

CS 6393 Lecture 2



Access Control Models

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Information needs to be protected

- In motion
- ✤ At rest
- In use

> Absolute security is impossible and unnecessary

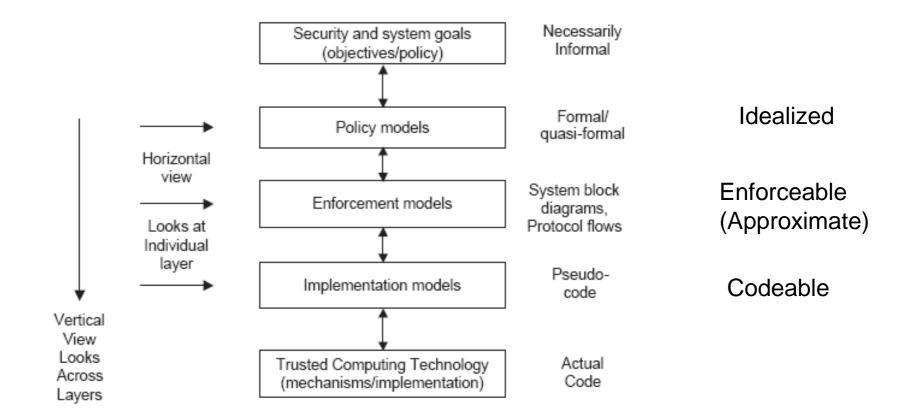
- Trying to approximate absolute security is a bad strategy
- Good enough" security is feasible and meaningful
- Better than "good enough" is bad
- Security is meaningless without application context
 Cannot know we have "good enough" without this context
- Models and abstractions are all important
 - Without a conceptual framework it is hard to separate "what needs to be done" from "how we do it"

We are not very good at doing any of this



PEI Models



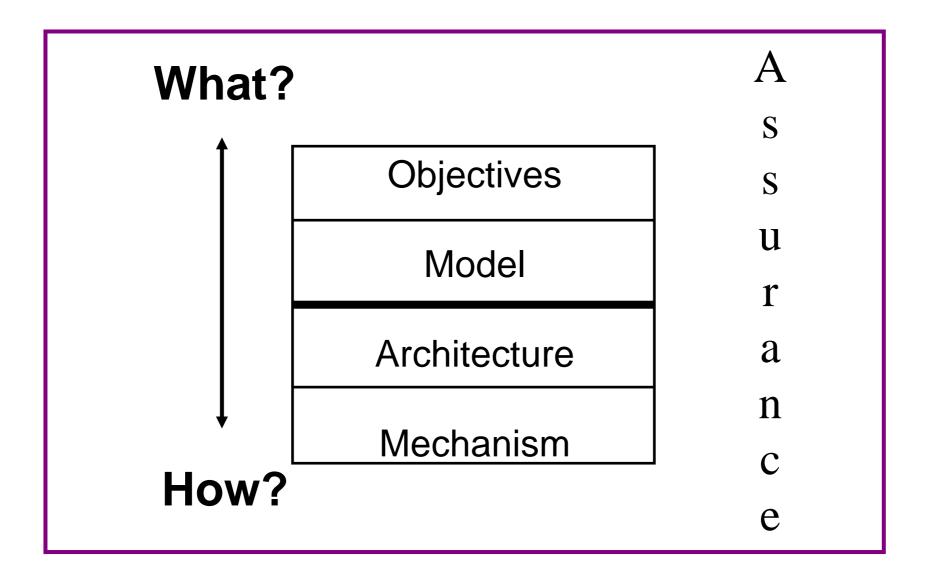


At the policy layer security models are essentially access control models



OM-AM (Older Version of PEI)





