The *RCL2000* Language for Specifying Role-Based Authorization Constraints

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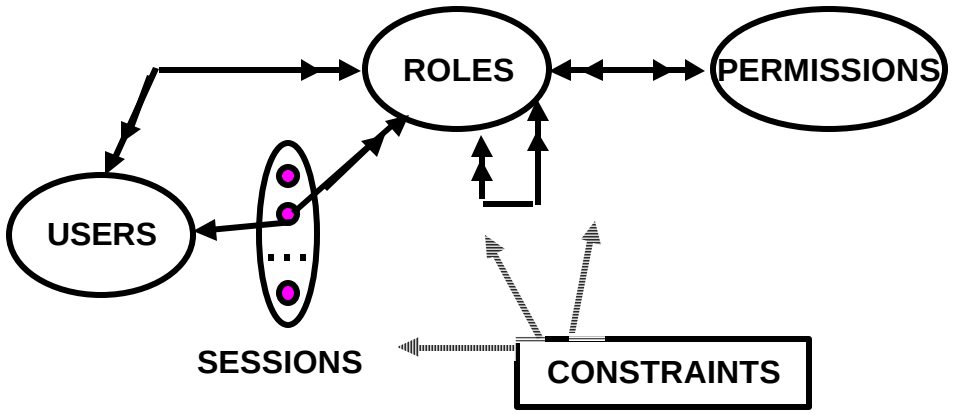
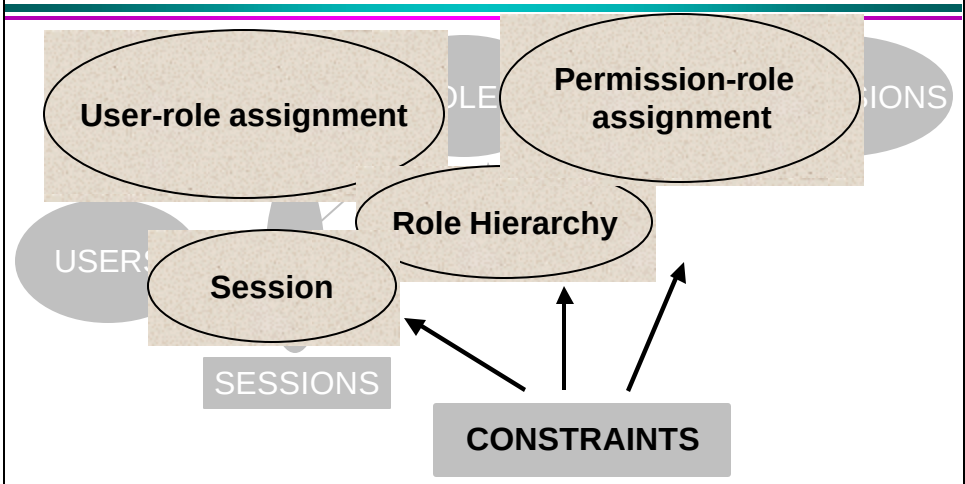
ABSTRACT

- **This presentation includes**
  - **The first formal (and intuitive) language for role-based authorization constraints**
  - **A formal semantics for this language**
  - **Demonstration of the expressive power of the language**
  - **Characterization of role-based constraints into prohibition and obligation constraints**

RBAC96



RBAC96



SEPARATION OF DUTY (1)

- **SOD is fundamental technique for preventing fraud and errors**
- **Related Work**
  - **Enumerate several forms of SOD**
  - **Little work on specifying SOD in a comprehensive way**

SEPARATION OF DUTY (2)

PURCHASING MANAGER

ACCOUNTING PAYABLE MANAGER

PROHIBITION

- **Separation of Duty constraints**

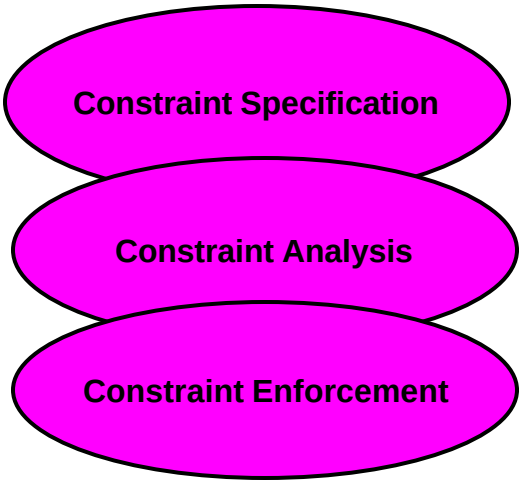
OBLIGATION

- **Every faculty member must be assigned to at least one departmental committee**

