

# Operation and Administration of Access Control in IoT Environments

## **Ph.D. Dissertation Defense**

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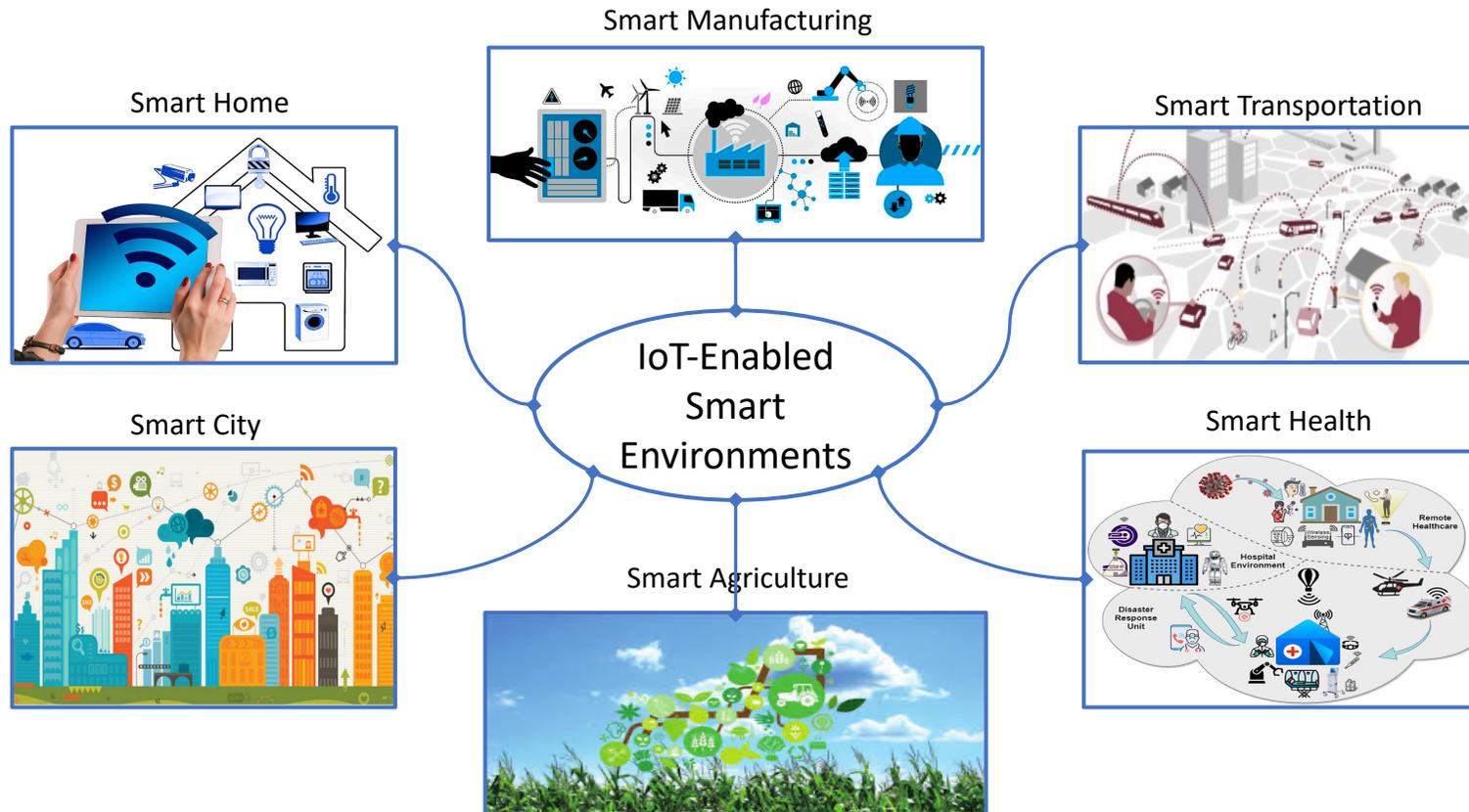
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**April 2022**

- The **Internet of Things (IoT)** denotes a network of evolving and expanding number of technologies embedded in **smart things** with at least one network interface to **connect, interact, and exchange data and information**.

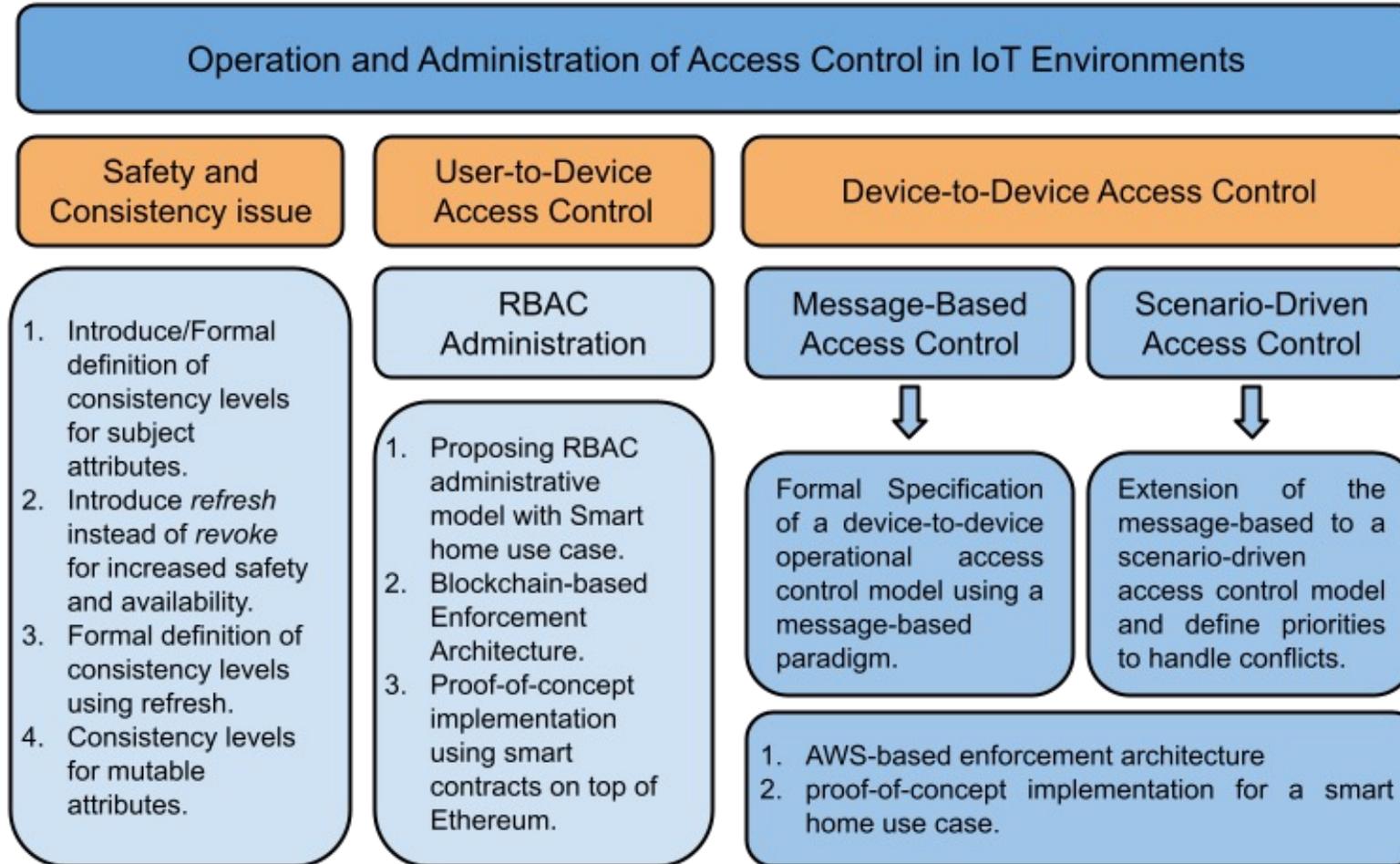


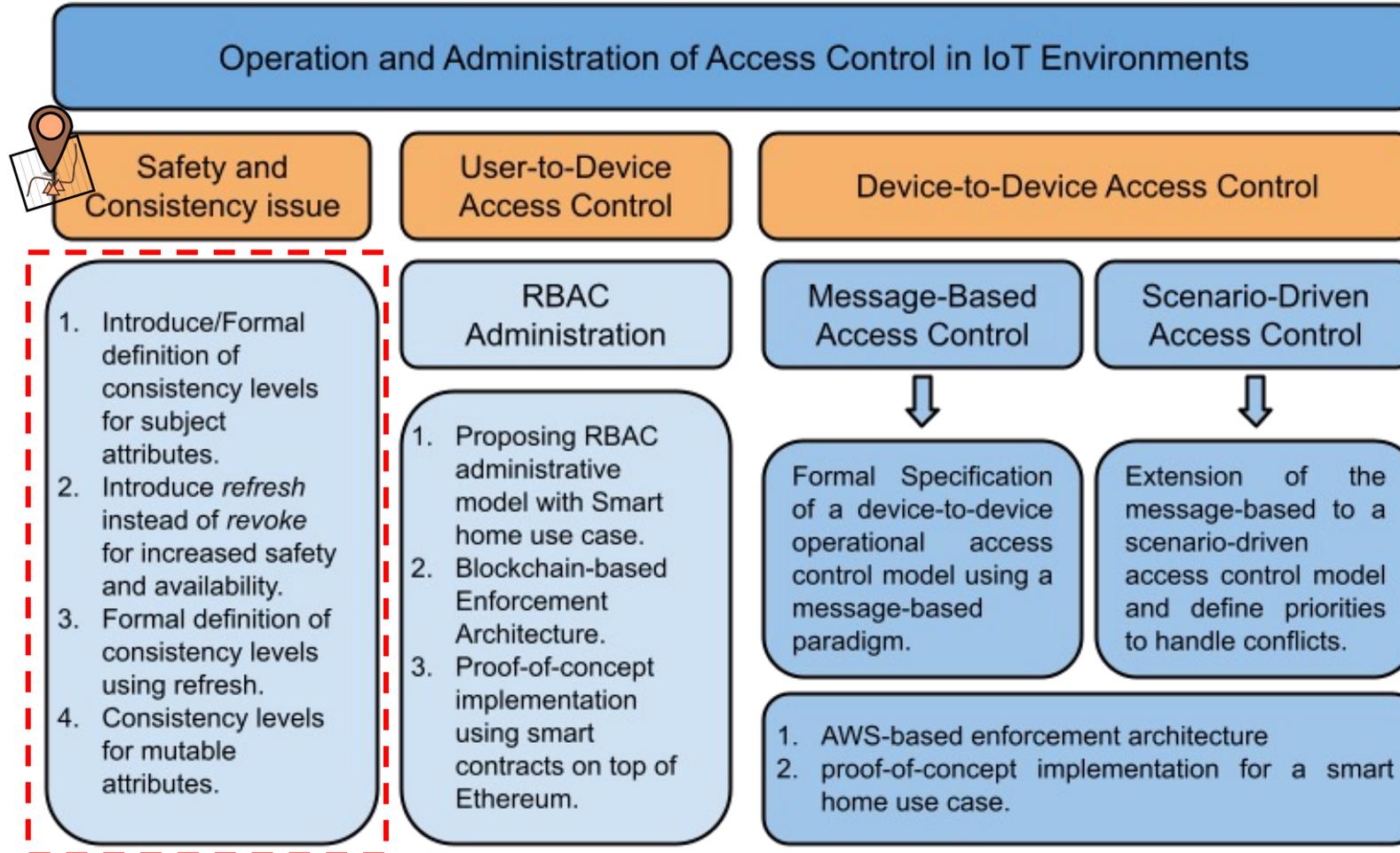
						
Policy Specification	Granularity	●	●	●	●	●
	Context Awareness	●	●	●	●	●
Policy Management	Handling the complexity of Environment	○	●	●	◐	●
	Usability	●	○	○	●	○
	Multi-domain Administration	○	●	●	◐	●
Policy Enforcement	Minimum user involvement	◐	●	●	●	●
	Light-weight	◐	◐	◐	●	●
	Reliability and Availability	●	●	●	●	●

- We recognize **smart home IoT unique characteristics** necessitate **oriented authorization models** to be particularly designed.
- Little attention has been paid to **administration of access** in IoT environments.
- The potential security violations which may happen in **device-to-device** interactions are largely uninvestigated.

- The uptick of **smart home** technology adoption while its authorization is less investigated compared to other domains.
- We investigate **three major access control-related topics** which **affect** or **directly provide authorization** in the home IoT environment.
- **ABAC** is a widely adopted paradigm to provide access control in different IoT access, including smart home. Using ABAC expose smart environments to the risk of **access violation due to inconsistency**.
- While overall system security is crucially dependent on both administrative and operational authorization models, **administration of access** in smart home environments is overlooked.
- There is a hype around utilizing **blockchain for access control** in IoT environments, but it is not clear if the provided benefits could overcome inherent blockchain drawbacks.
- A holistic view of home automation demands **catered access control specifications** to facilitate **device-to-device** interactions.

- *The established paradigm of role-based access control can be utilized for access **administration** of user-to-device in corresponding operational access control models, which could be based on either role-based or attribute-based access control.*
- *If the operational authorization is based on attribute-based access control, a detailed analysis of required **consistency** for both mutable and immutable attribute values can ultimately benefit the overall safety of the system by providing a decision point with most recent values of attribute credentials.*
- *The established paradigm of attribute-based access control can be adapted and extended to provide a **device-to-device** access control approach towards considering heterogeneous IoT devices in a smart home as an ecosystem with intercommunication.*







Consider a smart lock with its access control deployed on the vendor's cloud.

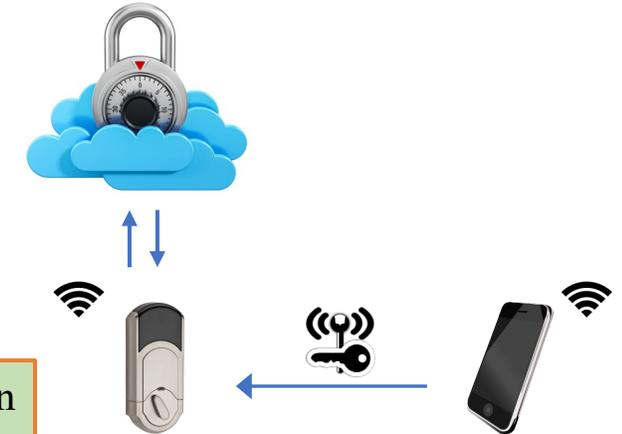


Smart lock would authenticate and pair with a user's mobile in its Bluetooth range, then receive the status of that specific user's access key from the cloud.

