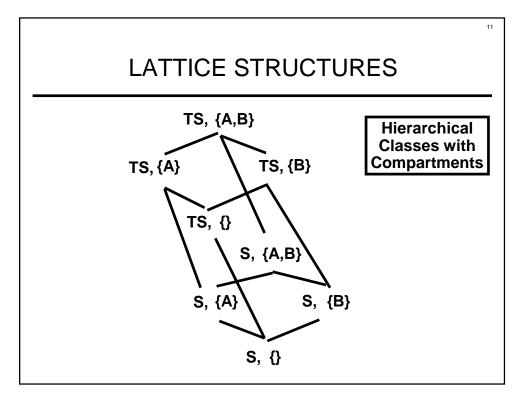




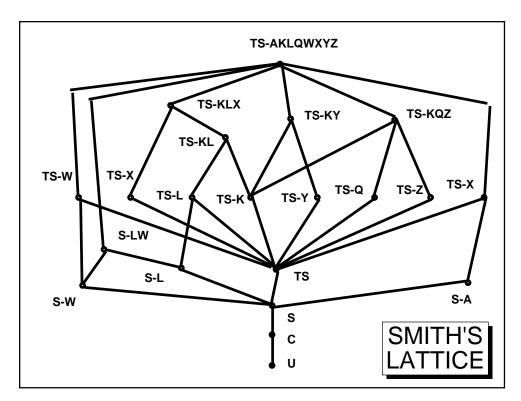








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#### SMITH'S LATTICE

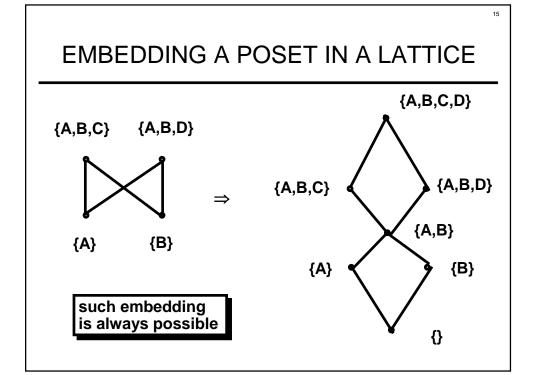
- With large lattices a vanishingly small fraction of the labels will actually be used
  - Smith's lattice: 4 hierarchical levels, 8 compartments, therefore
     number of possible labels = 4\*2^8 = 1024
     Only 21 labels are actually used (2%)
  - Consider 16 hierarchical levels, 64 compartments which gives 10<sup>2</sup>0 labels

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#### **EMBEDDING A POSET IN A LATTICE**

- Smith's subset of 21 labels do form a lattice. In general, however, selecting a subset of labels from a given lattice
  - may not yield a lattice, but
  - is guaranteed to yield a partial ordering
- Given a partial ordering we can always add extra labels to make it a lattice



# **BLP BASIC ASSUMPTIONS**

- SUB = {S1, S2, ..., Sm}, a fixed set of subjects
- OBJ = {O1, O2, ..., On}, a fixed set of objects
- $R \supseteq \{r, w\}$ , a fixed set of rights
- D, an m × n discretionary access matrix with D[i,j] ⊆ R
- M, an  $m \times n$  current access matrix with  $M[i,j] \subseteq \{r, w\}$

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#### **BLP MODEL**

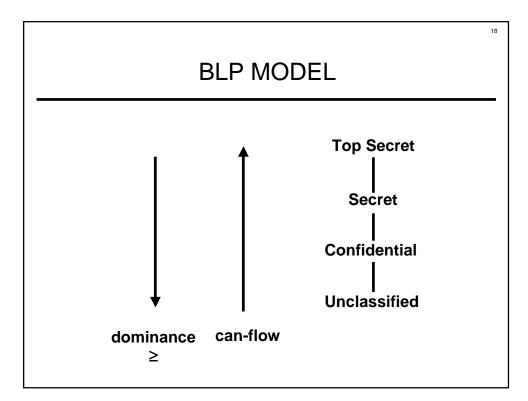
· Lattice of confidentiality labels

$$\Lambda = \{\lambda 1, \lambda 2, ..., \lambda p\}$$

• Static assignment of confidentiality labels

$$\lambda$$
: SUB  $\cup$  OBJ  $\rightarrow \Lambda$ 

- M, an m x n current access matrix with
  - $r \in M[i,j] \Rightarrow r \in D[i,j] \land \lambda(Si) \ge \lambda(Oj)$  simple security
  - $w \in M[i,j] \Rightarrow w \in D[i,j] \land \lambda(Si) \le \lambda(Oj)$  star-property



#### STAR-PROPERTY

- applies to subjects not to users
- users are trusted (must be trusted) not to disclose secret information outside of the computer system
- subjects are not trusted because they may have Trojan Horses embedded in the code they execute
- star-property prevents overt leakage of information and does not address the covert channel problem