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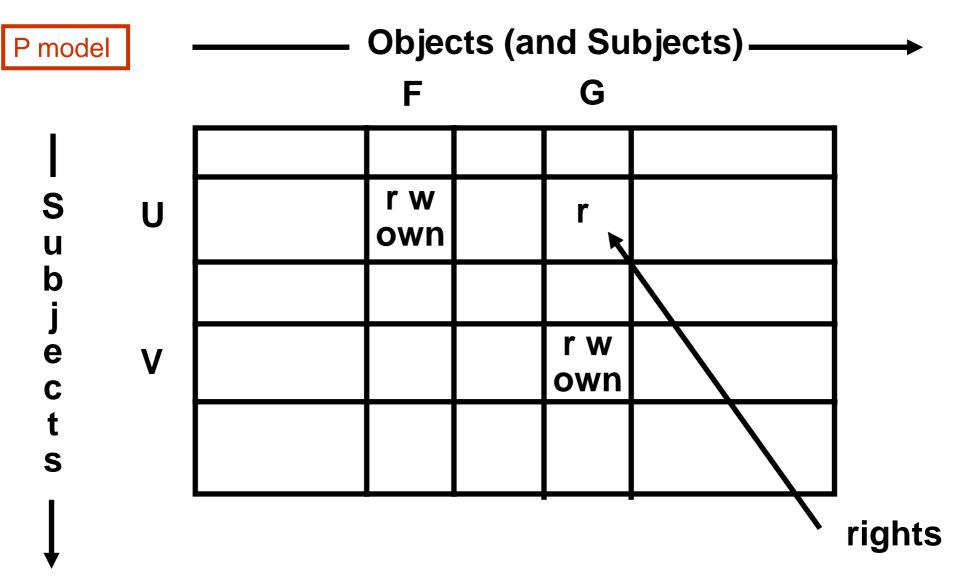




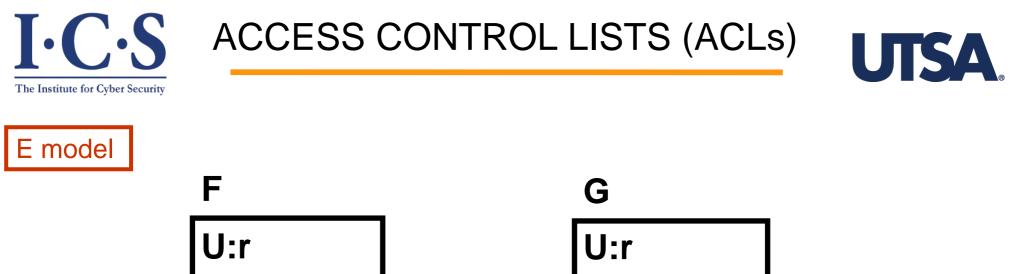
- Discretionary Access Control (DAC)
 - Owner controls access but only to the original, not to copies
- Mandatory Access Control (MAC)
 Same as Lattice-Based Access Control (LBAC)
 - Access based on security labels
 - Labels propagate to copies
- Role-Based Access Control (RBAC)
 - Access based on roles
 - Can be configured to do DAC or MAC
 - Generalizes to Attribute-Based Access Control (ABAC)

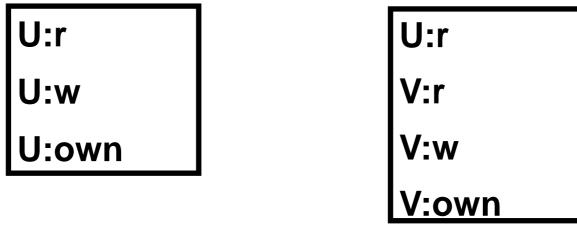
Numerous other models but only 3 successes





UTSA





each column of the access matrix is stored with the object corresponding to that column



CAPABILITY LISTS



E model

U F/r, F/w, F/own, G/r

V G/r, G/w, G/own

each row of the access matrix is stored with the subject corresponding to that row