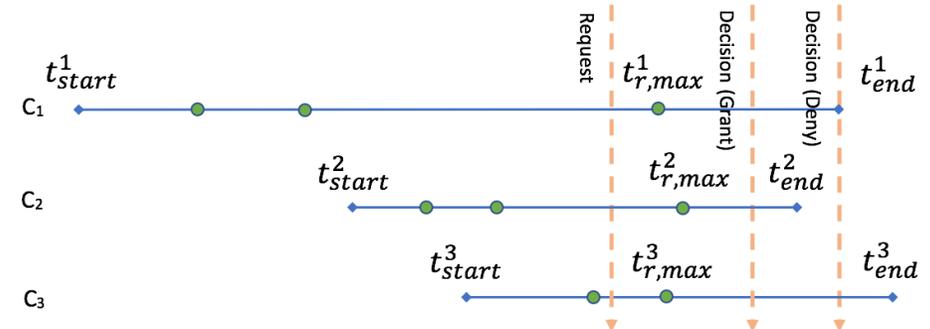


The key status would be also saved in the local database of the lock.

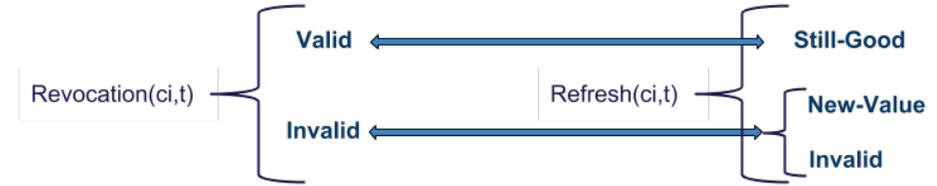
- Due to intermittent Internet connection and limited storage space of home IoT devices, required authorization information might not be available in real time.
- There is an increased risk of making access control decisions based on outdated information.
- This problem may arise in any attribute-based access control environment in which attributes are provided incrementally to the decision point.
- We investigate this problem in general, interpreted with use cases in a smart home IoT environment.



Exposure of decision point to stale attributes, hence access violation.

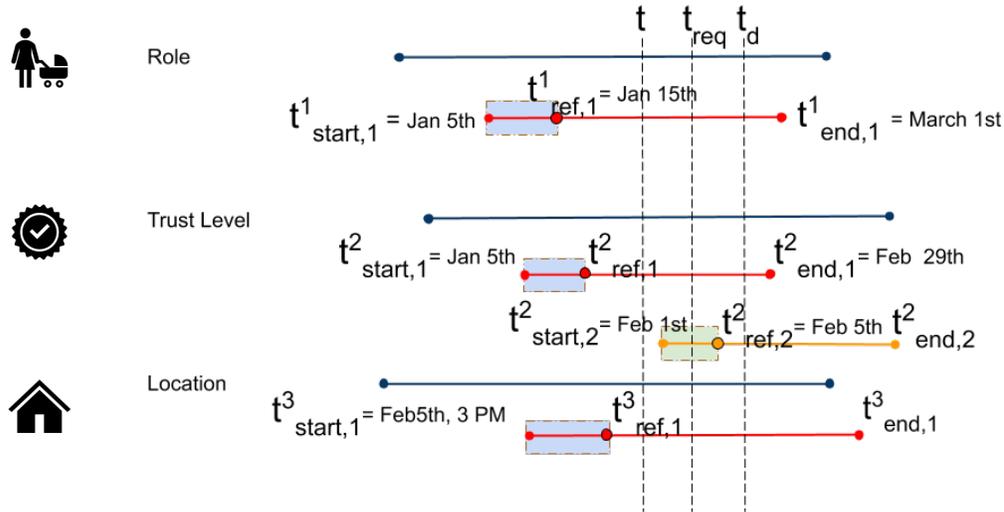


**Refresh** replaces an older value with a newer one, while **Revoke** simply invalidates the old value.



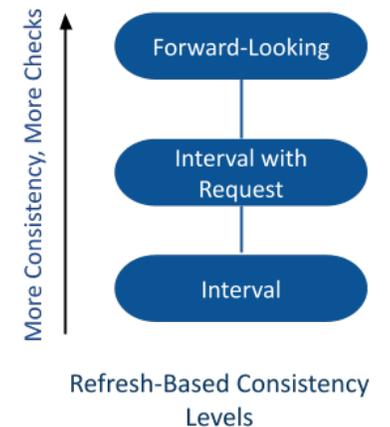
(s: user, o: thermostat, op: operation)

$$\leftrightarrow \left( (op.trustLevel \leq u.trustLevel) \wedge \left( (u.role == "parent") \vee ((u.role \in \{"parent", "kid"\}) \wedge (u.location == "home")) \right) \right)$$



Interval Consistency with Request Time

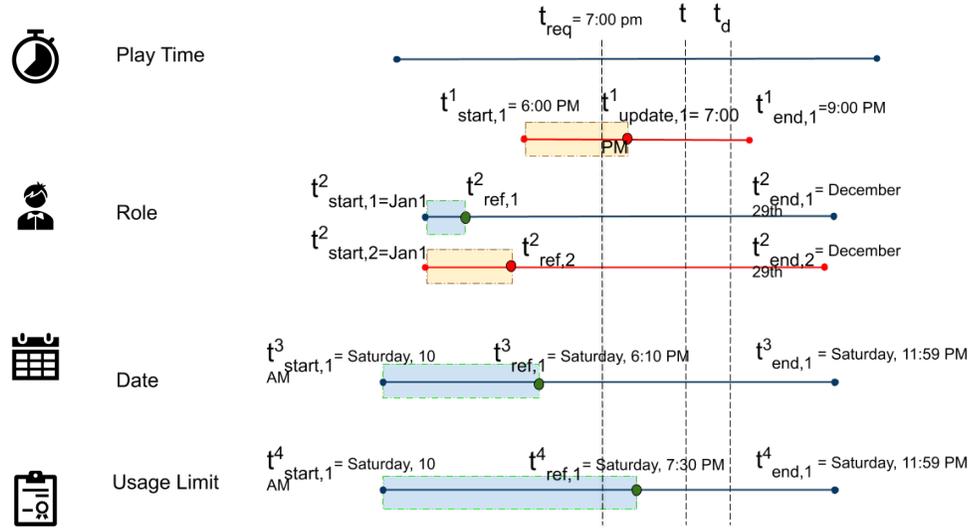
What if babysitter tries to adjust the thermostat after leaving home?



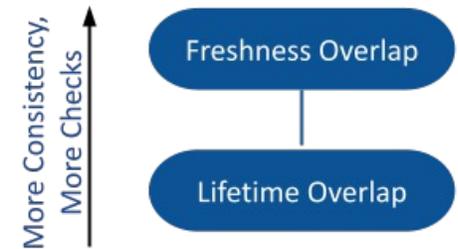
**Mutable Attribute:** Attribute changes are the consequence of access utilization by subject. We claim the revocation scenario to be inappropriate for mutable attributes.

(*s: user, o: playstation, op: operation*)

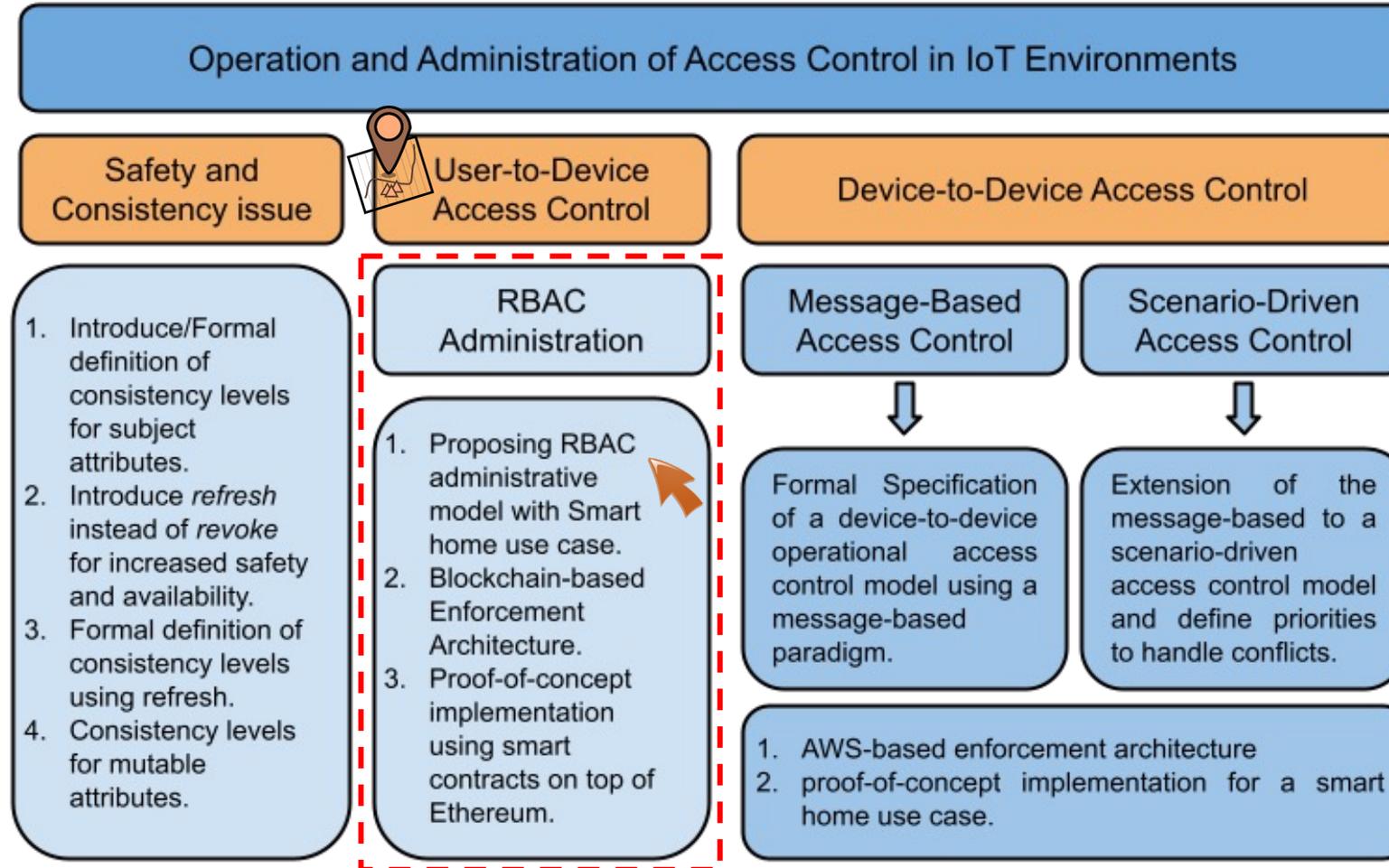
$$\leftrightarrow \left( (u.quota > 0) \wedge \left( (u.role == "parent") \vee ((u.rolr == "kid") \wedge (day\_of\_week \in \{ "Saturday", "Sunday" \})) \right) \right)$$



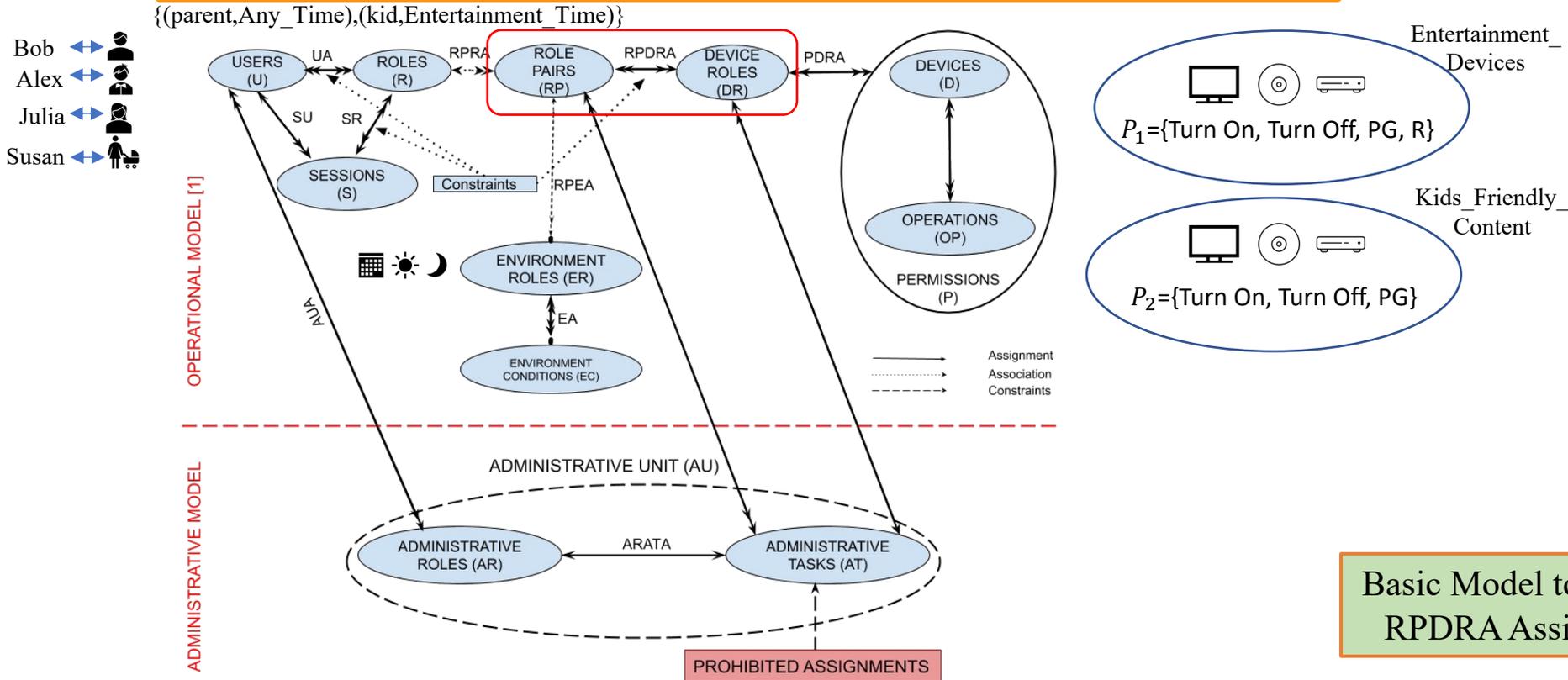
What if parents decide to reduce the usage limit?



