

Group-Centric Models for Secure and Agile Information Sharing

Ravi Sandhu
Executive Director and Endowed Professor
September 2010

ravi.sandhu@utsa.edu, www.profsandhu.com, www.ics.utsa.edu

Joint work with ICS colleagues
Ram Krishnan, Jianwei Niu and Will Winsborough

- 3 successful access control models in 40+ years
 - ❖ Discretionary Access Control (DAC)
 - ❖ Mandatory Access Control (MAC)
also called Lattice-Based Access Control (LBAC)
 - ❖ Role-base Access Control (RBAC)
- Numerous others defined and studied, implemented but no success
- Will Group Centric Models be the 4th element?
 - ❖ Strong mathematical foundations
 - ❖ Strong intuitive foundations
 - ❖ Significant real-world deployment

Goal: Share but protect

➤ Containment challenge

❖ Client containment

- High assurance infeasible (e.g., cannot close the analog hole)
- Low to medium assurance achievable

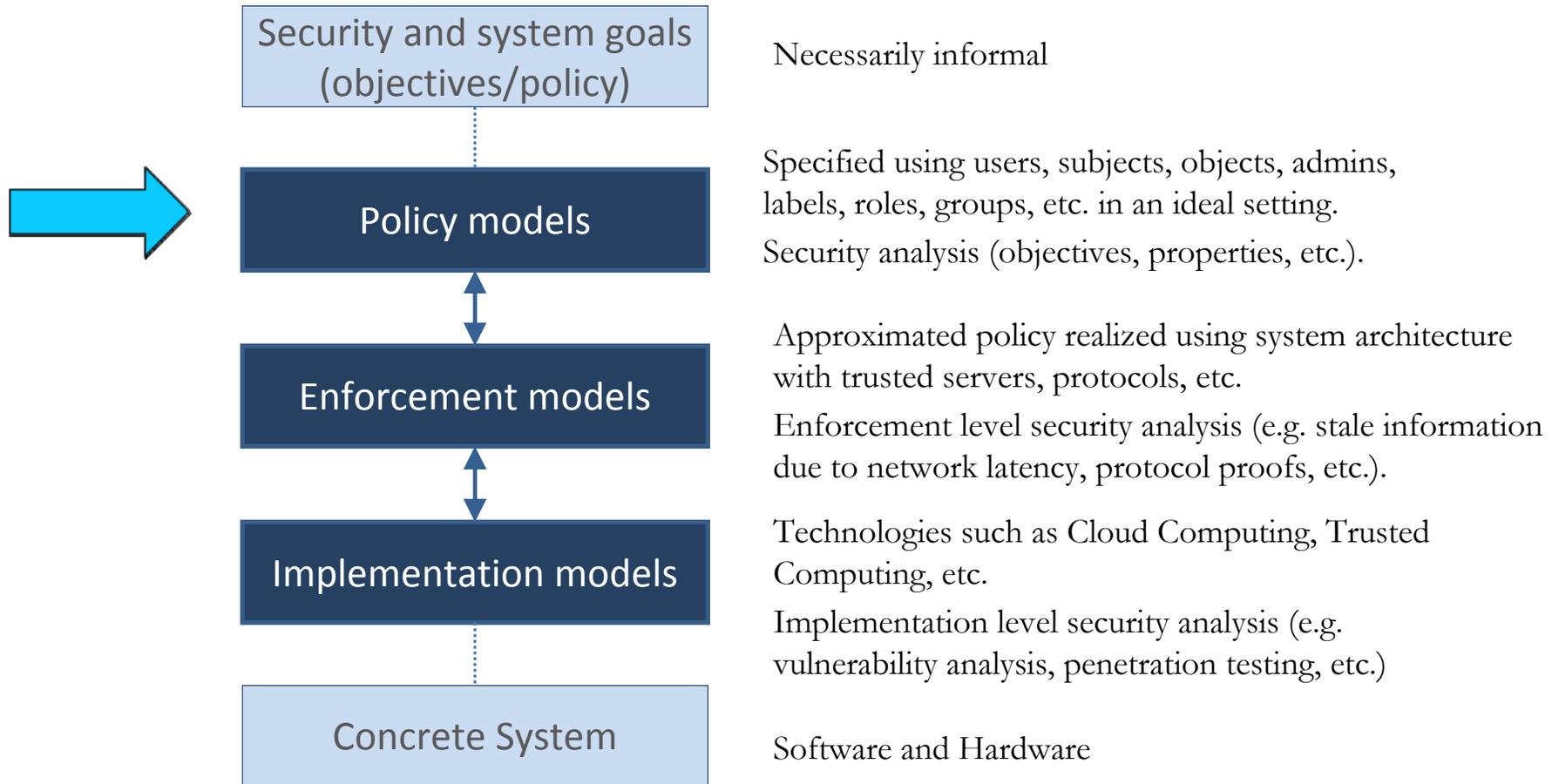
❖ Server containment

- Will typically have higher assurance than client containment

➤ Policy challenge

❖ How to construct meaningful, usable SIS policy

❖ How to develop an intertwined information and security model



Fundamental Goal: Share BUT Protect

I. Dissemination-Centric Sharing

- Digital Rights Management
- Enterprise Rights Management
- XrML
- Workflow-centric sharing