ACCESS CONTROL TRIPLES



E model

Subject	Access	Object
U	r	F
U	W	F
U	own	F
U	r	G
V	r	G
V	W	G
V	own	G

commonly used in relational database management systems



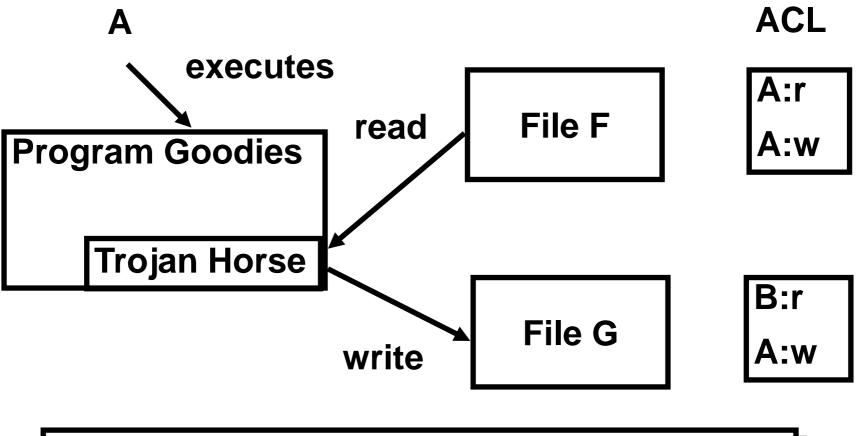




B cannot read file F

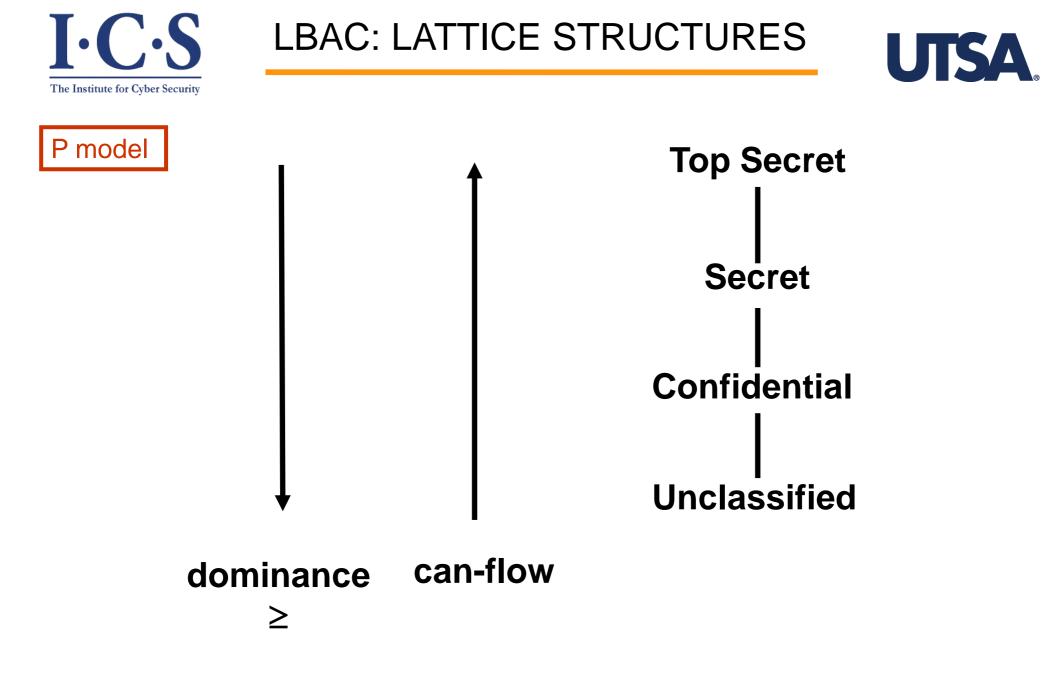


DAC: TROJAN HORSE EXAMPLE



B can read contents of file F copied to file G

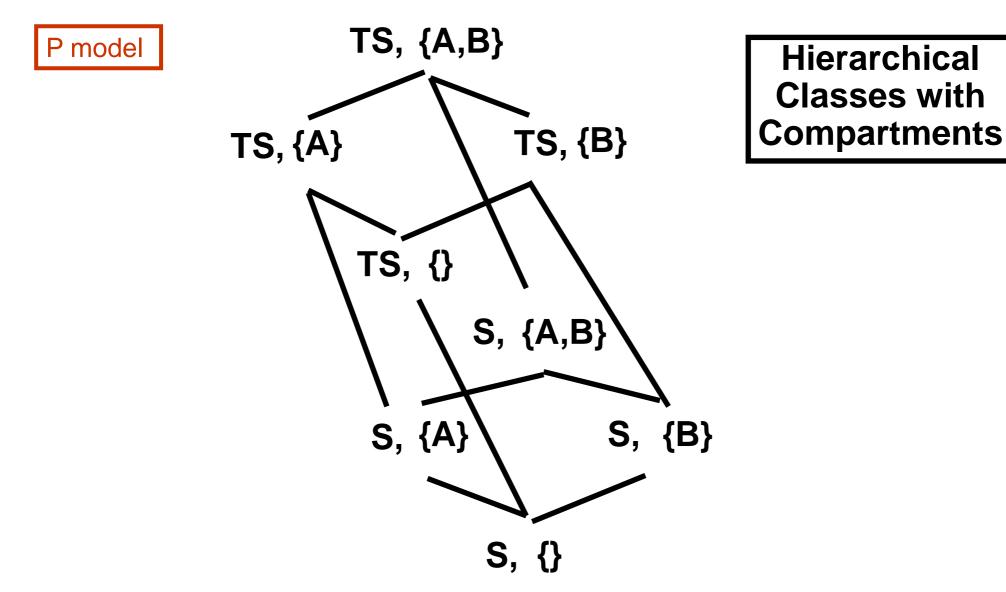
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LBAC: LATTICE STRUCTURES