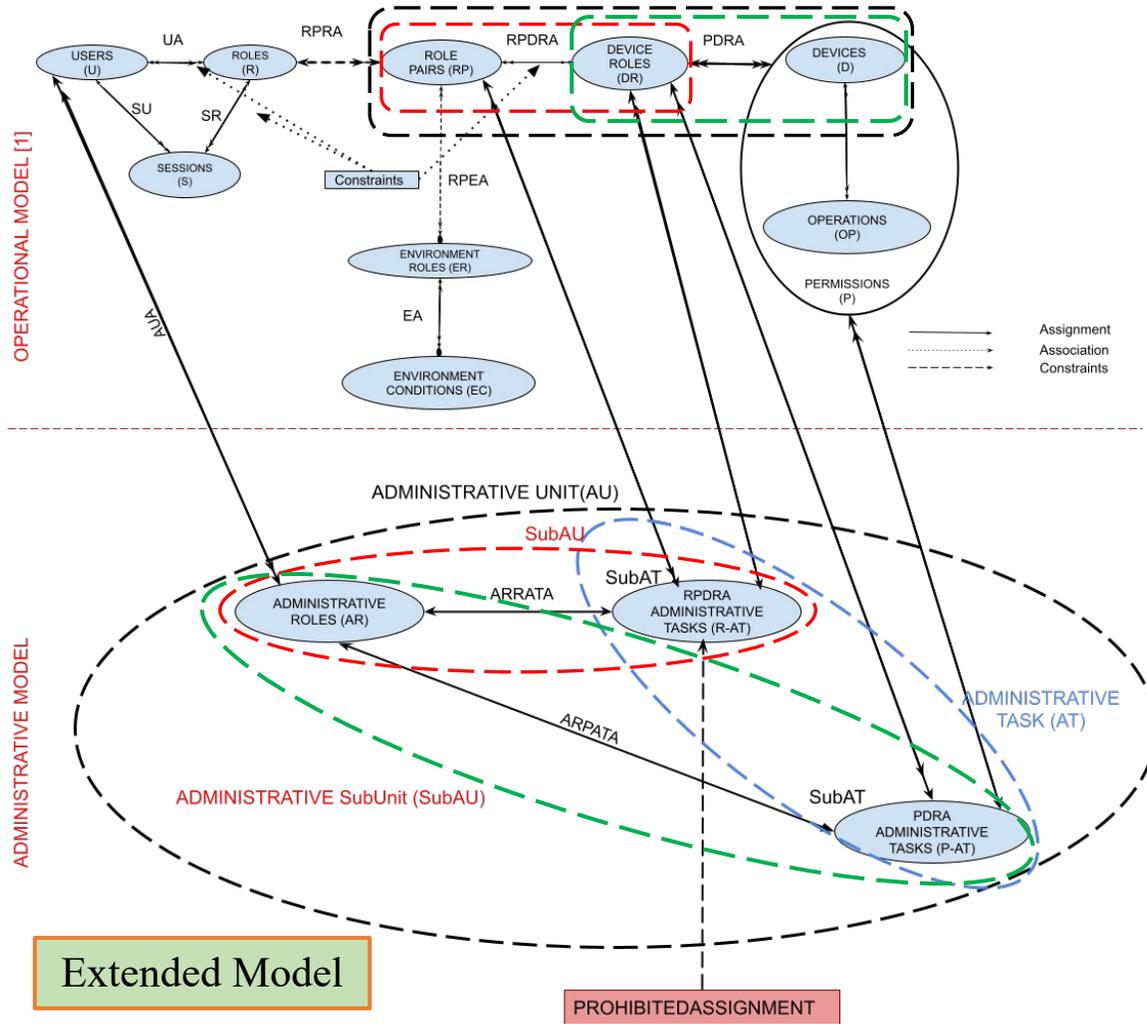
Lifetime Overlap Consistency Level



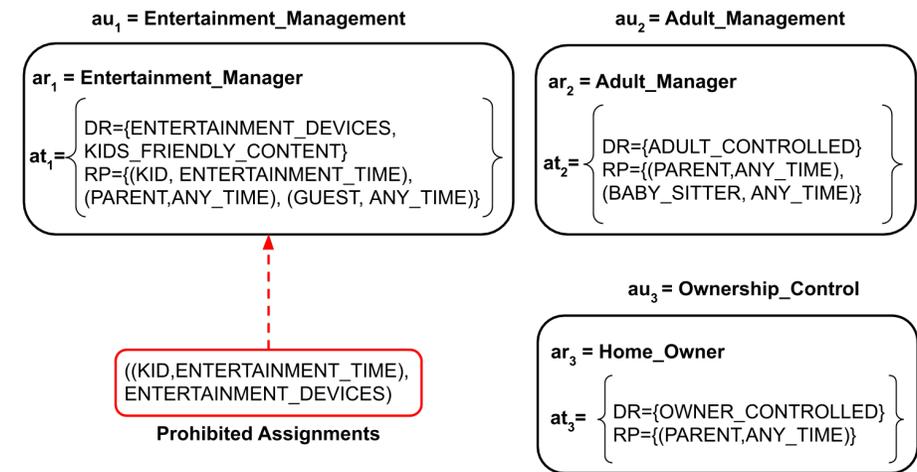
- **Access administration** has been overlooked.
  - Overall system security is crucially dependent on **both the administrative and operational models**.
  - Many **lack the support of a formal model** and rely on informally assumed policy objectives.
- Administration of access must be carefully crafted to ensure that **policy does not drift away from its original objectives**.
- Access administration in a smart home environment is a **particular** access management problem due to:
  - lack of expertise in home users.
  - shared ownership of IoT devices.
- We follow **PEI (Policy, Enforcement Architecture, Implementation)** as our reference model.
- We opt for **EGRBAC** as our underlying operational model, for it being:
  - Granular at permission level, instead of device level.
  - Capture the environment conditions and provide contextuality.

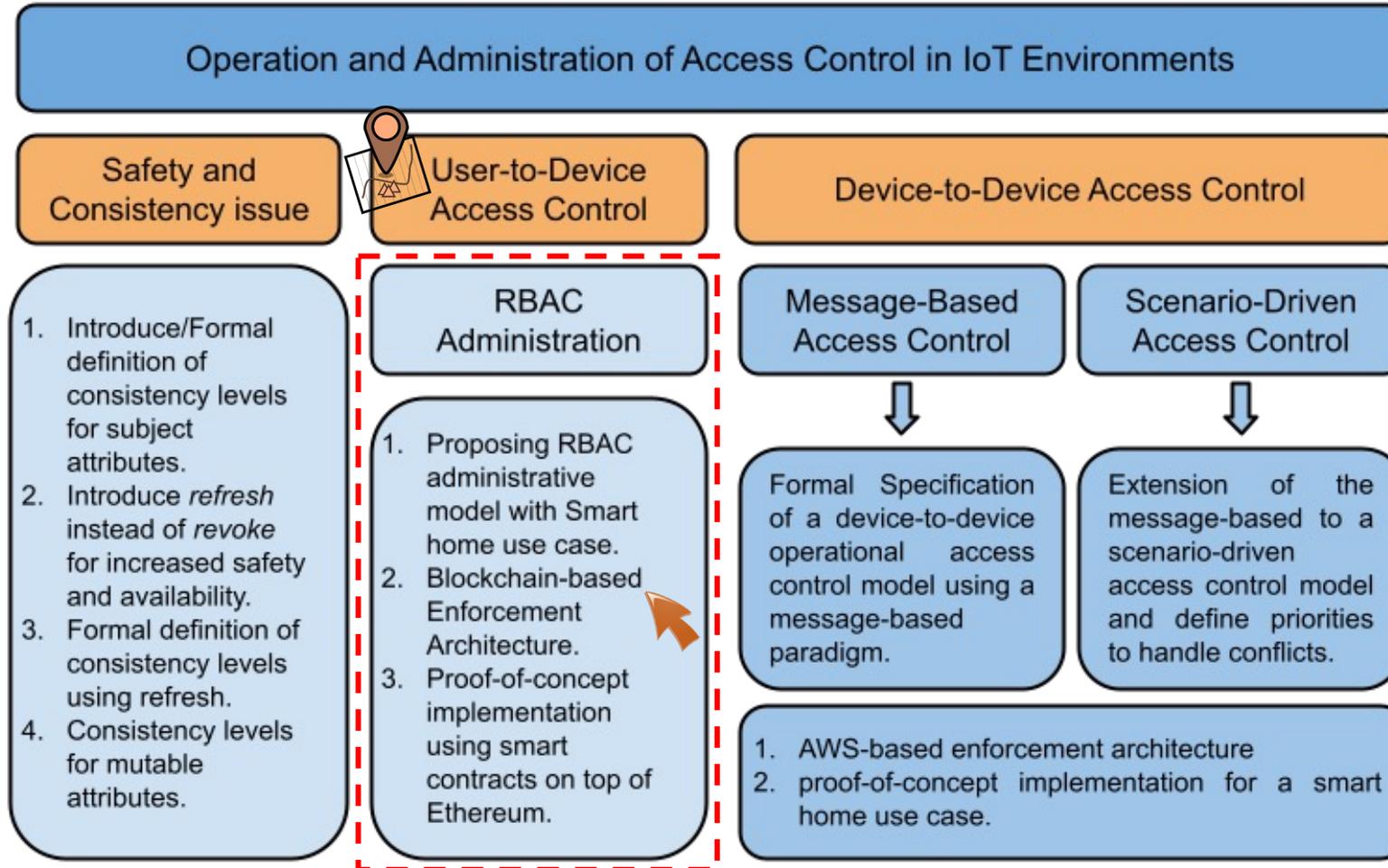


- We recognize administration is best to be done **decentralized**. Decentralization provided through **Administrative Units (AU)**.
- We define one administrative unit per operational assignment to be managed, which includes a unique **administrative role (AR)** and a set of **administrative tasks (AT)**.
- Authorization is scoped as a set of administrative tasks defined to manage corresponding assignments in the operational model.



- We extended our administrative model by defining **one administrative unit per operational assignment** to be managed.
- Each administrative unit includes a **unique administrative role** which **controls a predefined set of administrative tasks** which represents its **scope of administration**.





- Blockchain for Access Control:

### Benefits:

- Decentralized Control
- Transparency and Auditability
- Distributed Information
- Tamper-proof

### Why **NOT** Blockchain for Operational Access Control:

- IoT Constraints
- Long Transaction Confirmation Time
- Financially Prohibitive

### Why Blockchain for Administrative Access Control:

- Less Frequency of Administrative Tasks
- Posteriori Analysis
- Scalable
- No need for IoT devices to be engaged in blockchain

- Threat Model and Blockchain Benefits:

### Threat Model:

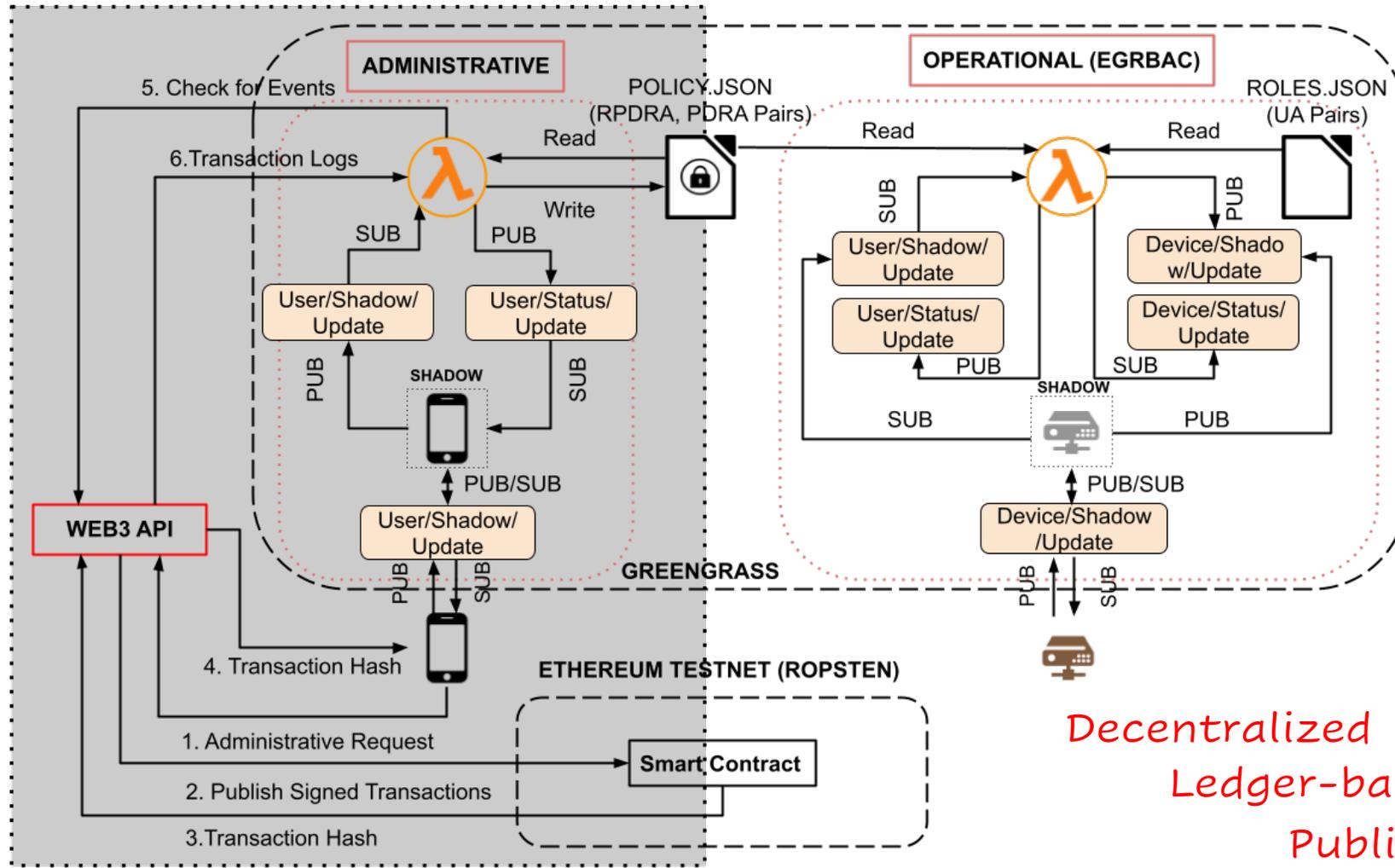
- **Insider Attack:**
  - Spoofing, Tampering, Privilege Escalation, Repudiation