II. Query-Centric Sharing

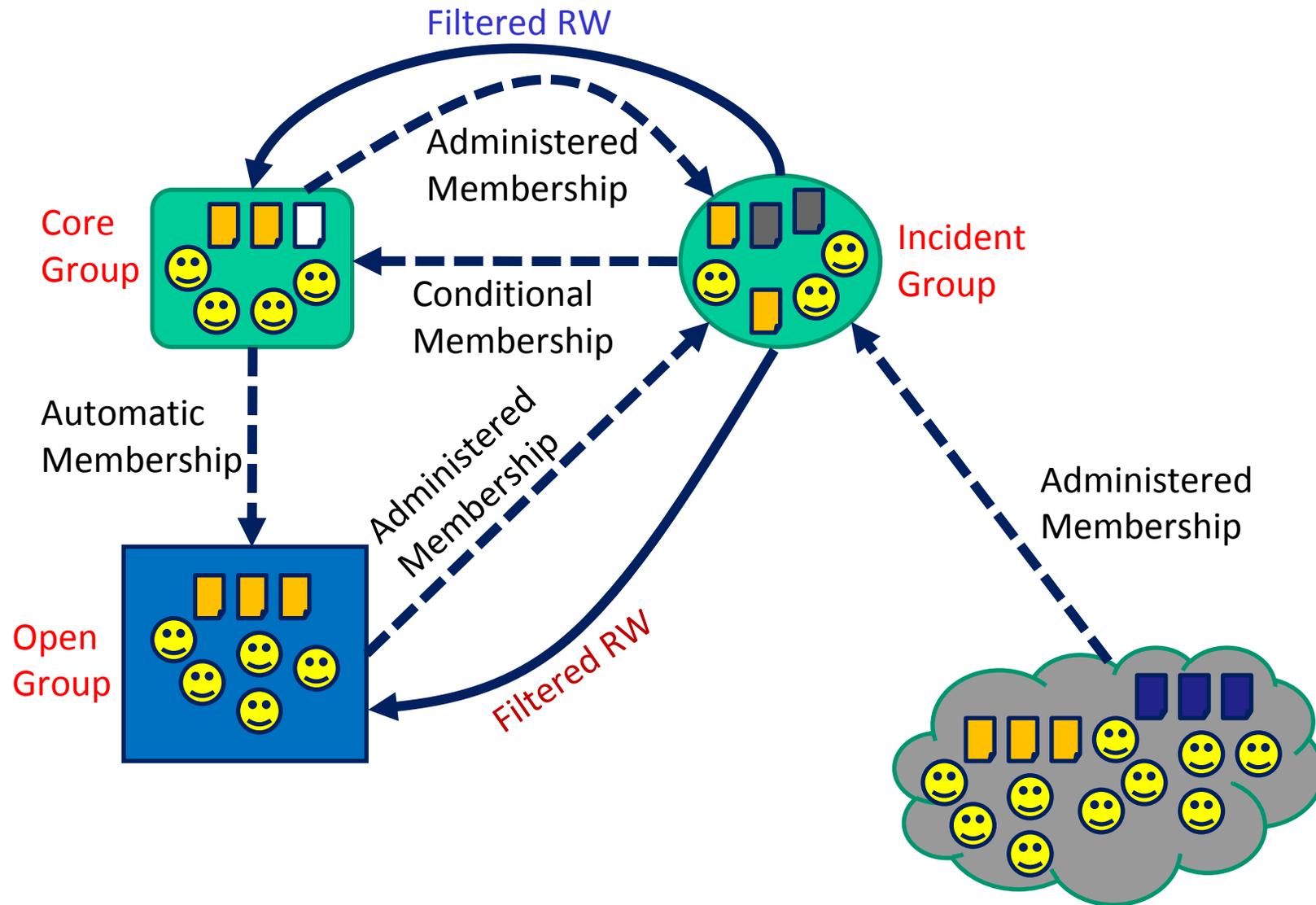
- Queries wrt a protected dataset
- Privacy/confidentiality protection
- More generally de-aggregation/inference protection