SIMPLE-SECURITY

Subject S can read object O only if

label(S) dominates label(O)

STAR-PROPERTY (LIBERAL)

Subject S can write object O only if

label(O) dominates label(S)

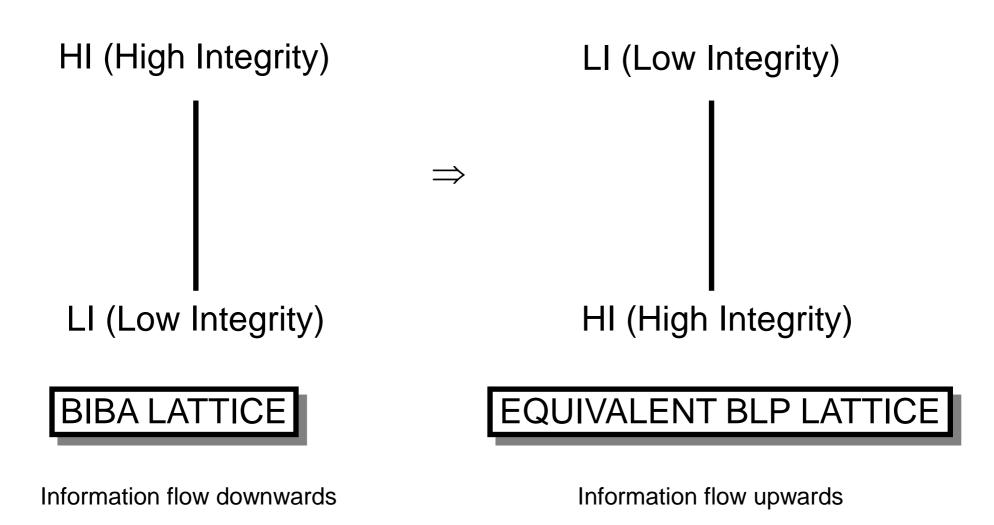
STAR-PROPERTY (STRICT)

Subject S can write object O only if

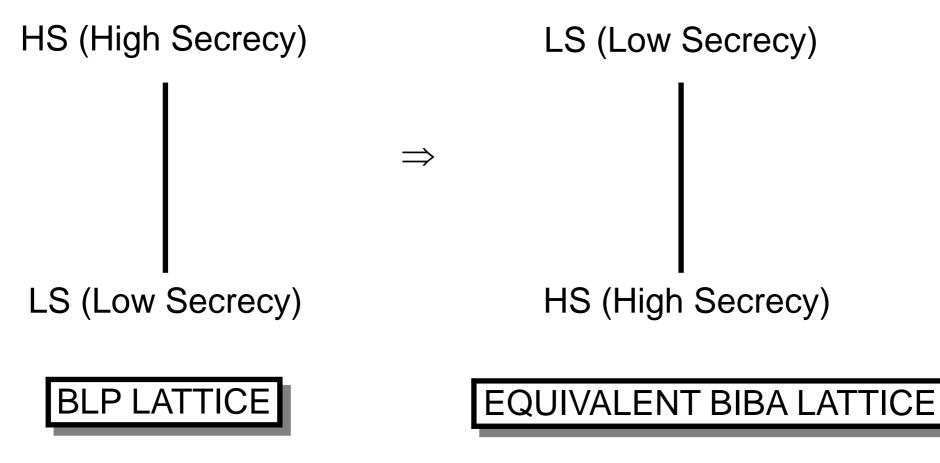
label(O) equals label(S)