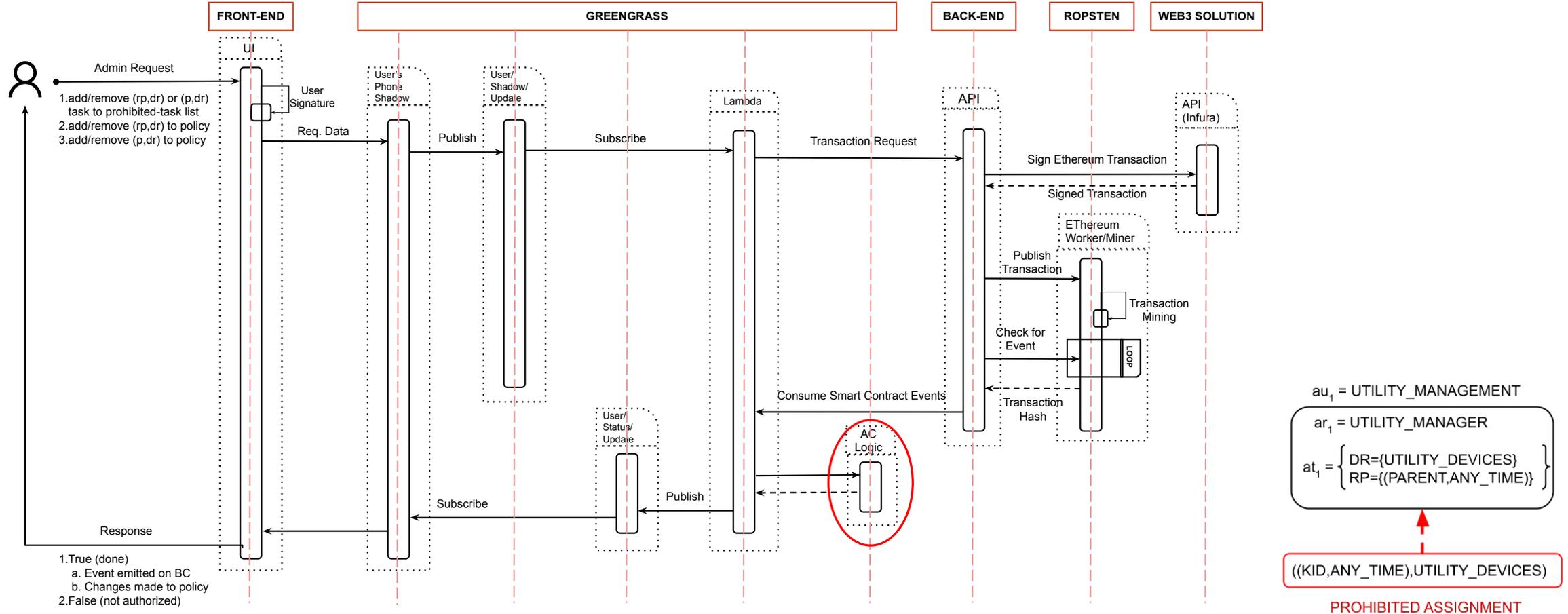
### Assumptions:

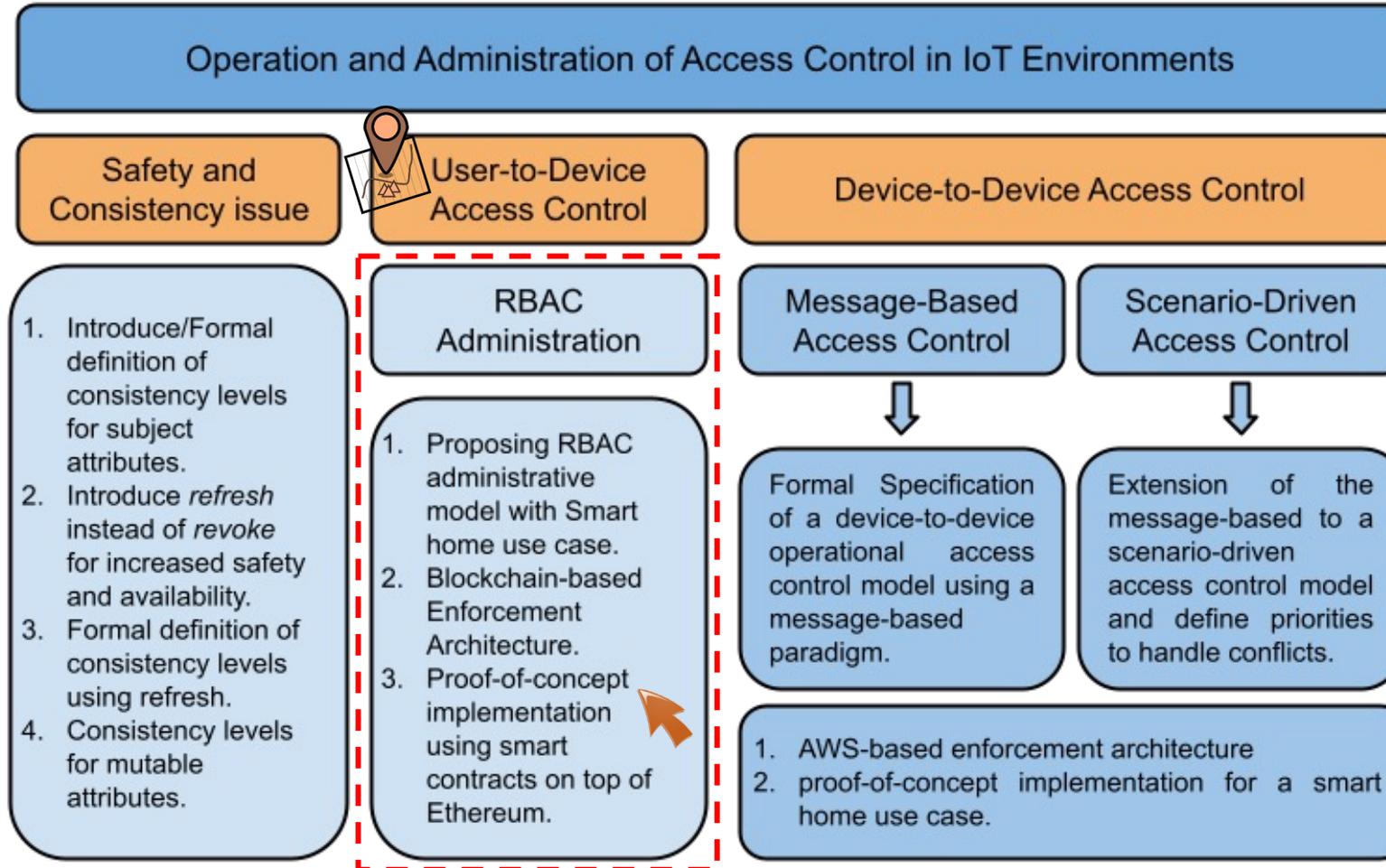
- Users' communication with edge is secure over local network.
- Routing attacks are out-of-scope
- Attacks against Web3 API are **out-of-scope**
- API attacks against user's private keys in their wallets considered to be **out-of-scope**

### Blockchain Security Benefits:

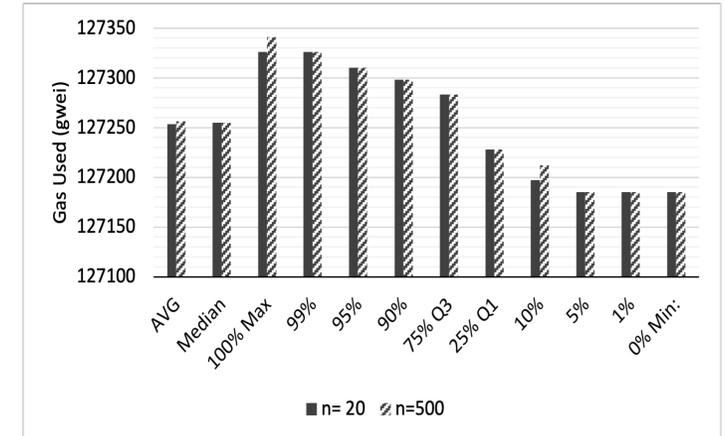
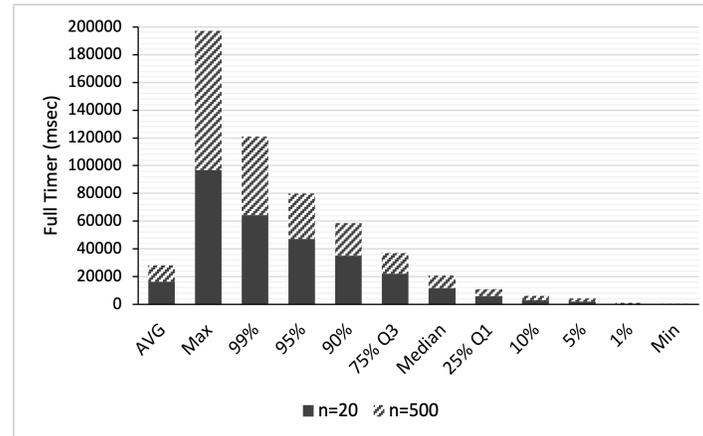
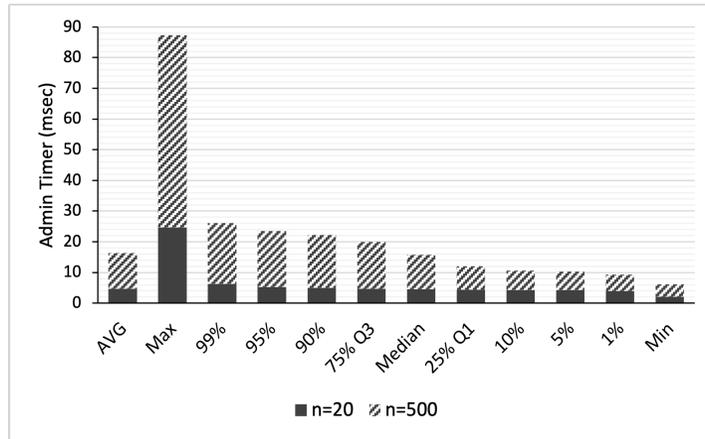
- Administrator account cannot be faked
- Administrative policy is encoded in a smart contract recorded to ledger via consensus
- System is equipped with transparency and auditability







- Administrative access control policy implemented in a [single smart contract](#) on the [Ropsten](#).
- Different [administrative controls](#) are coded as [functions](#), which would be triggered by transactions.
- Smart contract is programmed in [Solidity](#) and tested it on [Remix IDE](#).
- [Infura](#) is used as [web3.0 API](#) to interact with blockchain.
- Experiment Environment:
  - AWS IoT Greengrass v1.
  - Greengrass runs on a dedicated virtual machine: one virtual CPU, 2 GB of RAM and 20 GB hard drive.
  - The virtual machine's operating system is Ubuntu 20.4.2 LTS and it is connected to a 1 Gbps network.
  - Our AWS lambda code on the Greengrass is written in Python 3.8 and is running in a long-lived isolated runtime environment with limited RAM of 256 MB
- Experiments are done for a [normal distribution with a 99.9% confidence interval](#).
- We ran our experiments in two settings with the policy sizes of  $n=20$  and  $n=500$ .
  - Both experiments were run for a total of [500 times](#).



**Admin Timer:** After a transaction has been successfully mined, Lambda checks the logs to search out the succeeded transactions. Then, it makes appropriate changes to the “policy.json” file and publishes the results to the User/Status/Update to inform the user about his/her administrative request.

**Full Timer:** Complete cycle of an administrator submitting a request, to that request being mined, and the lambda function processing the results and updating as necessary.