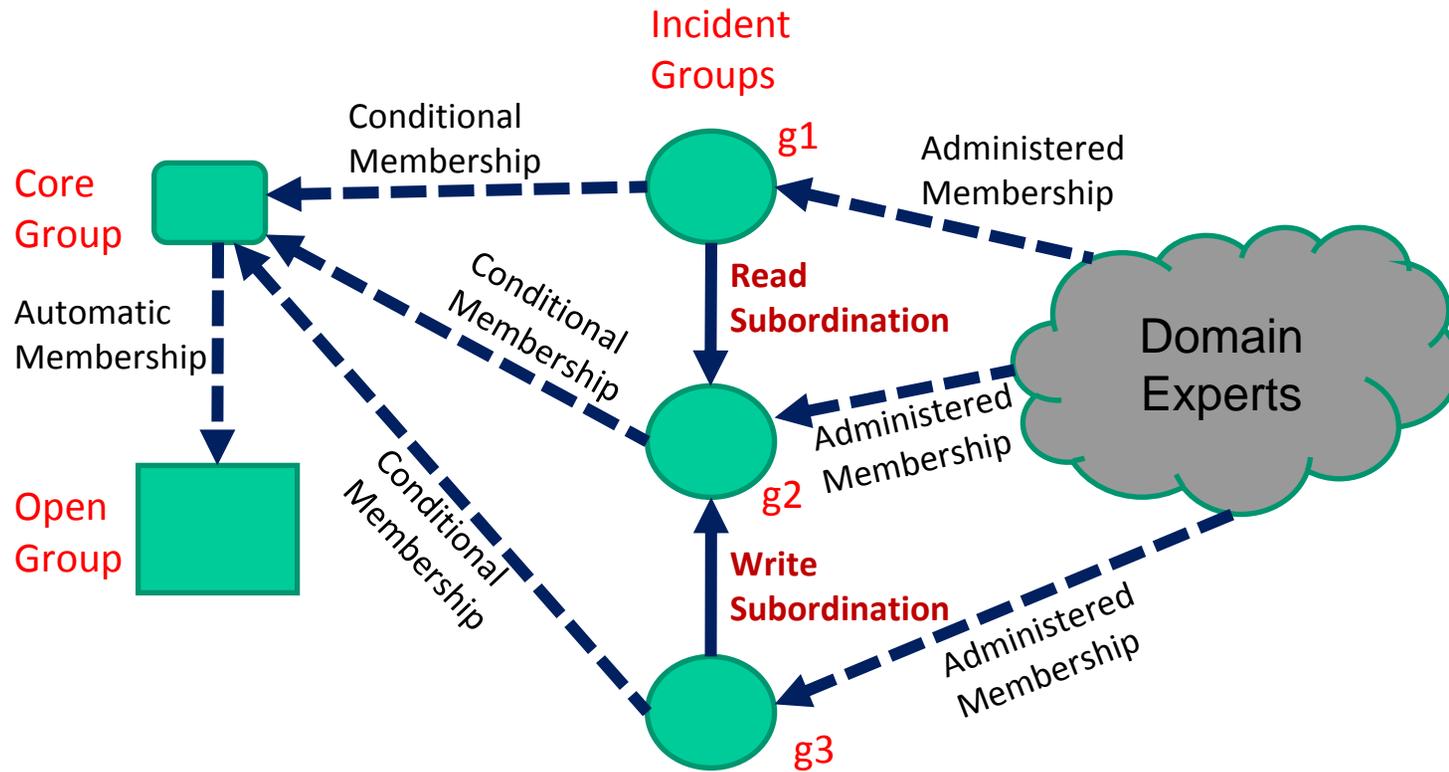
III. Group-Centric Sharing

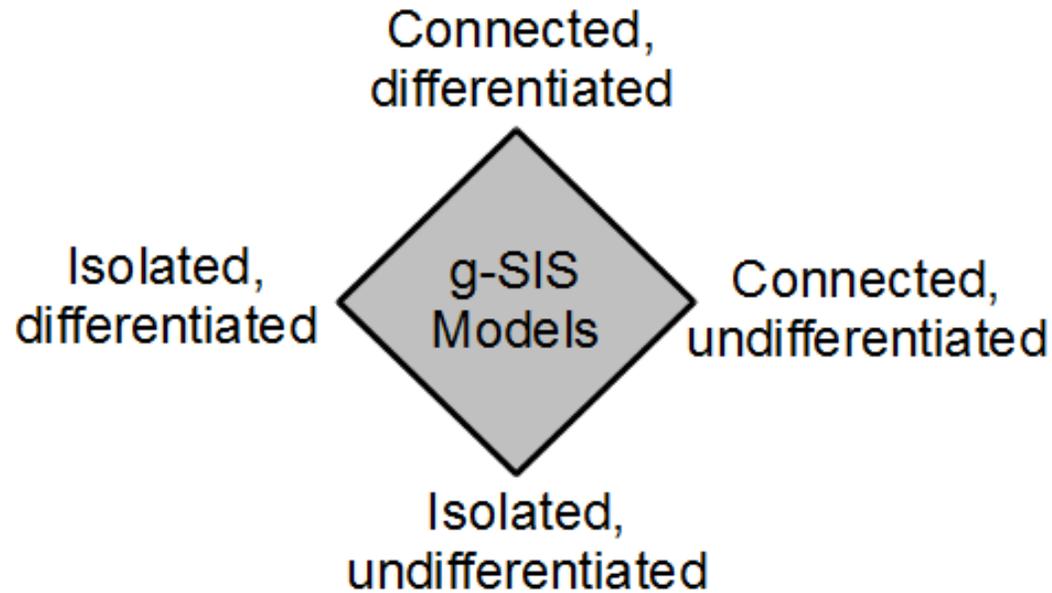
- Sharing for a purpose
- Mission-centric sharing
- Purpose-centric sharing

- A community is a county or larger city size unit
 - ❖ Clearly demarcated geographical boundary
 - ❖ More or less aligned with governance boundary

- The ICS Center for Infrastructure Assurance and Security has a decade long experience conducting cyber security exercises and training for communities all across USA
 - ❖ Community cyber security incident life cycle







- Formal stateless behavioral model with
 - ❖ Provable security properties
- Formal stateful enforceable model with
 - ❖ Proof of correspondence between stateless and stateful models

➤ Operational aspects

❖ Group operation semantics

- Add, Join, Leave, Remove, etc
- Multicast group is one example

❖ Object model

- Read-only
- Read-Write (no versioning vs versioning)

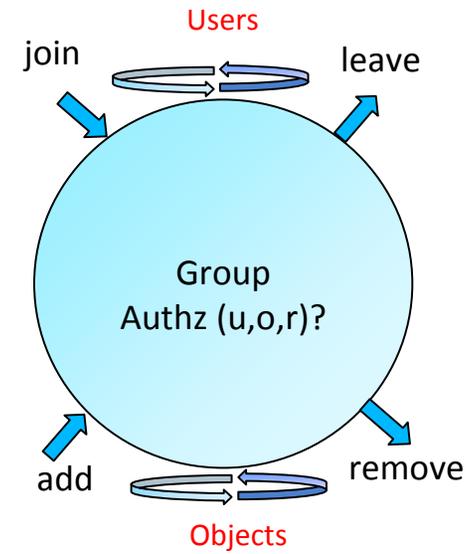
❖ User-subject model

- Read-only Vs read-write

❖ Policy specification

➤ Administrative aspects

- ❖ Authorization to create group, user join/leave, object add/remove, etc.