RESEARCH PLAN

- Need to specify these constraints
  - Language
- Show the meaning of expression
  - Formal semantics
- Expressive power of the language
  - Well-known constraints and simulations
- Analysis of the work
  - Characterization

BIG PICTURE



WHO IS THE USER

- Security Researcher
- Security Policy Designer
- Security Architect

RCL 2000

- RCL 2000 (Role-based Constraints Language 2000)
- Specification Language
  - to formally express constraints in role-based systems
- Most components are built upon RBAC96

BASIC ELEMENT (from RBAC96)

- U : a set of users
- R : a set of roles
  - RH Í R ´ R : role hierarchy
- OBJ : a set of objects
- OP : a set of operations
- P = OP ´ OBJ : a set of permissions
- S : a set of sessions

BASIC ELEMENT (from RBAC96)

- UA : a many-to-many user-to-role assignment relation
- PA : a many-to-many permissions-to-role assignment relation

SYSTEM FUNCTIONS (from RBAC96)

- user : $R \rightarrow 2^U$
- roles : $U \times P \times S \rightarrow 2^R$
- sessions : $U \rightarrow 2^S$
- permissions : $R \rightarrow 2^P$
- operations : $R \times OBJ \rightarrow 2^{OP}$
- object : $P \rightarrow 2^{OBJ}$

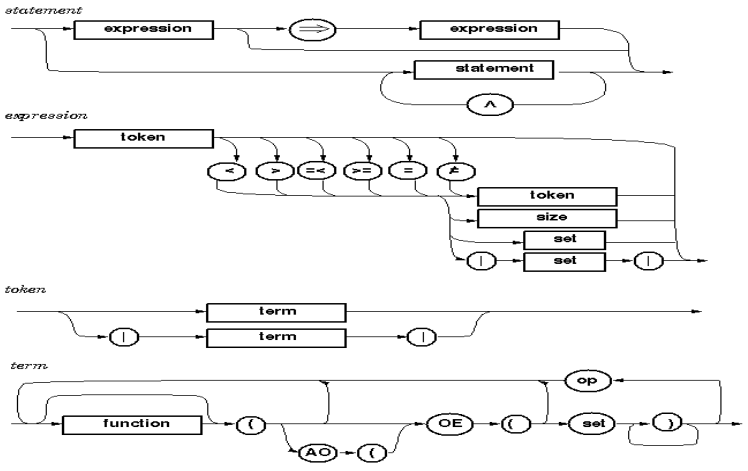
BASIC ELEMENT (beyond RBAC96)

- CR : all conflicting role sets
- CU : all conflicting user sets
- CP : all conflicting permission sets

NON-DETERMINISTIC FUNCTIONS (beyond RBAC96)

- **introduced by Chen and Sandhu (1995)**
- **oneelement (OE)**
  - **oneelement(X) = x_i, where x_i ∈ X**
- **allother (AO)**
  - **allother(X) = X - {OE(X)}**
    = X - {x_i}
    - **should occur along with OE function**

SYNTAX



EXAMPLES OF CONSTRAINT EXPRESSION

Conflicting roles cannot have common users

- **|roles(OE(U)) Ç OE(CR)| £1**

Conflicting users cannot have common roles

- **roles(OE(OE(CU))) Ç roles(AO(OE(CU))) = f**

Users cannot activate two conflicting roles

- **|roles(sessions(OE(U))) Ç OE(CR)| £1**

Users cannot activate two conflicting roles in a single session

- **| roles(OE(sessions(OE(U)))) Ç OE(CR)| £1**

FORMAL SEMANTICS

- **Reduction Algorithm**
  - **to convert a constraint expression to a restricted form of first order predicate logic (RFOPL)**
- **Construction Algorithm**
  - **to construct a constraint expression from RFOPL**