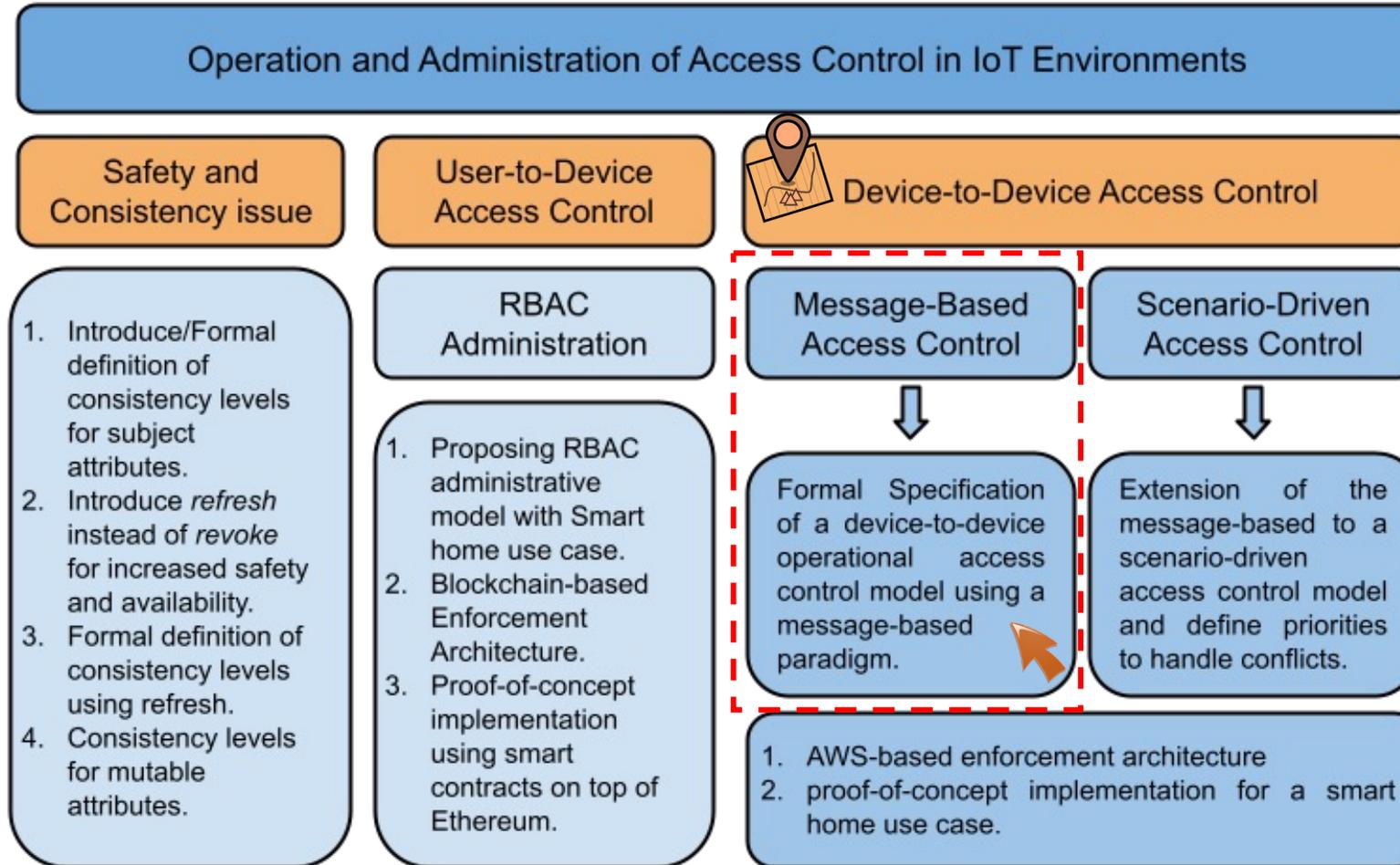
**Gas Used:** the actual amount of gas which was used during execution. Gas prices are denoted in GWEI, which equals to  $10^{-9}$  ETH. We calculated the monetary cost of each transaction to be 28 cents.

- Our administrative model **features**:
  - Decoupled Assignment and Revocation
  - Symmetric Assignment and Revocation
  - Generalizability
  - Transparency and Auditability
  - Privacy
- **Security considerations** specific to our architecture:
  - Smart Contract Security
  - Device-Cloud Communications
- **Limitations**:
  - Continuous Access control and Mutability
  - Handling Conflicts
- Our implementation results are reassuring that although the use of **blockchain for operational access control** is **NOT promising**, **BUT** it is **promising** to be utilized at administrative level.

### Blockchain Hype

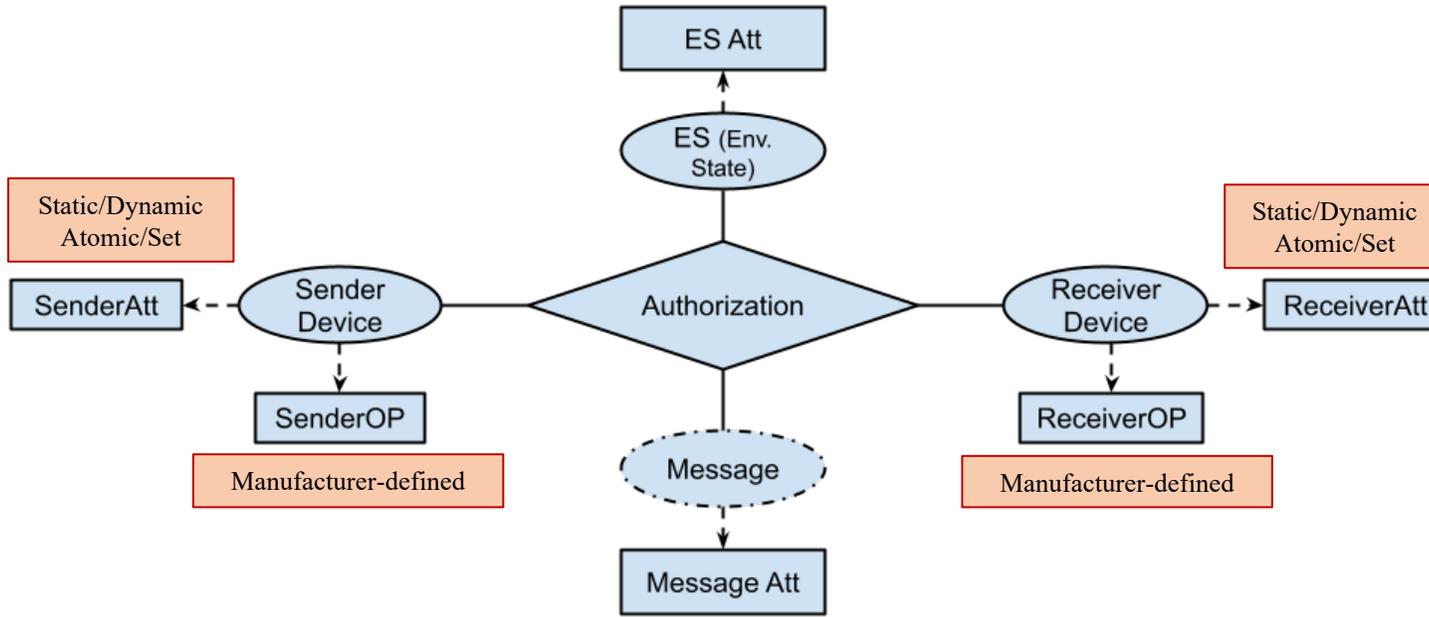


EGRBAC is not chosen as a de-facto!



- Seamless **interoperability** among IoT devices, **device-to-device communication**, is imperative for incipient evolution of the IoT ecosystem.
- Many standardization efforts as well as proposals for integration of heterogeneous IoT platforms are going on in both academic and industry, but there is **no de-facto standard** and **there would be none in foreseeable future**.
- There is **heterogeneity in all levels of IoT** technology, including **device, networking, middleware, application, and data/semantics** of IoT scenarios, makes heterogeneous IoT devices cumbersome.
- There is no access control model specification to provide an **access control model** for **device-to-device** interoperability.
- Scenarios of **device-to-device interactions** are inevitable for real-life **home automation**, which brings **added convenience** but also **inherent security risks**.

We present an **access control model** which governs **authorized flow of information** for **device-to-device** interactions in a **smart home IoT** using **Attribute-Based Access Control (ABAC)** for the first time, utilizing a **message passing paradigm**.



$$M = \{m_i: m_i = ((att_1 = value_1), (att_2 = value_2), \dots, (att_n = value_n))\}$$

$$CheckAtt = True \leftrightarrow typeSetAtt(m) = \begin{cases} \subseteq 2^{DAA(r)} & \text{if } m.value_1 = \text{"query"} \\ \in DOA(r) & \text{if } m.value_1 = \text{"command"} \\ \subseteq 2^{DAA(s)} & \text{if } m.value_1 = \text{"info"} \end{cases}$$

### Core Components

- $D$  is a set of smart home IoT devices deployed by homeowner.
- $OP$  is a set of operations available on different devices in the system (manufacturer specified).
- $ES = \{current\}$  is the singleton set, representing the environment state at the current time instant.
- $Ent = D \cup ES \cup M$  is the set of entities in the system, where the set of messages  $M$  is defined below.
- $DOA : D \rightarrow 2^{OP}$  is a one to many relation which associates a device to its available operations as specified by the device manufacturer.

### Attribute Functions

- $DAA, EAA$  are respectively sets of attribute functions which associate a device or the current environment state with attribute values.
- $attValue\ Type : DAA \cup EAA \rightarrow \{atomic, set\}$
- $\forall att \in DAA \cup EAA, Range(att)$ , is the attribute range, a finite set of atomic values.
- Each  $att \in DAA \cup EAA$  maps a device/environment to a single *atomic* value or to a finite *set* of values, as follows:

$$-att : DAA \cup EAA \rightarrow \begin{cases} Range(att) & \text{if } attValue\ Type(att) = atomic \\ 2^{Range(att)} & \text{if } attValue\ Type(att) = set \end{cases}$$

- $attAssign\ Type : DAA \cup EAA \rightarrow \begin{cases} static & \text{set/changed via administrative actions} \\ dynamic & \text{set/changed automatically by deployed sensors in home IoT} \end{cases}$

### Message and Message Functions

- $M = \{m\}$  is the set of all messages in the system.
- $m = \{(att_1, value_1), (att_2, value_2), \dots, (att_n, value_n)\}$ , represents any single message in the system with  $n$  different attributes, each of which is indicated as a (*key, value*) pair.
- $typeSet = \{\text{"query"}, \text{"command"}, \text{"info"}\}$  is a mandatory first attribute in every message which indicates its *type* and thereby the rest of message attributes.
- For each  $m \in M$ , we assume the first attribute determines the type of the message:  $att_1 \in typeSet$
- $typeSetAtt : M \rightarrow 2^{DAA} \cup 2^{DOA}$ , is a function which indicates the set of attribute keys required to be communicated based on the message type, supposed to be communicated via  $\{att_2, \dots, att_n\}$  in each message.

### Check Access Predicate

- $CheckAccess$  is evaluated when a sender device ( $s$ ) tries to send a message ( $m$ ) to a receiver device ( $r$ ) in context of current environment state ( $current$ ) and is evaluated based on following formula:

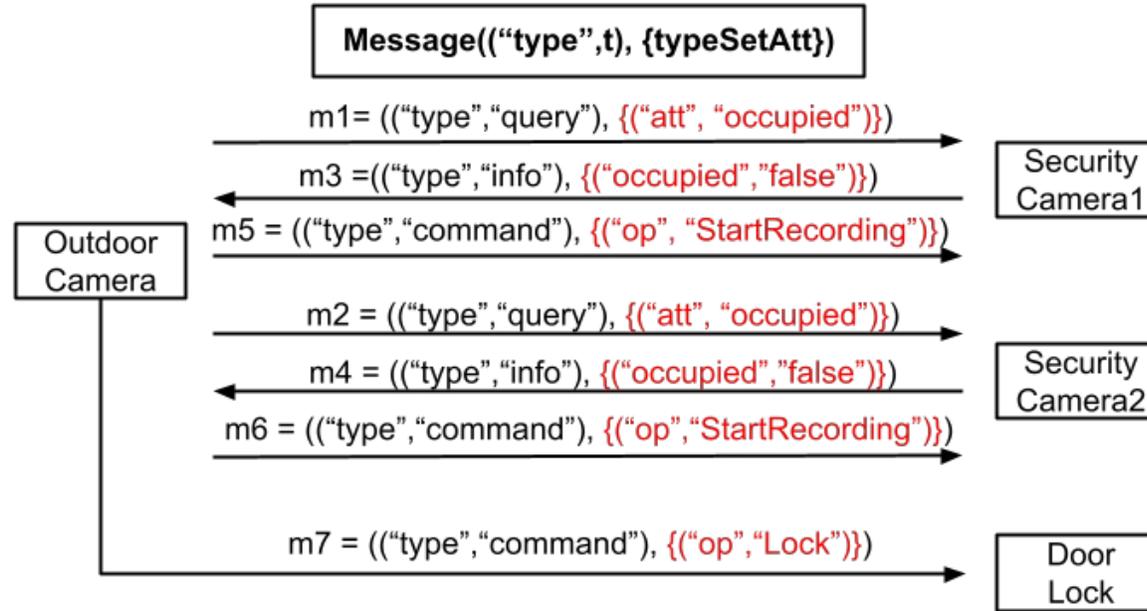
$$-CheckAccess(s : D, m : M, r : D, current : ES) \equiv CheckAtt(s : D, m : M, r : D, current : ES) \wedge Authorization(s : D, m : M, r : D, current : ES)$$

$$-CheckAtt = True \iff typeSetAtt(m) = \begin{cases} \subseteq 2^{DAA(r)} & \text{if } m.value_1 = \text{"query"} \\ \in DOA(r) & \text{if } m.value_1 = \text{"command"} \\ \subseteq 2^{DAA(s)} & \text{if } m.value_1 = \text{"info"} \end{cases}$$

- $Authorization(s : D, m : M, r : D, current : ES)$  is a logical proposition which could be evaluated to either True or False and is created using following policy rules.