➤ Authorization Persistence

- ❖ *Authorization cannot change unless some group event occurs*

$$\kappa_0 = \forall u : U. \forall o : O. \forall v : V. \forall g : G.$$

$$\square(\text{Authz}(u, o, v, g, \mathbf{r}) \rightarrow (\text{Authz}(u, o, v, g, \mathbf{r}) \mathcal{W} (\text{Join}(u, g) \vee \text{Leave}(u, g) \vee \text{Add}(o, v, g) \vee \text{Remove}(o, v, g))))))$$

$$\kappa_1 = \forall u : U. \forall o : O. \forall v : V. \forall g : G.$$

$$\square(\text{Authz}(u, o, v, g, \mathbf{w}) \rightarrow (\text{Authz}(u, o, v, g, \mathbf{w}) \mathcal{W} \text{Leave}(u, g)))$$

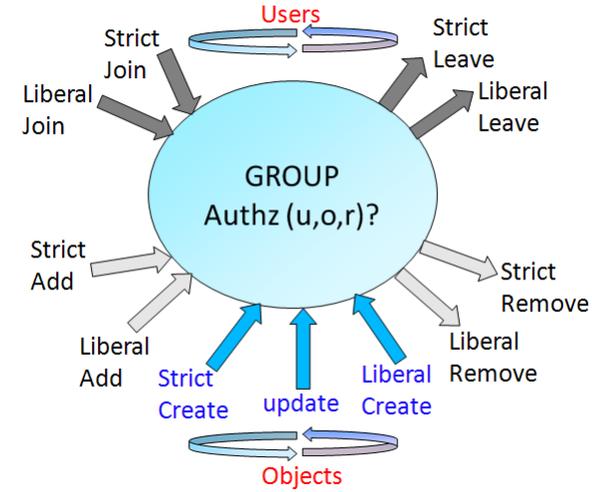
$$\kappa_2 = \forall u : U. \forall o : O. \forall v_1 : V. \forall g : G. \exists s : S. \exists v_2 : V.$$

$$\square(\neg \text{Authz}(u, o, v_1, g, \mathbf{r}) \rightarrow (\neg \text{Authz}(u, o, v_1, g, \mathbf{r}) \mathcal{W} (\text{Join}(u, g) \vee \text{Leave}(u, g) \vee \text{Add}(o, v_1, g) \vee \text{Remove}(o, v_1, g) \vee \text{CreateO}(o, v_1, g) \vee \text{update}(s, o, v_2, v_1, g))))))$$

$$\kappa_3 = \forall u : U. \forall o : O. \forall v_1 : V. \forall g : G. \exists s : S. \exists v_2 : V.$$

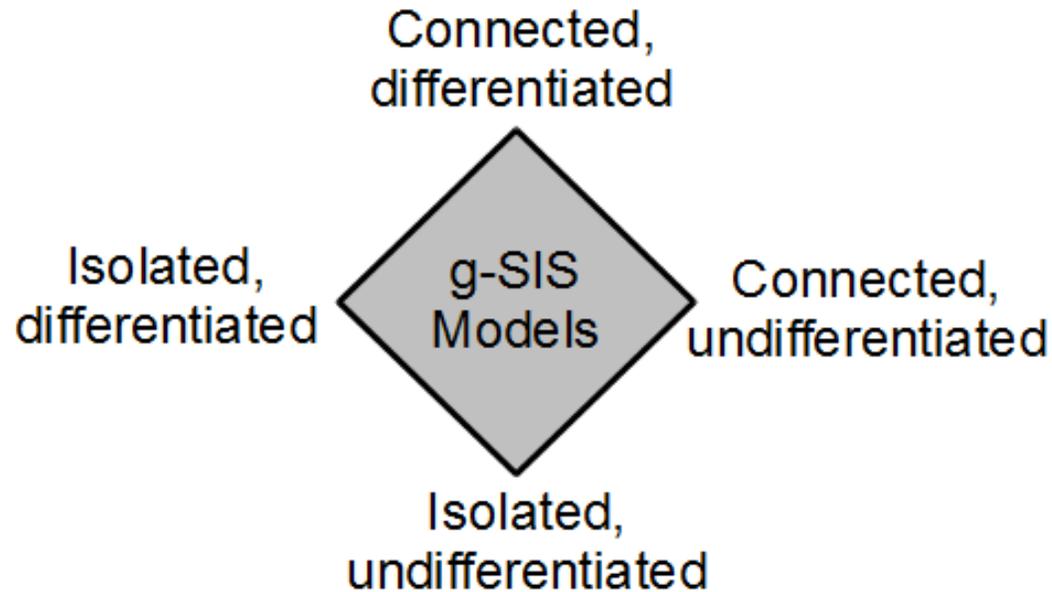
$$\square(\neg \text{Authz}(u, o, v_1, g, \mathbf{w}) \rightarrow (\neg \text{Authz}(u, o, v_1, g, \mathbf{w}) \mathcal{W} (\text{Join}(u, g) \vee \text{CreateO}(o, v_1, g) \vee \text{update}(s, o, v_2, v_1, g))))))$$

Table 1: The π -system.