REDUCTION ALGORITHM

$OE(OE(CR)) \in roles(OE(U)) \Rightarrow AO(OE(CR)) \cap roles(OE(U)) = \emptyset$

1. $OE(OE(CR)) \in roles(OE(U)) \Rightarrow (OE(CR) - \{OE(OE(CR))\}) \cap roles(OE(U)) = \emptyset$

2. $\forall cr \in CR : OE(cr) \in roles(OE(U)) \Rightarrow (cr - \{OE(cr)\}) \cap roles(OE(U)) = \emptyset$

3. $\forall cr \in CR, \forall r \in cr : r \in roles(OE(U)) \Rightarrow (cr - \{r\}) \cap roles(OE(U)) = \emptyset$

4. $\forall cr \in CR, \forall r \in cr, \forall u \in U : r \in roles(u) \Rightarrow (cr - \{r\}) \cap roles(u) = \emptyset$

RFOPL STRUCTURE

- sequence part : predicate
- $\forall r \in R, \forall u \in U : r \in roles(u)$
- $\forall x_2 \in x_1, \forall x_3 \in x_2, \forall x_4 \in x_3$: predicate

CONSTRUCTION ALGORITHM

$\forall cr \in CR, \forall r \in cr, \forall u \in U : r \in roles(u) \Rightarrow (cr - \{r\}) \cap roles(u) = \emptyset$

1. $\forall cr \in CR, \forall r \in cr : r \in roles(OE(U)) \Rightarrow (cr - \{r\}) \cap roles(OE(U)) = \emptyset$

2. $\forall cr \in CR : OE(cr) \in roles(OE(U)) \Rightarrow (cr - \{OE(cr)\}) \cap roles(OE(U)) = \emptyset$

3. $OE(OE(CR)) \in roles(OE(U)) \Rightarrow (OE(CR) - \{OE(OE(CR))\}) \cap roles(OE(U)) = \emptyset$

4. $OE(OE(CR)) \in roles(OE(U)) \Rightarrow AO(OE(CR)) \cap roles(OE(U)) = \emptyset$

SOUNDNESS AND COMPLETENESS

- **Theorem 1** *Given RCL2000 expression* **a**, **a** *can be translated into RFOPL expression* **b**. *Also* **a** *can be reconstructed from* **b**.

  $C(R(a)) = a$

- **Theorem 2** *Given RFOPL expression* **b**, **b** *can be translated into RCL2000 expression* **a**. *Also* **b´** *which is logically equivalent to* **b** *can be reconstructed from* **a**.

  $R(C(b)) = b'$

SEPARATION OF DUTY CONSTRAINTS

- **Specification of SOD constraints identified by Simon and Zurko (1997) and formulated by Virgil et al (1998)**
- **Identify new SOD properties**
  - **Role-centric**
  - **User-centric**
  - **Permission-centric**

ROLE-CENTRIC SOD CONSTRAINT EXPRESSION

- **Static SOD**

  : Conflicting roles cannot have common users

  U = $\{u_1,u_2,\ldots u_n\}$, R = $\{r_1,r_2,\ldots r_n\}$,
  CR = $\{cr_1,cr_2\}$: $cr_1 = \{r_1,r_2,r_3\}$, $cr_2 = \{r_a,r_b,r_c\}$

  - **|roles(OE(U)) Ç OE(CR)| £1**

PERMISSION-CENTRIC SOD CONSTRAINT EXPRESSION

- **SSOD-CP**
  - **|permissions(roles(OE(U))) Ç OE(CP)| £1**
- **Variations of SSOD-CP**
  - **SSOD-CP Ù**
    |permissions(OE(R)) Ç OE(CP)| £1

USER-CENTRIC SOD CONSTRAINT EXPRESSION

- **SSOD-CU (User-centric)**
  - **SSOD-CR Ù |user(OE(CR)) Ç OE(CU)| £1**

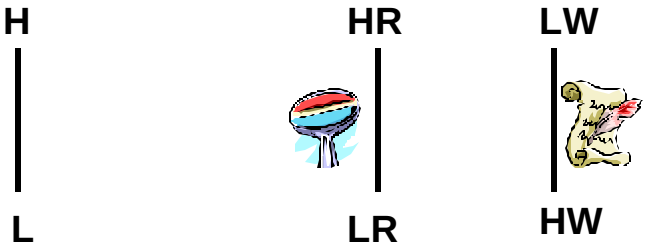
DYNAMIC SOD

- User-based DSOD
  - |roles(sessions(OE(U))) Ç OE(CR)| £1
- User-based DSOD with CU
  - |roles(sessions(OE(OE(CU)))) Ç OE(CR)| £1
- Session-based DSOD
  - |roles(OE(sessions(OE(U)))) Ç OE(CR)| £1
- Session-based DSOD with CU
  - |roles(OE(sessions(OE(OE(CU))))) Ç OE(CR)| £1