$$-p \equiv (p) \mid \neg p \mid p \wedge p \mid p \vee p \mid \exists x \in set.p \mid \forall x \in set.p \mid set \Delta \ set \mid atomic \in set$$

$$-\Delta \equiv \subseteq \mid \subset \mid \supseteq \mid \supset \mid \cup$$

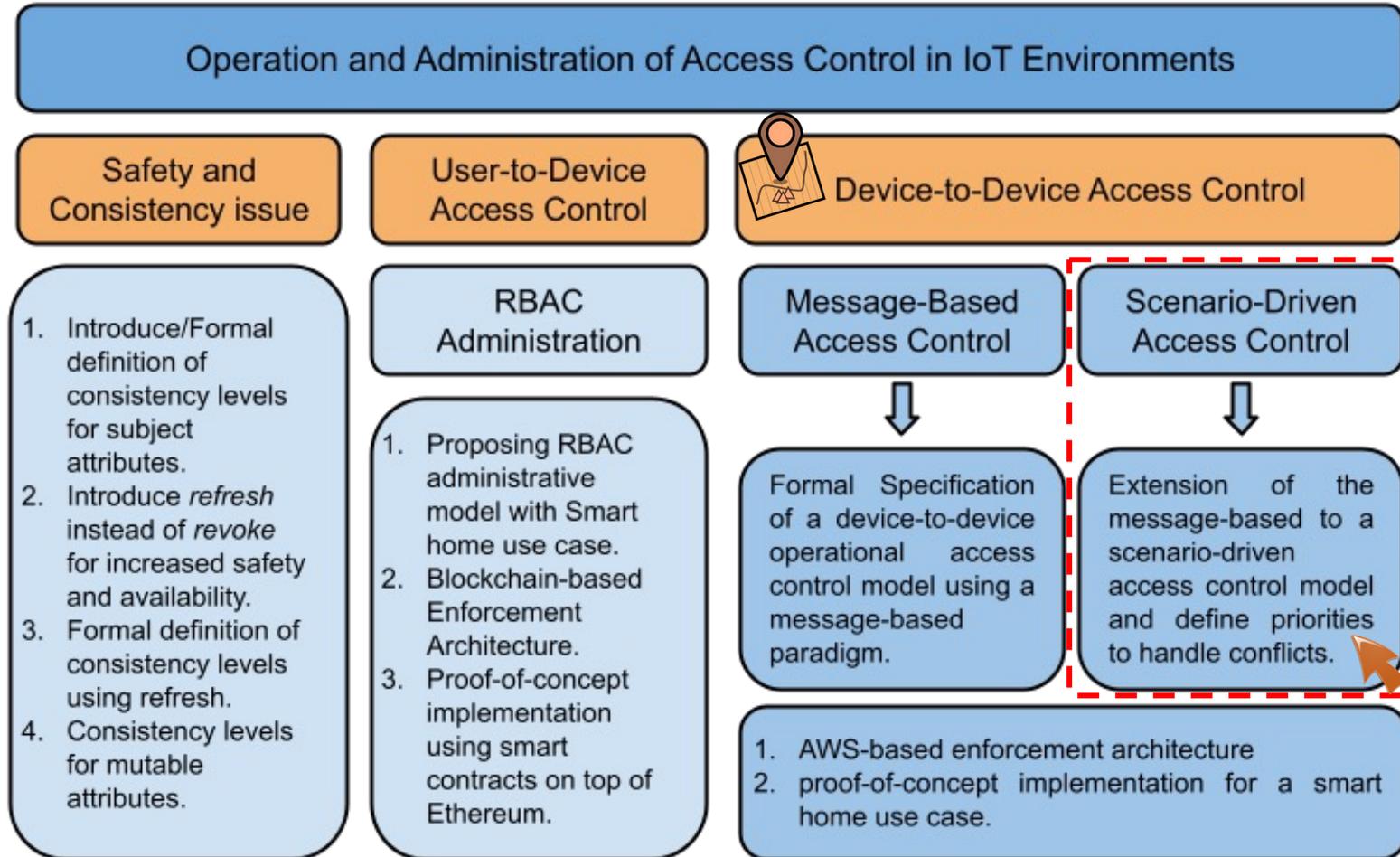


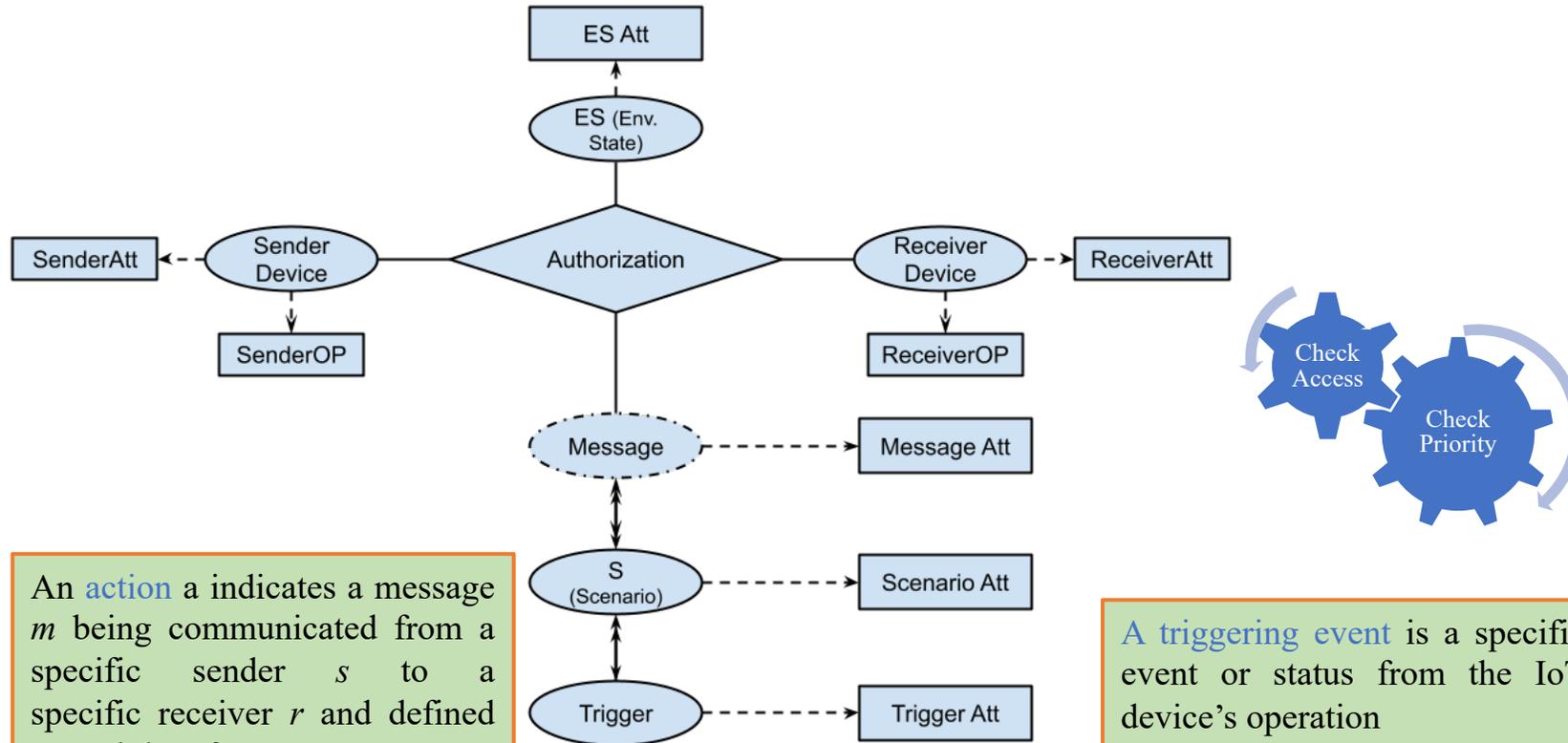
$CheckAccess(s:D, m:M, r:D, current:ES) \equiv CheckAtt(s: D, m: M, r: D, current: ES) \wedge Authorization(s: D, m: M, r: D, current: ES)$

$Authorization(s: D, m: M, r: D, current: ES) = ((m.att_1 = "query") \wedge (typeSetAtt(m) \in \{ "recording", "occupied" \}) \wedge (type(s) = type(r) = "cameras") \wedge (location(s) = "outdoor") \wedge (location(r) = "indoor"))) \vee ((m.att_1 = "info") \wedge (typeSetAtt(m) \in \{ "recording", "occupied" \}) \wedge (type(r) = type(s) = "cameras") \wedge (location(s) = "indoor") \wedge (location(r) = "outdoor"))) \vee ((m.att_1 = "command") \wedge (typeSetAtt(m) \in \{ "StartRecording", "StopRecording" \}) \wedge (type(r) = type(s) = "cameras") \wedge (location(s) = "outdoor") \wedge (location(r) = "indoor"))) \vee ((m.att_1 = "command") \wedge (typeSetAtt(m) \in \{ "Lock", "Unlock" \}) \wedge (type(s) = "cameras") \wedge (type(r) = "locks") \wedge (location(s) = "outdoor") \wedge (location(r) = "mainEntrance")))$

- IoT devices' **susceptibility to cyberattacks** could make them disturbing security holes.
  - A hacker with access to your thermostat could fiddle with it, causing your HVAC system to malfunction.
  - An attacker may lock you out of your home in case the door lock is hacked.
- We adopt the **Dolev-Yao (DY) threat model** in which communicating endpoints **cannot** be assumed as trusted nodes in the network.
  - An adversary can **tamper with the data** through **modification**, **deletion**, or **insertion of fake information** as the communications rely on **wireless medium**.
- Our **Assumptions**:
  - As many IoT devices are not IP-enabled, using a gateway (GW) node in the network is inevitable. We **assume the GW node in our model is trustworthy and available**, which is a common assumption.
  - Attacker is assumed to be an **outsider** to the network with the goal of obtaining illegitimate access to available functionalities/operations of smart home IoT devices.
  - We do **not** consider adversaries to have **physical access** to IoT devices.

Our access control approach provides a **defense-in-depth prevention/protection** against any outsider attack, through restricting the set of authorized messages being communicated among IoT devices.





An **action**  $a$  indicates a message  $m$  being communicated from a specific sender  $s$  to a specific receiver  $r$  and defined as a triplet of:

$$a = (s: D, m: M, r: D)$$

A **scenario** is defined as a set of actions to be done in the smart home

**Priority** is a totally ordered set relation, depicted as  $(pr, <)$  between any two triggering events  $tr_i$  and  $tr_j$  and is reflected in their (administratively) assigned priority values. In a smart home environment, the priority values could be defined as:

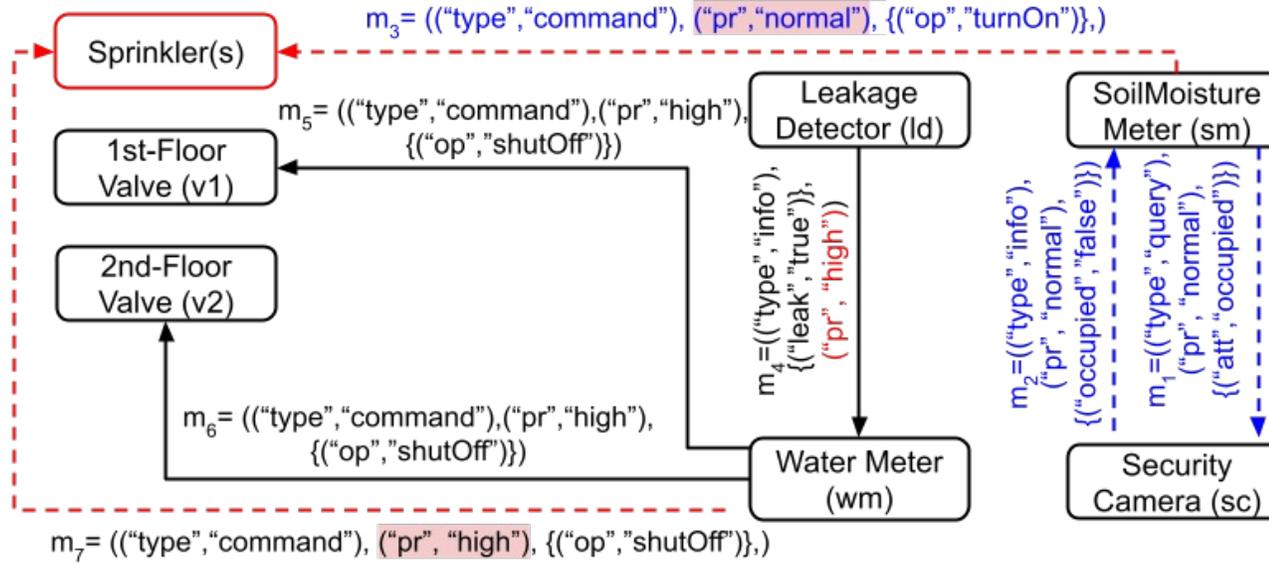
$$(pr, <) = (\perp < low < medium < high <)$$

A **triggering event** is a specific event or status from the IoT device's operation

–  $TeSA: TE \rightarrow 2^S$  is a one-to-many relation which defines a (set) of scenario(s) that would be provoked by a triggering event  $te \in TE$ .  
 – For each  $s \in S: s = (Action_s \subseteq A, tr_s: TE, pr_s: (pr, <), active: \{ "true", "false" \}, id)$ , and  $pr_s = pr_A(tr_s)$ .  
 –  $active(s: S) = \begin{cases} "true" & \text{while } tr_s \text{ is in effect.} \\ "false" & \text{As soon as } tr_s, \text{ triggering event, reverts.} \end{cases}$   
 –  $typeSet = \{ "query", "command", "info" \}$  is a mandatory attribute of every message which indicates its *type* and thereby the rest of message attributes.  
 – For each  $m \in M$ , we assume the first attribute must determine the type of the message and the second attribute must determine its priority:  $att_i \in typeSet, att_2 \in (pr, <)$   
 –  $typeSetAtt: M \rightarrow 2^{DAA} \cup 2^{DOA}$ , is a function which indicates the set of attribute keys required to be communicated based on the message type, supposed to be communicated via  $\{att_1, \dots, att_n\}$  in each message.  
 –  $msgPrA: D \times M \times D \rightarrow (pr, <)$  is a function which is checked each time a device  $s_d$  wants to create a command message  $m$  and assigns proper priority to it, in order to be sent to device  $r_d$ .  
 –  $msgPrA(s: D, m: M, r: D) = \begin{cases} pr_s & \text{if } (m.att_1 = "command") \wedge (\exists s \in S: active(s) = "true") \wedge \\ & (\exists a \in Action_s: a.s = s_d \wedge a.r = r_d \wedge typeSetAtt(a,m) = typeSetAtt(m)) \\ \perp & \text{otherwise} \end{cases}$

**Check Access Predicate**  
 – *CheckAccess* is evaluated when a sender device ( $s$ ) wants to send a message ( $m$ ) to a receiver device ( $r$ ) in the context of current environment state (*current*) and is evaluated based on following formula:  
 –  $CheckAccess(s: D, m: M, r: D, current: ES) \equiv CheckAtt(s: D, m: M, r: D, current: ES) \wedge CheckPriority(s: D, m: M, r: D, current: ES) \wedge Authorization(s: D, m: M, r: D, current: ES)$   
 –  $CheckAtt = True \iff typeSetAtt(m) = \begin{cases} \subseteq 2^{DAA(r)} & \text{if } m.value_1 = "query" \\ \subseteq DOA(r) & \text{if } m.value_1 = "command" \\ \subseteq 2^{DAA(s)} & \text{if } m.value_1 = "info" \end{cases}$   
 –  $CheckPriority(s: D, m: M, r: D, current: ES) \equiv \begin{cases} "false" & \text{if } (m.att_1 = "command") \wedge \\ & [(m."op", currentOP(r)) \in conflict(r)] \wedge (m.value_2 < currentPR(r)) \\ "true" & \text{otherwise} \end{cases}$   
 – *Authorization*( $s: D, m: M, r: D, current: ES$ ) is a logical proposition which could be evaluated to either True or False and is created using following policy rules.  
 –  $p \equiv (p) \mid \neg p \mid p \wedge p \mid p \vee p \mid \exists x \in set.p \mid \forall x \in set.p \mid set.D \mid atomic \in set$   
 –  $\Delta \equiv \subseteq \mid \supseteq \mid \sqsubset \mid \sqsupset \mid \cup$

Scenario-driven access control model resolves a class of conflicts due to receiving two conflicting command messages by the the same device .