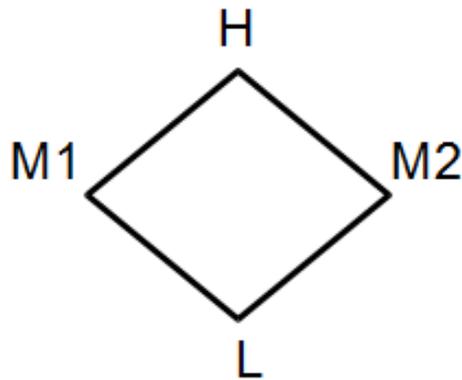
$$\begin{aligned} \chi_0 &= \forall u : U. \forall o : O. \forall v : V. \forall g : G. \\ &\quad \Box (\text{Authz}(u, o, v, g, \mathbf{r}) \leftrightarrow \exists v_1 : V. \exists s : S. (\lambda_0(u, o, v, g) \vee \dots \vee \lambda_3(u, s, o, v_1, v, g) \vee \\ &\quad \lambda'_0(u, o, v, g) \vee \dots \vee \lambda'_4(u, s, o, v_1, v, g))) \\ \chi_1 &= \forall u : U. \forall o : O. \forall v : V. \forall g : G. \\ &\quad \Box (\text{Authz}(u, o, v, g, \mathbf{w}) \leftrightarrow \text{Authz}(u, o, v, g, \mathbf{r}) \wedge (\neg \text{Leave}(u, g) \mathcal{S} \text{Join}(u, g)) \wedge \\ &\quad (\exists v_1 : V. \exists s : S. \blacklozenge \text{update}(s, o, v_1, v, g) \vee \blacklozenge (\text{LC}(o, v, g) \vee \text{SC}(o, v, g)))) \\ \chi_2 &= \forall u : U. \forall s : S. \forall g : G. \\ &\quad \Box (\text{createS}(u, s, g) \rightarrow \blacklozenge \text{Join}(u, g)) \\ \chi_3 &= \forall s : S. \forall o : O. \forall v : V. \forall g : G. \\ &\quad \Box (\text{AuthzS}(s, o, v, g, \mathbf{r}) \leftrightarrow \exists u : U. (\text{Authz}(u, o, v, g, \mathbf{r}) \wedge \\ &\quad (\neg \text{kill}(u, s, g) \mathcal{S} \text{createS}(u, s, g)))) \\ \chi_4 &= \forall s : S. \forall o : O. \forall v : V. \forall g : G. \\ &\quad \Box (\text{AuthzS}(s, o, v, g, \mathbf{w}) \leftrightarrow \exists u : U. (\text{Authz}(u, o, v, g, \mathbf{w}) \wedge \\ &\quad (\neg \text{kill}(u, s, g) \mathcal{S} \text{createS}(u, s, g)))) \\ \chi_5 &= \forall s : S. \forall o : O. \forall v_1, v_2 : V. \forall g : G. \\ &\quad \Box (\text{read}(s, o, v_1, g) \rightarrow \text{AuthzS}(s, o, v_1, g, \mathbf{r})) \wedge \\ &\quad \Box (\text{update}(s, o, v_1, v_2, g) \rightarrow \text{AuthzS}(s, o, v_1, g, \mathbf{w})) \\ \chi_6 &= \forall u_1, u_2 : U. \forall s_1, s_2 : S. \forall o : O. \forall v_1, v_2, v_3 : V. \forall g_1, g_2 : G. \\ &\quad \tau_0(u_1, s_1, s_2, o, v_1, v_2, v_3, g_1) \wedge \dots \wedge \tau_3(u_1, s_1, o, v_1, g_1) \end{aligned}$$


$$\pi = \chi_0 \wedge \chi_1 \wedge \chi_2 \wedge \chi_3 \wedge \chi_4 \wedge \chi_5 \wedge \chi_6$$