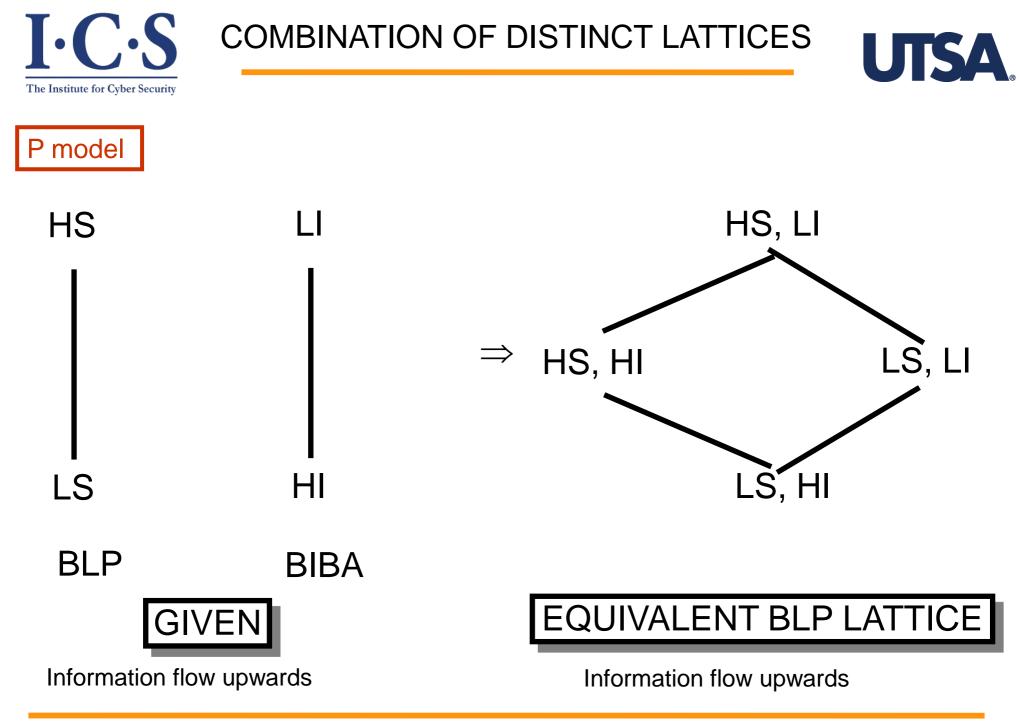






Information flow downwards

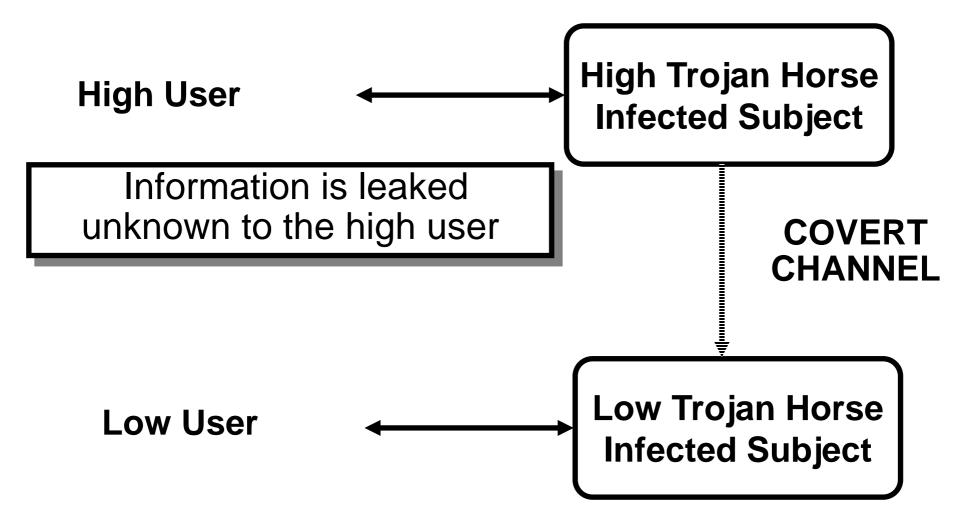
Information flow upwards





LBAC: LATTICE STRUCTURES









- Access is determined by roles
- A user's roles are assigned by security administrators

First emerged: mid 1970s First models: mid 1990s

 A role's permissions are assigned by security administrators

Is RBAC MAC or DAC or neither?