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# **BIBA MODEL**

Lattice of integrity labels

$$\Omega = \{\omega_1, \omega_2, ..., \omega_q\}$$

Assignment of integrity labels

$$\omega$$
: SUB  $\cup$  OBJ  $\rightarrow \Omega$ 

- M, an m × n current access matrix with
  - $r \in M[i,j] \Rightarrow r \in D[i,j] \land \omega(Si) \le \omega(Oj)$  simple integrity
  - $w \in M[i,j] \Rightarrow w \in D[i,j] \land \omega(Si) \ge \omega(Oj)$  integrity confinement

#### EQUIVALENCE OF BLP AND BIBA

- Information flow in the Biba model is from top to bottom
- Information flow in the BLP model is from bottom to top
- Since top and bottom are relative terms, the two models are fundamentally equivalent

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EQUIVALENCE OF BLP AND BIBA

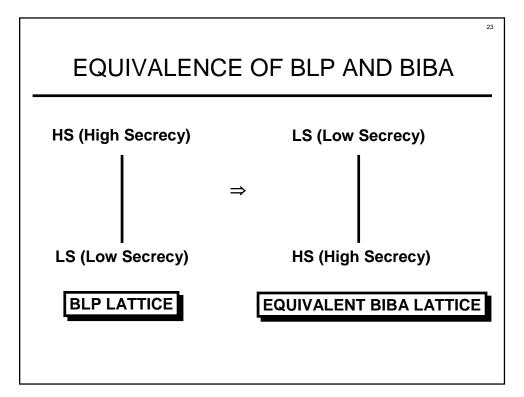
HI (High Integrity)

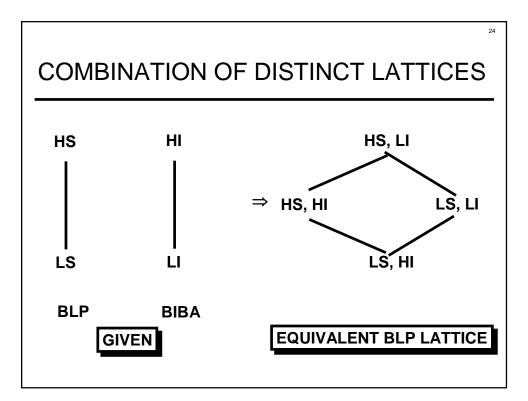
LI (Low Integrity)

LI (Low Integrity)

HI (High Integrity)

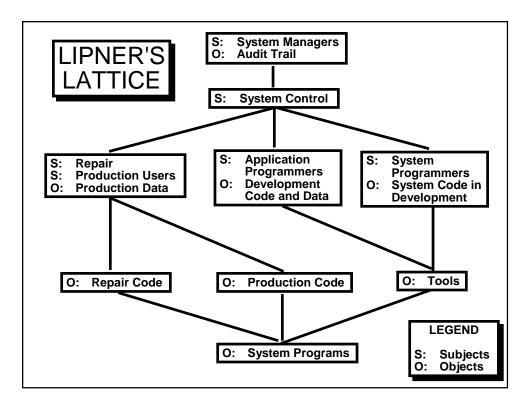
EQUIVALENT BLP LATTICE





#### **BLP AND BIBA**

- BLP and Biba are fundamentally equivalent and interchangeable
- Lattice-based access control is a mechanism for enforcing one-way information flow, which can be applied to confidentiality or integrity goals
- We will use the BLP formulation with high confidentiality at the top of the lattice, and high integrity at the bottom



#### LIPNER'S LATTICE

- Lipner's lattice uses 9 labels from a possible space of 192 labels (3 integrity levels, 2 integrity compartments, 2 confidentiality levels, and 3 confidentiality compartments)
- The single lattice shown here can be constructed directly from first principles

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## LIPNER'S LATTICE

- The position of the audit trail at lowest integrity demonstrates the limitation of an information flow approach to integrity
- System control subjects are exempted from the star-property and allowed to
  - write down (with respect to confidentiality) or equivalently
  - write up (with respect to integrity)

## DYNAMIC LABELS IN BLP

- Tranquility (most common):
   λ is static for subjects and objects
- BLP without tranquility may be secure or insecure depending upon the specific dynamics of labelling
- Noninterference can be used to prove the security of BLP with dynamic labels

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## DYNAMIC LABELS IN BLP

- High water mark on subjects:
   λ is static for objects
  - $\lambda$  may increase but not decrease for subjects

Is secure and is useful

- High water mark on objects:
  - $\lambda$  is static for subjects
  - $\boldsymbol{\lambda}$  may increase but not decrease for subjects

Is insecure due to disappearing object signaling channel

# **REFERENCES**

Ravi Sandhu, "Lattice-Based Access Control Models."
 Manuscript handed out in class