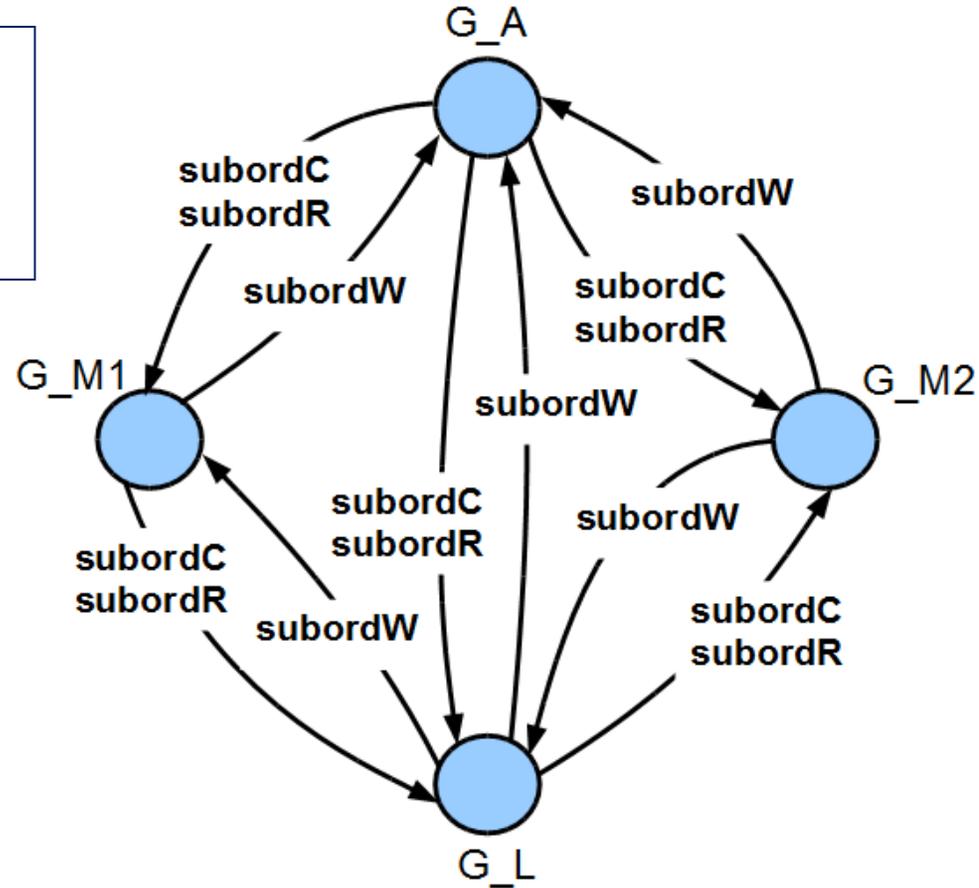
- Formal stateless behavioral model with
 - ❖ Provable security properties
- Formal stateful enforceable model with
 - ❖ Proof of correspondence between stateless and stateful models