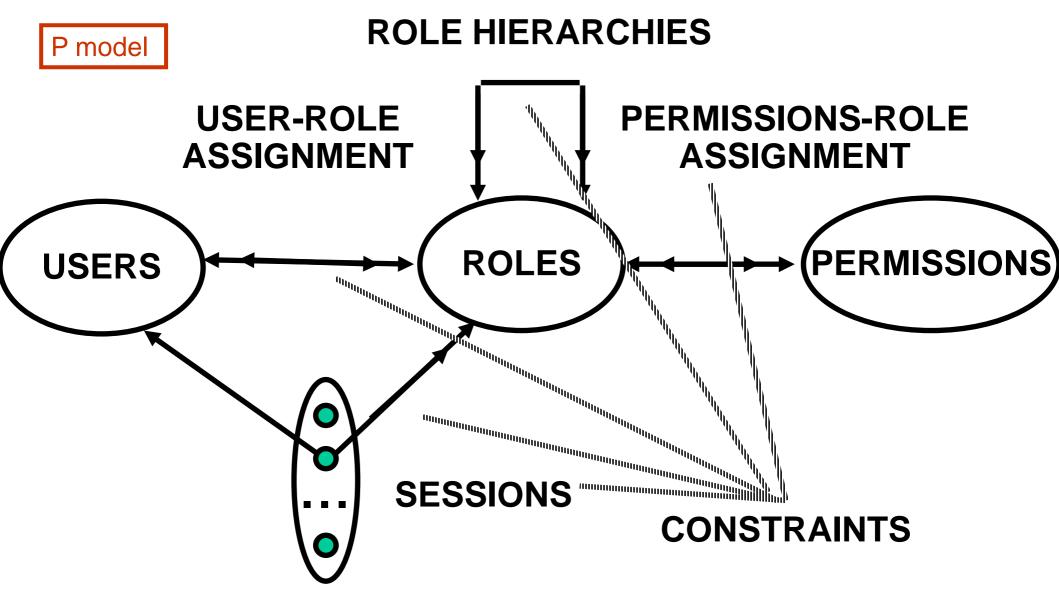
- RBAC can be configured to do MAC
- RBAC can be configured to do DAC
- RBAC is policy neutral

RBAC is neither MAC nor DAC!