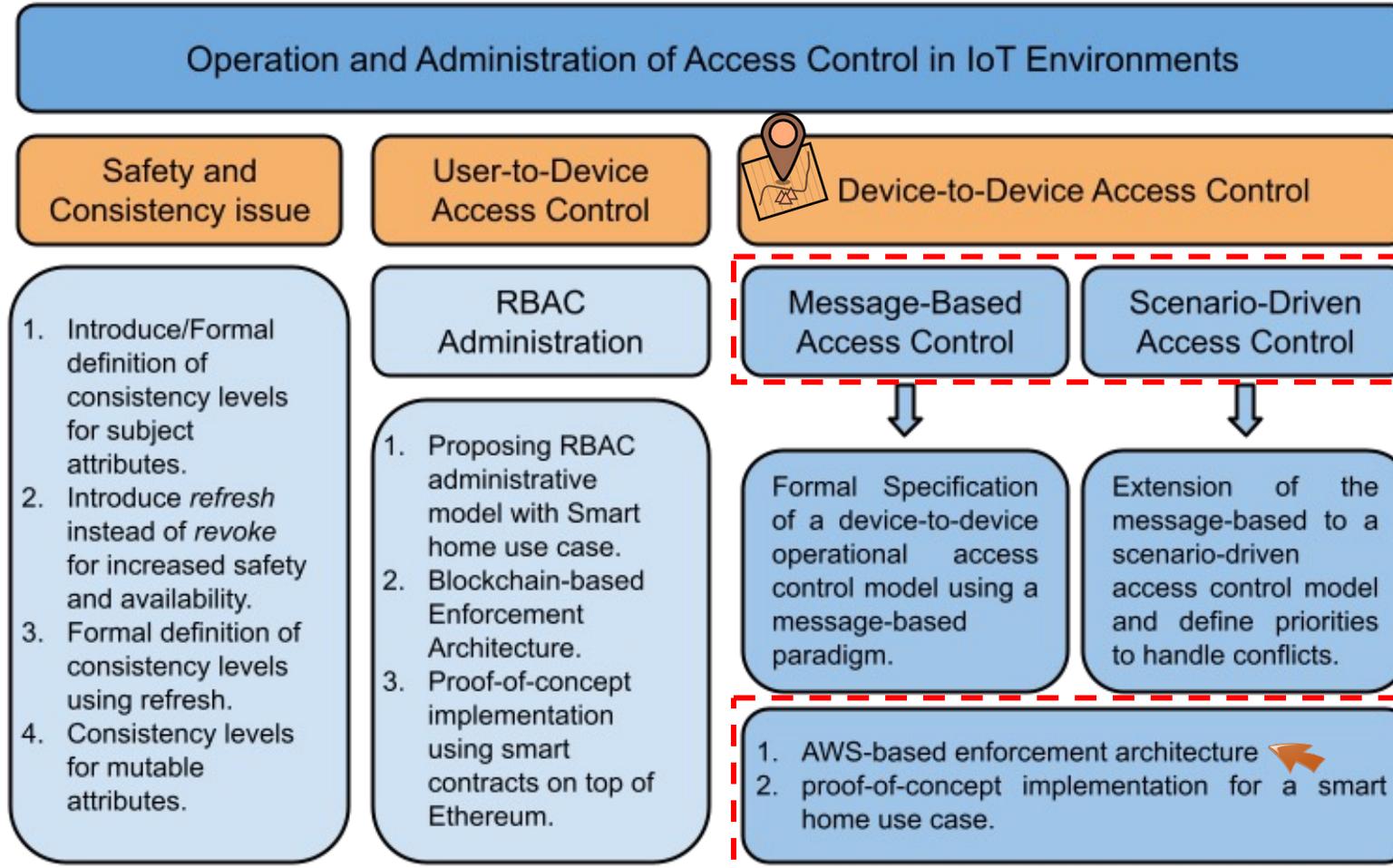
$$s_2 = \{m_4, m_5, m_6, m_7\}, pr_{s_2} = \text{"high"}$$

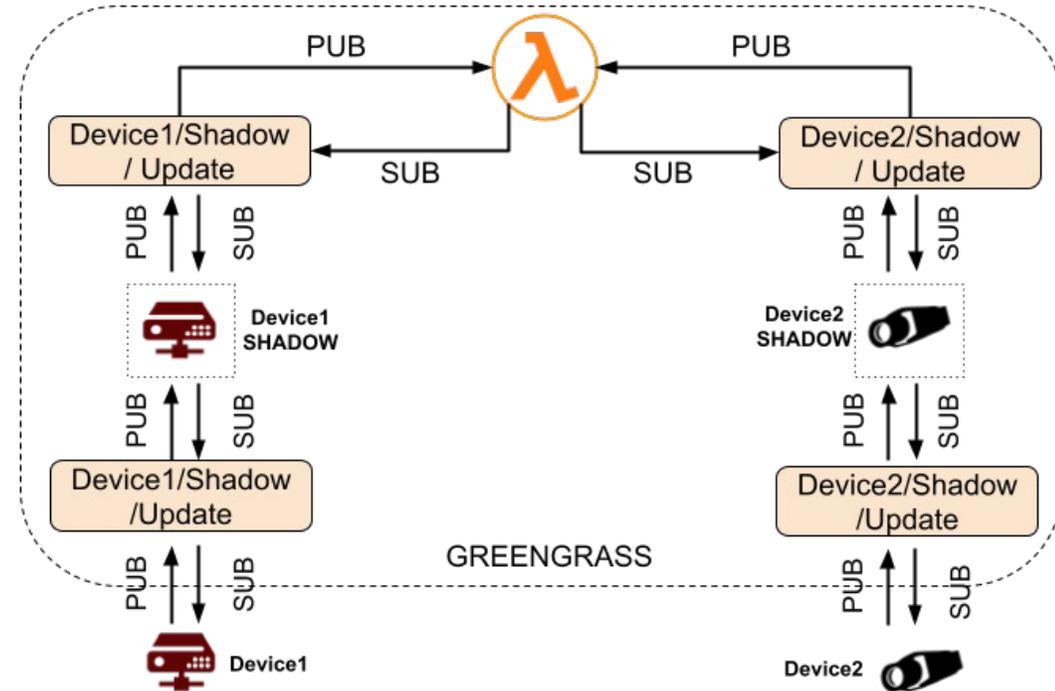
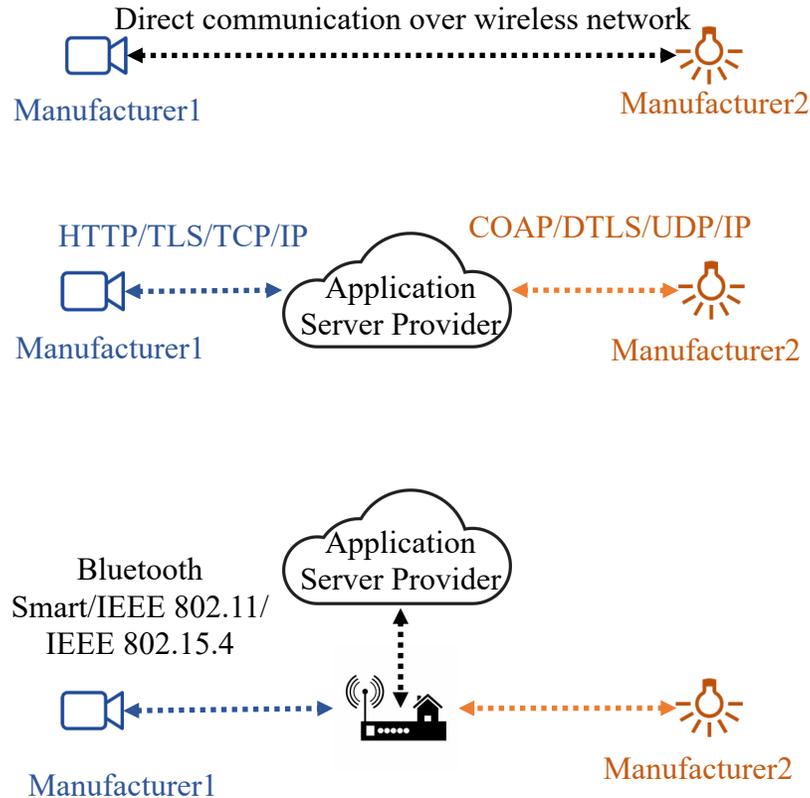
$$s_1 \leq s_2$$

$$s_1 = \{m_1, m_2, m_3\}, pr_{s_1} = \text{"normal"}$$

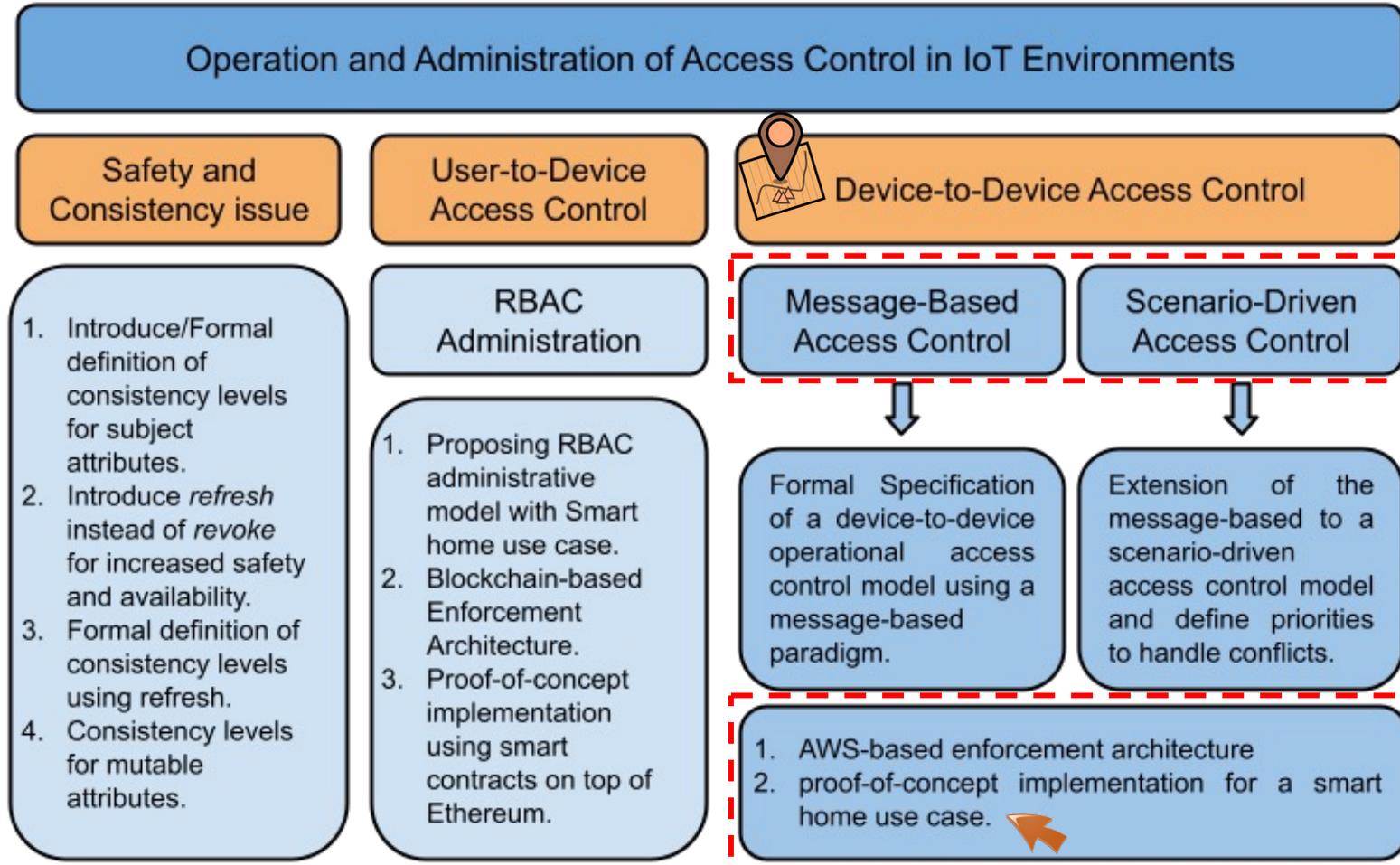
$CheckAccess(s:D, m:M, r:D, current:ES) \equiv CheckAtt(s: D, m: M, r: D, current: ES) \wedge CheckPriority(s: D, m: M, r: D, current: ES) \wedge Authorization(s: D, m: M, r: D, current: ES)$

$Authorization(s: D, m: M, r: D, current: ES) = ((m.att_1 = \text{"info"}) \wedge (typeSetAtt(m) \in \{\text{"leak"}\}) \wedge (type(r) = \text{"valves"}) \wedge (subtype(r) = \text{"watermeter"})) \vee ((m.att_1 = \text{"command"}) \wedge (typeSetAtt(m) \in \{\text{"ShutOff"}, \text{"TurnOn"}\}) \wedge (type(r) = type(s) = \text{"valves"}) \wedge (subtype(s) = \text{"watermeter"})) \vee ((m.att_1 = \text{"query"}) \wedge (typeSetAtt(m) \in \{\text{"occupied"}\}) \wedge (type(s) = \text{"detectors"}) \wedge (subtype(s) = \text{"soil"}) \wedge (type(r) = \text{"cameras"}) \wedge (location(r) = \text{"outdoor"})) \vee ((m.att_1 = \text{"info"}) \wedge (typeSetAtt(m) \in \{\text{"occupied"}\}) \wedge (type(s) = \text{"camera"}) \wedge (location(s) = \text{"outdoor"}) \wedge (type(r) = \text{"detectors"}) \wedge (subtype(r) = \text{"soil"}))$