1. Read Subordination
2. Write Subordination
3. Subject Create Subordination

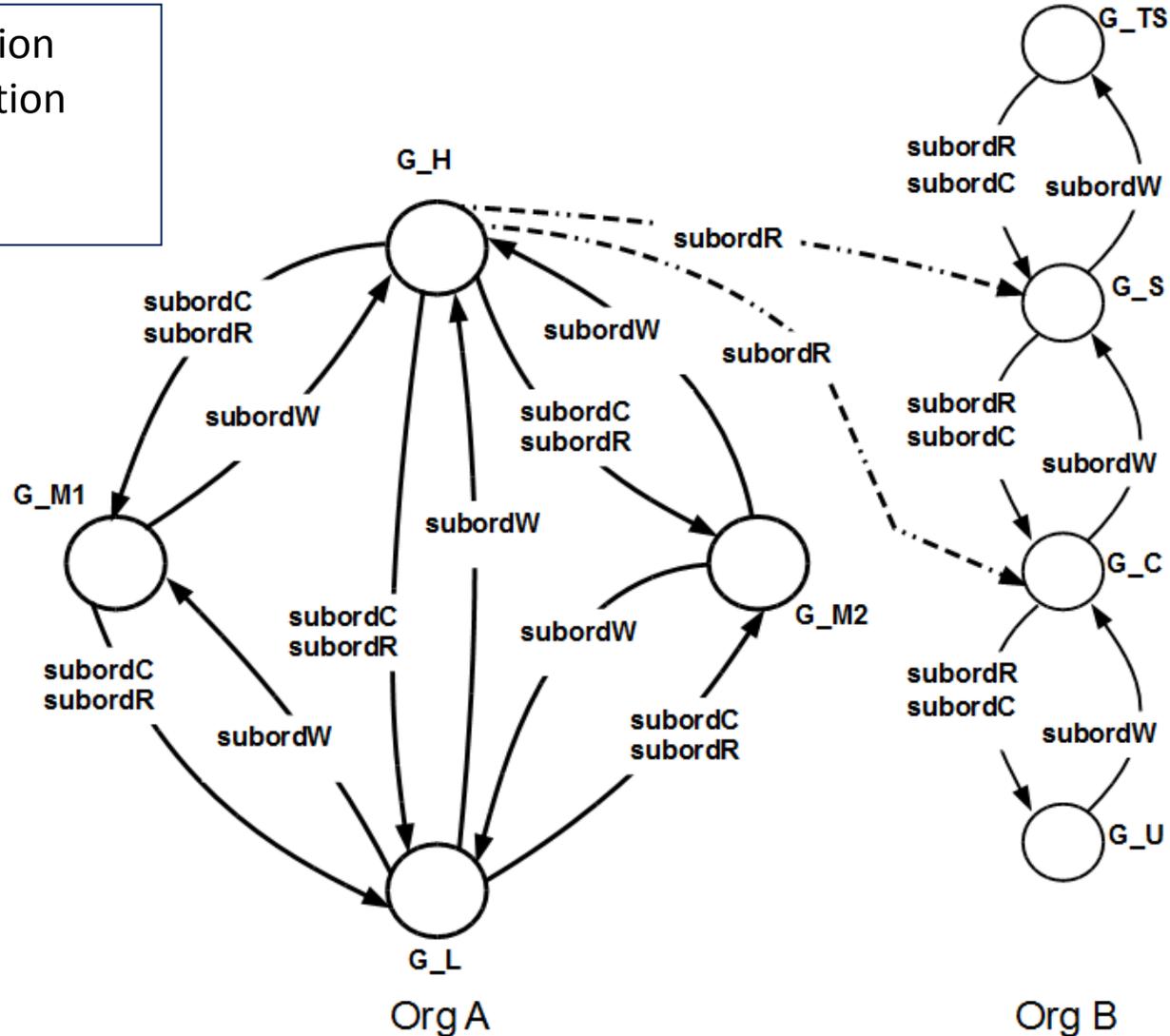


A sample lattice for one directional information flow



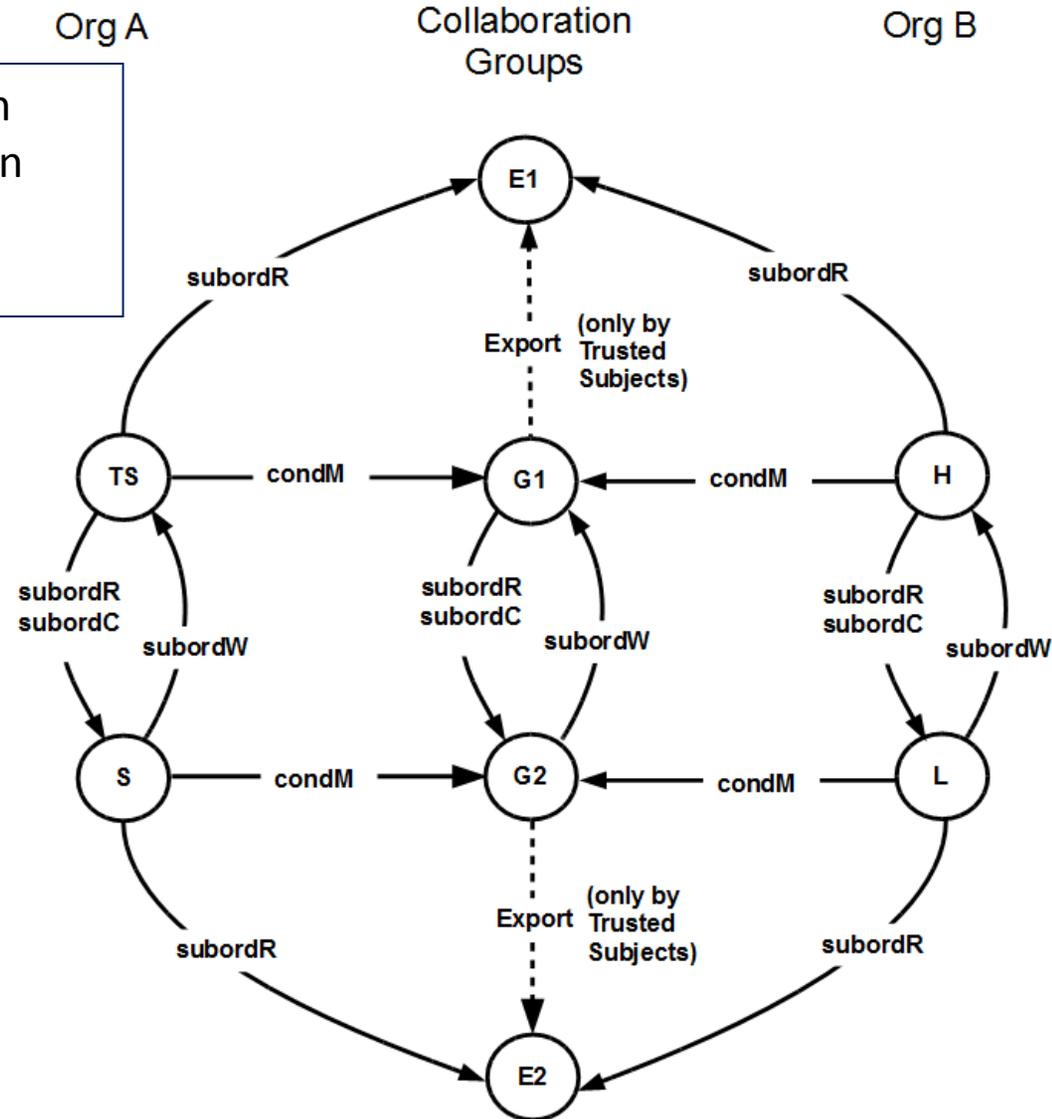
Equivalent g-SIS configuration of Org A lattice