RBAC: RBAC96 Model

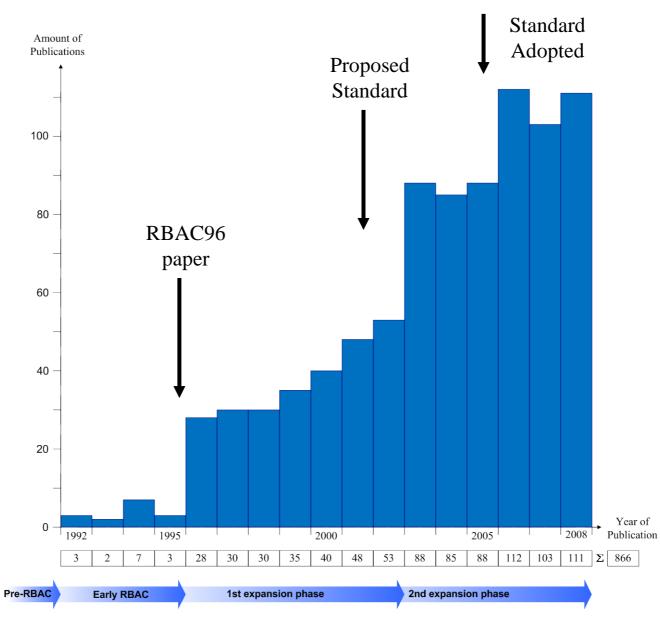






The RBAC Story



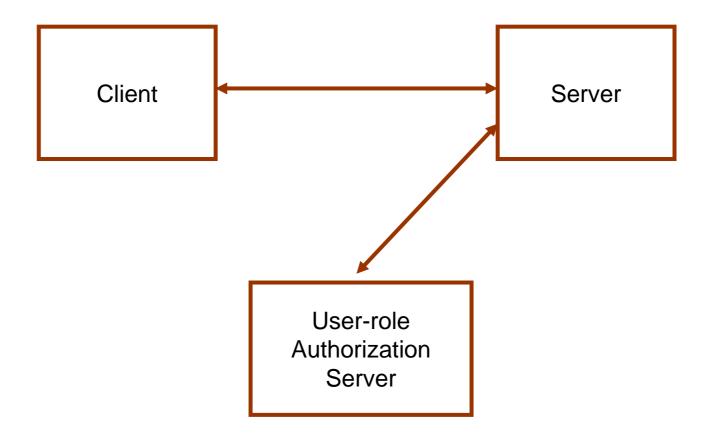




RBAC: SERVER PULL



E model

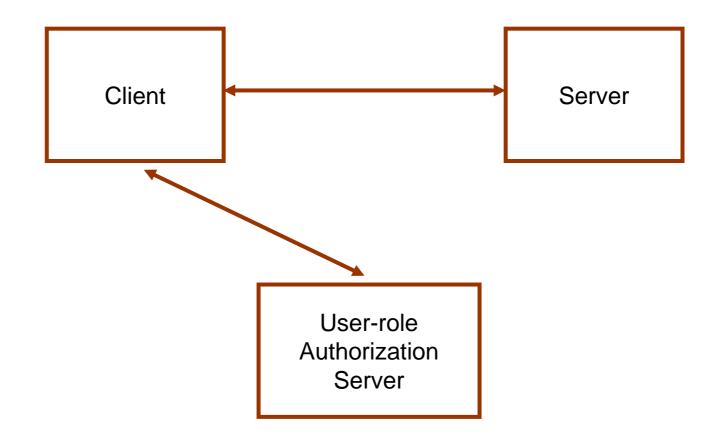




RBAC: CLIENT PULL



E model

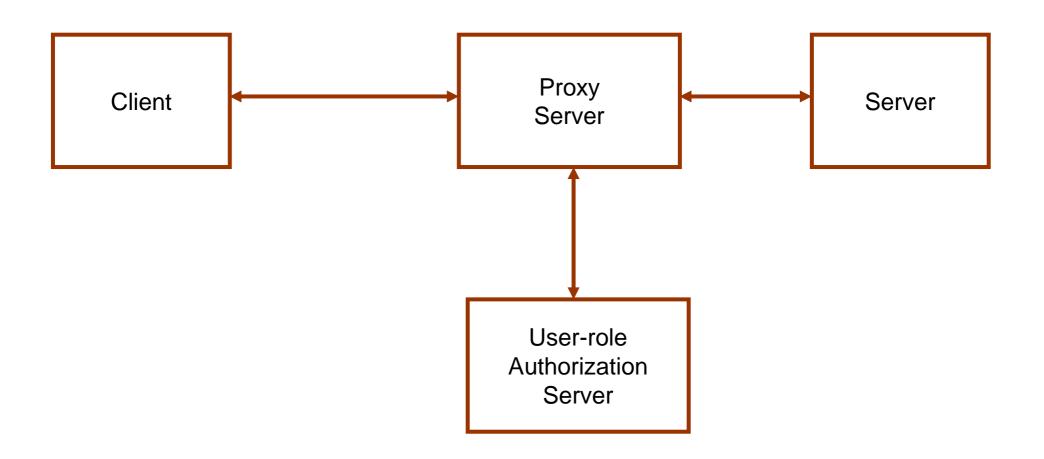




RBAC: PROXY-BASED



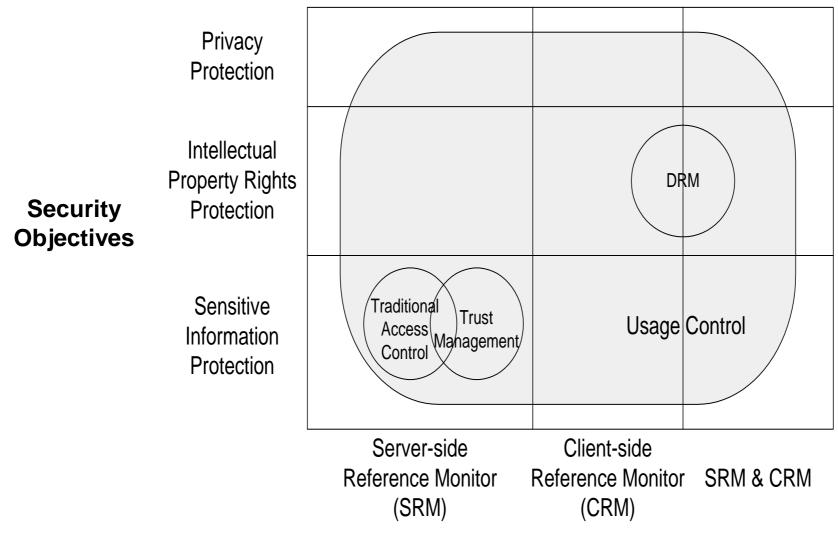
E model





UCON: Usage Control Scope



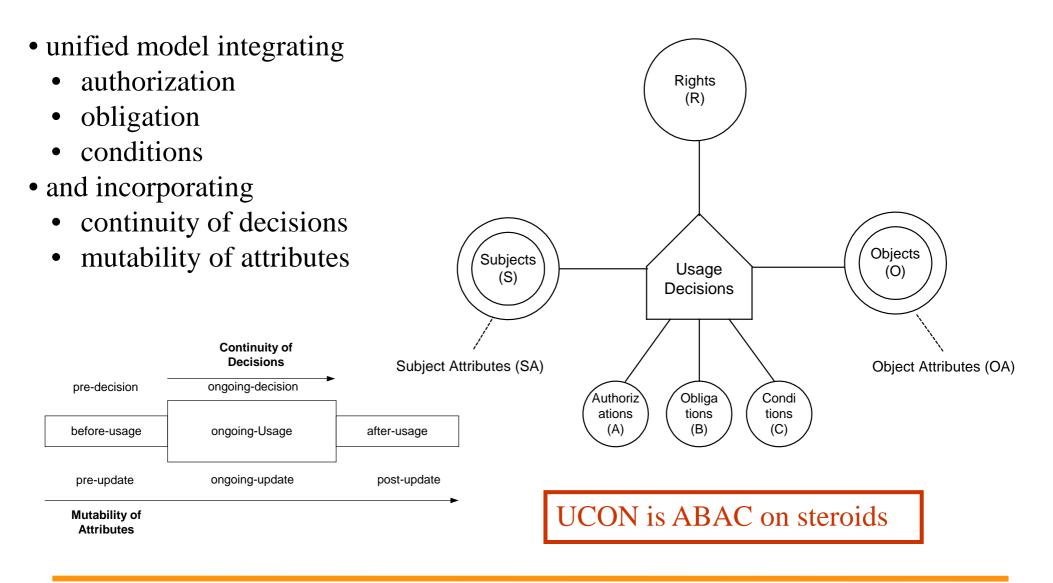


Security Architectures



UCON: Usage Control Model







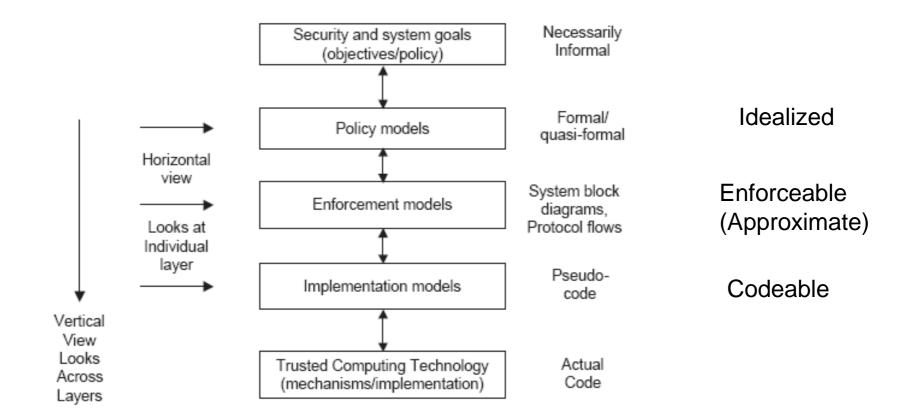


- Our Basic Premise
 - There can be no security model without application context
- So how does one customize an application-centric security model?
 - Meaningfully combine the essential insights of
 - > DAC, LBAC, RBAC, ABAC, UCON, etcetera
 - Directly address the application-specific trade-offs
 - Within the security objectives of confidentiality, integrity and availability
 - > Across security, performance, cost and usability objectives
 - Separate the real-world concerns of
 - > practical distributed systems and ensuing staleness and approximations (enforcement layer) from
 - > policy concerns in a idealized environment (policy layer)



PEI Models





At the policy layer security models are essentially access control models