CASE STUDIES

- Lattice-based access control
  - Ravi Sandhu (1993, 1996)
- Chinese Wall policy
  - Ravi Sandhu (1992)
- Discretionary access control
  - Sandhu and Munawer (1998)

LATTICE-BASED ACCESS CONTROL



- Subject *s* can write object *o* only if $\lambda(s) \geq \lambda(o)$
- Subject *s* can read object *o* only if $\lambda(o) \geq \lambda(s)$

Constraints on UA: ***Each user is assigned to exactly two roles xR and LW***

LATTICE-BASED ACCESS CONTROL

- AR = {ar1, ar2}
  - ar1={HR, HW}, ar2={LR, LW}
- ASR = {asr1, asr2}
  - asr1={HR, LW}, asr2={LR, LW}

- **Constraint on UA:**
  - roles(OE(U)) = OE(ASR)
- **Constraint on sessions:**
  - roles(OE(sessions(OE(U)))) = OE(AR)

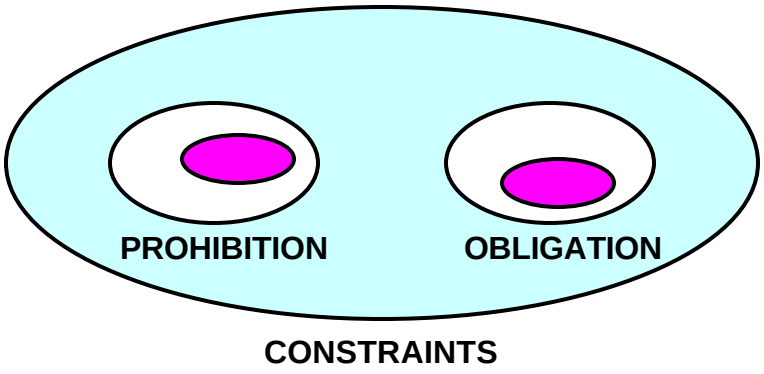
PROHIBITION CONSTRAINTS

- **Forbid the RBAC component from doing (or being) something which is not allowed to do (or be)**
  - **Separation of duty constraints**

OBLIGATION CONSTRAINTS

- **Force the RBAC component to do (or be) something**
  - **LBAC-RBAC, Chinese Wall-RBAC simulation**

CONSTRAINTS CHARACTERIZATION



SIMPLE PROHIBITION CONSTRAINTS

- **Type 1**
  - ***½expr½£ 1***
- **Type 2**
  - ***expr* = f or ½*expr*½= 0**
- **Type 3**
  - ***½expr1½<½expr2½***

SIMPLE OBLIGATION CONSTRAINTS

- **Type 1**
  - ***expr* ¹ 0 or ½*expr*½> 0**
- **Type 2**
  - **Set X = Set Y**
- **Type 3**
  - **obligation constraints Þ obligation constraints**
- **Type 4**
  - **½*expr*½ = 1**
    - **½*expr*½ = 1 • ½*expr*½£ 1 Ù ½*expr*½> 0**

CONTRIBUTIONS

- **Developed the first formal and intuitive language for role-based authorization constraints**
- **Provided a formal semantics for this language**
- **Demonstrated the expressive power of the language by**
  - **specifying well-known separation of duty constraints**
  - **identifying new role-based SOD constraints**
  - **showing how to specify constraints identified in the simulations of other policies in RBAC**
- **Characterized role-based constraints into prohibition and obligation constraints**

FUTURE WORK

- **Extension of RCL 2000**
  - **Applying it the formalization of some realistic security policies**
- **Implementation Issue**
  - **Tool for checking syntax and semantic as well as visualization of specification**
- **Enforcement of constraints**