1. Read Subordination
2. Write Subordination
3. Subject Create Subordination



Agile collaboration in LBAC enabled by g-SIS

1. Read Subordination
2. Write Subordination
3. Subject Create Subordination



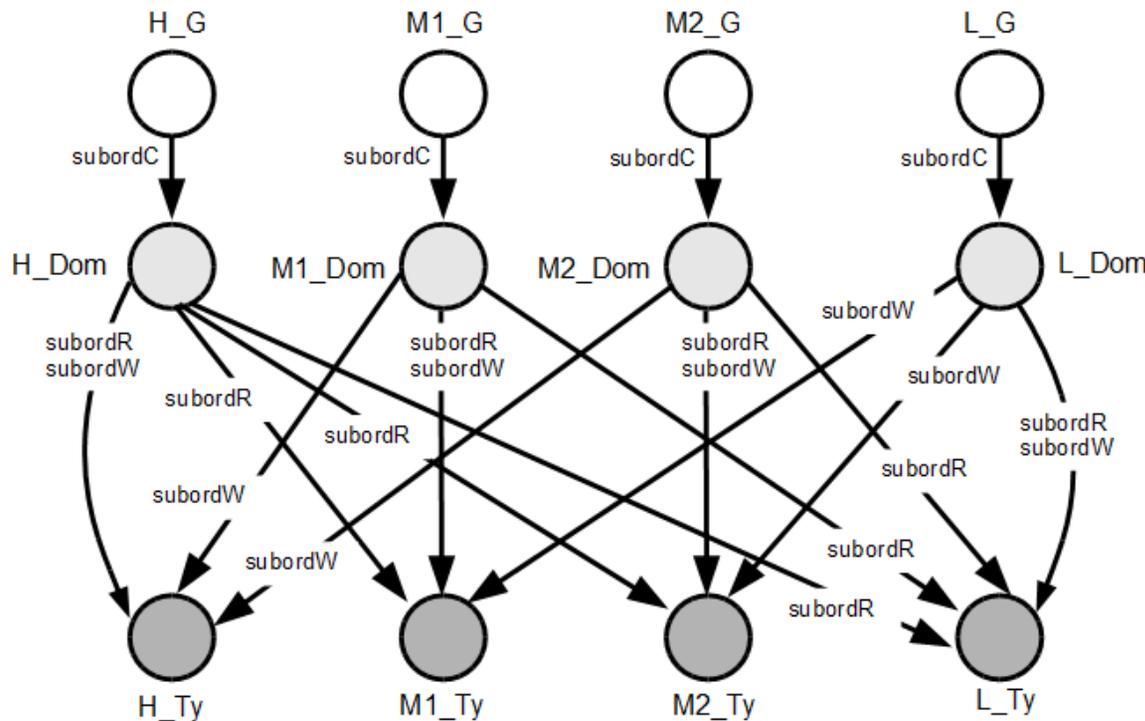
Collaboration groups established between two different lattices

Objects
→

Subjects
↓

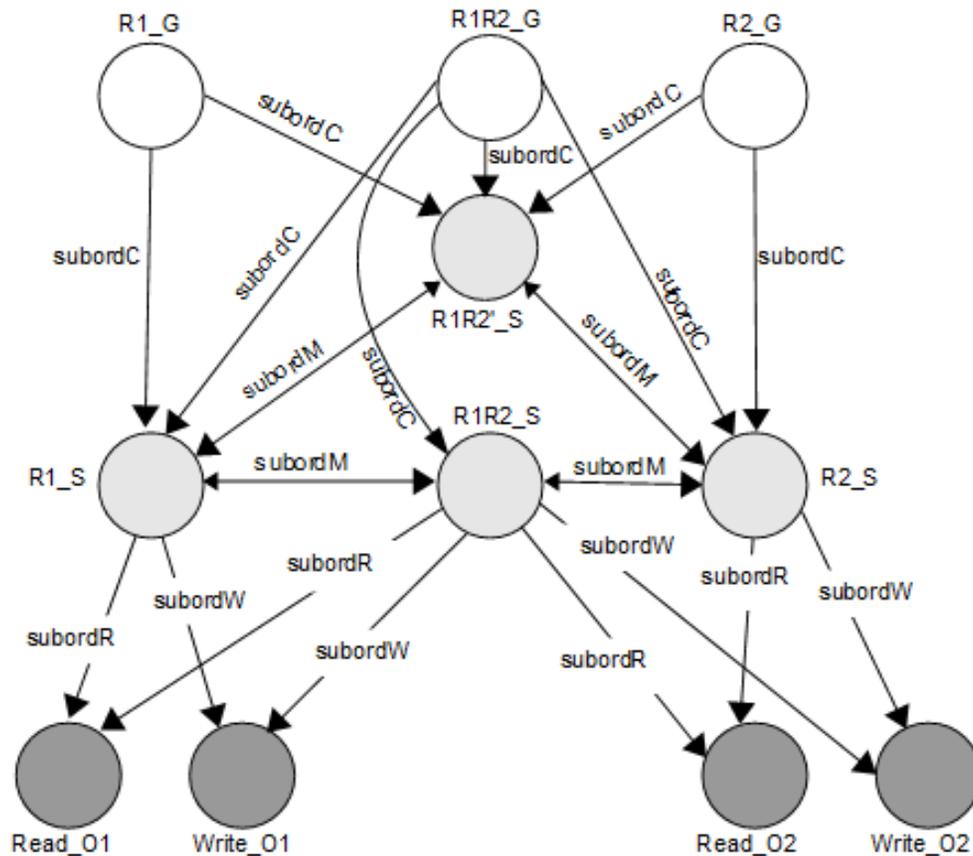
Domain \ Type	H_Ty	M1_Ty	M2_Ty	L_Ty
H_Dom	<i>rw</i>	<i>r</i>	<i>r</i>	<i>r</i>
M1_Dom	<i>w</i>	<i>rw</i>	-	<i>r</i>
M2_Dom	<i>w</i>	-	<i>rw</i>	<i>r</i>
L_Dom	-	<i>w</i>	<i>w</i>	<i>rw</i>

A sample DTE matrix



Equivalent g-SIS configuration

1. Read Subordination
2. Write Subordination
3. Subject Create Subordination



1. Read Subordination
2. Write Subordination
3. Subject Create Subordination
4. **Subject Move Subordination**

RBAC₀ with RW permissions in g-SIS

- 3 successful access control models in 40+ years
 - ❖ Discretionary Access Control (DAC)
 - ❖ Mandatory Access Control (MAC)
also called Lattice-Based Access Control (LBAC)
 - ❖ Role-base Access Control (RBAC)
- Numerous others defined and studied, implemented but no success
- Will Group Centric Models be the 4th element?
 - ❖ Strong mathematical foundations
 - ❖ Strong intuitive foundations
 - ❖ Significant real-world deployment