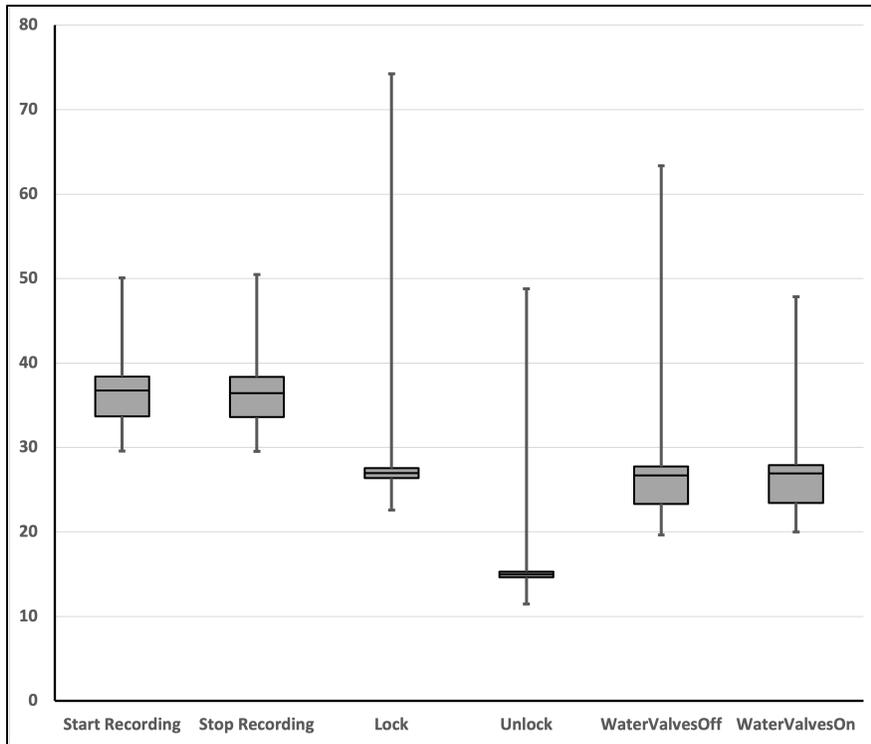




- Greengrass utilization enables devices to autonomously react to local events and securely communicate with each other over the local network.
- Using MQTT, devices will communicate on their private topics, `device/shadow/update`, to update their shadows, and trigger the local computation function, known as Lambda function.
- When Lambda is triggered, it executes the code we developed to check the policy.



Experiment	Min	10%	25%	Median	75%	90%	95%	99%	Max
Full Timer	24.2	25.4	26.0	34.9	36.7	39.1	42.9	47.2	67.2
State Update	3.5	3.6	3.7	5.7	6.6	6.9	7.4	10.2	21.8



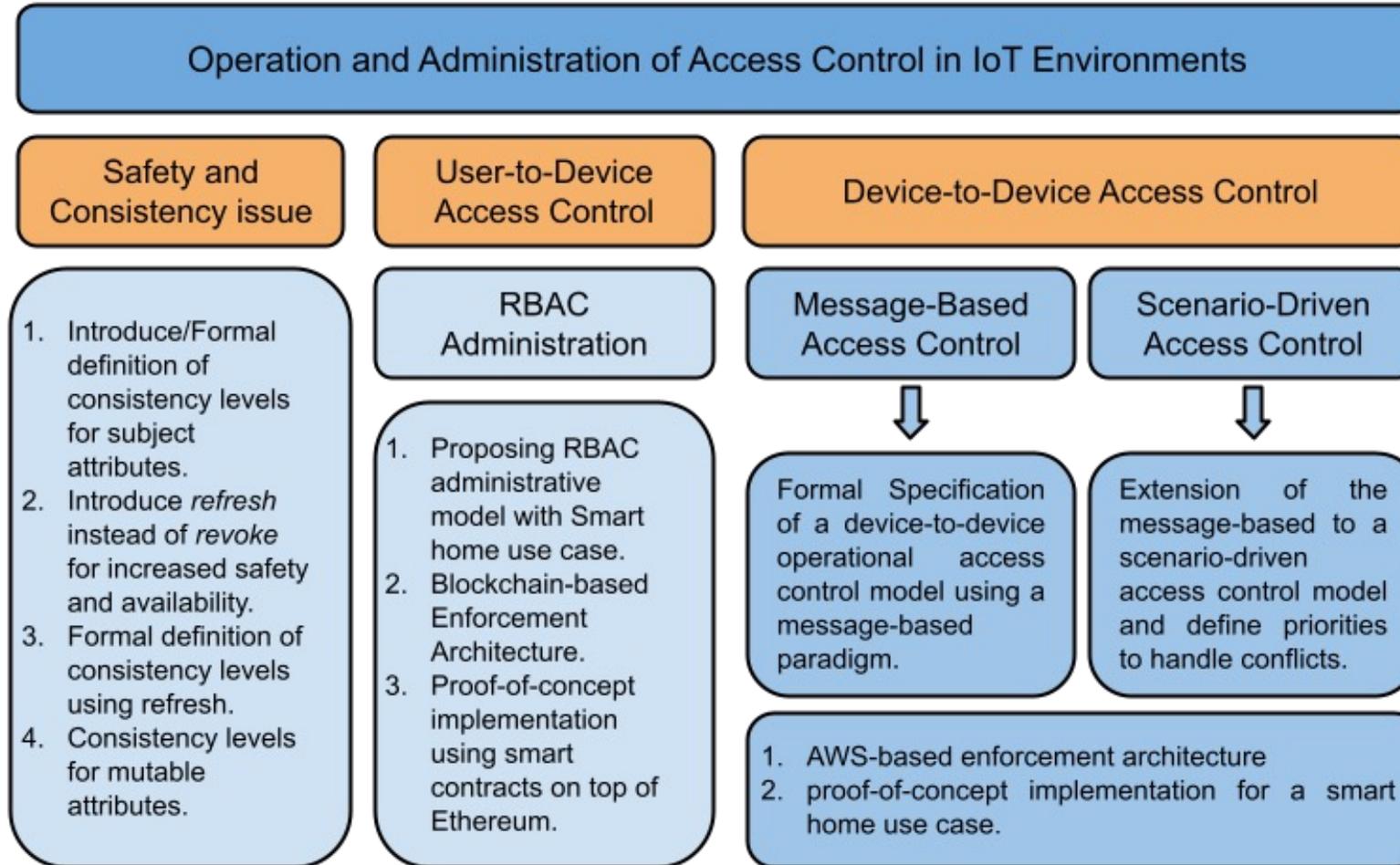
A proof-of-concept implementation of two previously discussed smart home use cases. Statistics collected across 500 trials. All statistics are in milliseconds.

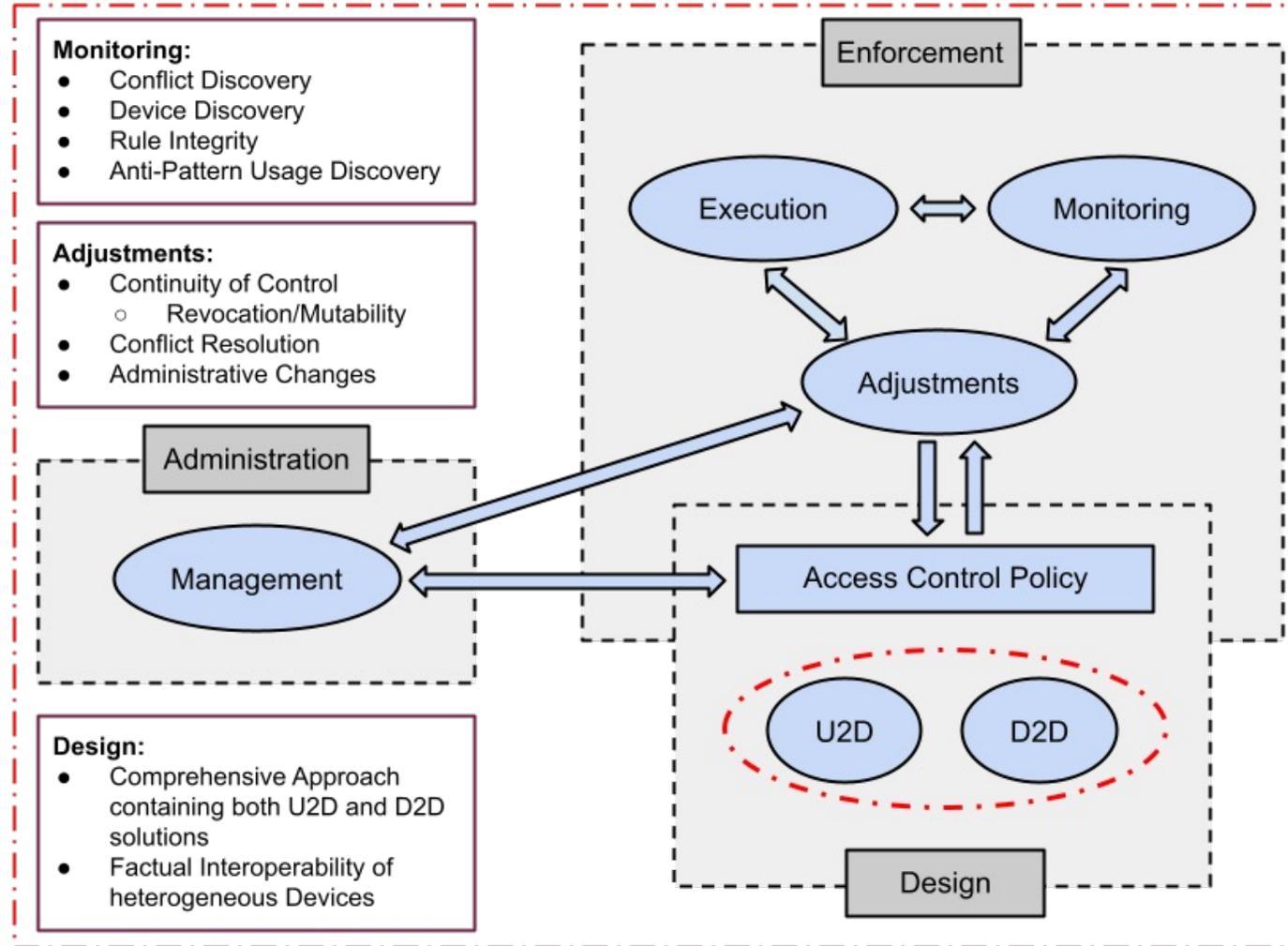
**Full timer:** a device to execute an action using our model/architecture, have lambda process it, and update the respective devices. The average full timer is 35 milliseconds.

**State Update:** The time to update a device's state, e.g., a door lock going from unlocked to locked, is on average 6 milliseconds.

**Time per Action:** depicted in the box-whisker plot. The box shows the 25 percentile and 75 percentile of the data, with median as the line contained inside the box.

- Our model provides specification for mediating access in device-to-device communications for the **first time**. Presented model is **context-aware**, **dynamic** and **lightweight**.
- **Continuity of Access**: Any change in participating devices' attributes **triggers the lambda** which then re-evaluates the policy and **adjusts the access authorizations** accordingly.
  - In order to provide continuity as one of the model's elements, regardless of its enforcement method, the **continuous retrieval and evaluation** of entity/environment attributes **should be incorporated in the model**.
  - Continuity of access control could be concluded when the proposed model is able to **revoke the previously granted access** in case of any unintended change in attributes.
- **Post-Authorization**: Any message communication may result in a change of its sender/receiver attributes which is indicated through **Impact** function in our model.
- **Architecture Agnostic**: Our model is **not peculiar to Amazon AWS**.
  - We use AWS because of its simplicity, security and flexibility and being agnostic to device type and OS.
  - AWS IoT Device Management is agnostic to device type and OS.





### Conference/Workshop Papers:

#### Chapter 3

1. Mehrnoosh Shakarami, and Ravi Sandhu. "Safety and Consistency of Subject Attributes for Attribute-Based Pre-Authorization Systems", In Proceedings of National Cyber Summit (NCS'19), pp. 248-263. Springer, Cham, 2019. **(Published)**
2. Mehrnoosh Shakarami, and Ravi Sandhu. "Refresh instead of revoke enhances safety and availability: A formal analysis." In IFIP Annual Conference on Data and Applications Security and Privacy (DBSec'19), pp. 301-313. Springer, Cham, 2019. **(Published)**
3. Mehrnoosh Shakarami, and Ravi Sandhu. "Safety and Consistency of Mutable Attributes Using Quotas: A Formal Analysis." In 2019 First IEEE International Conference on Trust, Privacy and Security in Intelligent Systems and Applications (TPS'19), pp. 1-9. IEEE, 2019. **(Published)**

#### Chapter 4

4. Mehrnoosh Shakarami, and Ravi Sandhu. "Role-Based Administration of Role-Based Smart Home IoT", In Proceedings of the 2021 ACM Workshop on Secure and Trustworthy Cyber-Physical Systems (SaT-CPS'21), pp. 49-58. 2021. **(Published)**
5. Mehrnoosh Shakarami, James Benson, and Ravi Sandhu. "Blockchain-Based Administration of Access in Smart Home IoT", In Proceedings of the 2022 ACM Workshop on Secure and Trustworthy Cyber-Physical Systems (SaT-CPS'22), 2022. **(Published)**

#### Chapter 5

6. Mehrnoosh Shakarami, James Benson, and Ravi Sandhu. "Scenario-Driven Device-to-Device Access Control in Smart Home IoT". IFIP Annual Conference on Data and Applications Security and Privacy (DBSec'22), Springer, Cham, 2022. **(To be Submitted at Mid-April)**

THANKS FOR TAKING THE TIME TO BE HERE!

