### **INFS 767 Fall 2003**

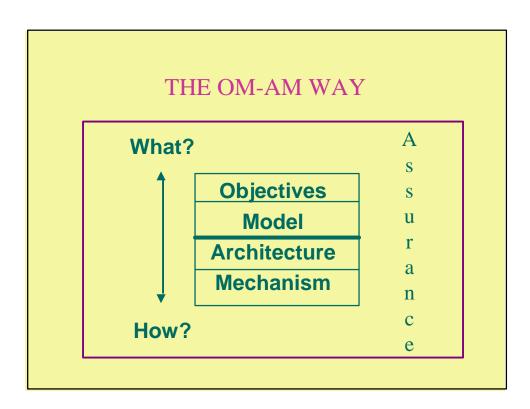
### **RBAC Architectures and Mechanisms**

Prof. Ravi Sandhu

# AUTHORIZATION, TRUST AND RISK

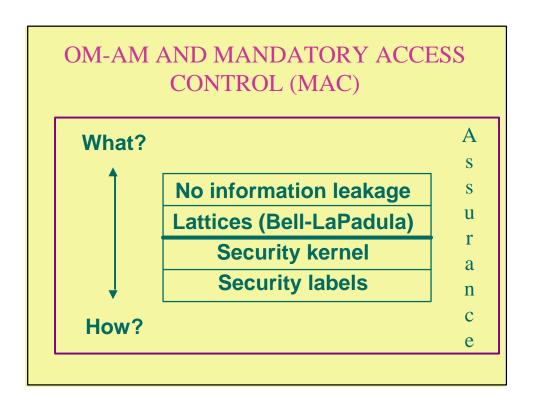
- Information security is fundamentally about managing
  - > authorization and
  - > trust

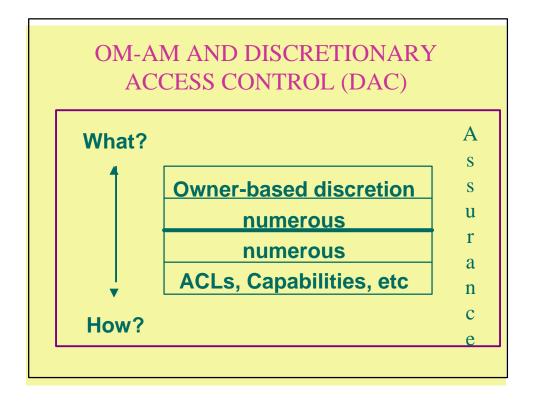
so as to manage risk



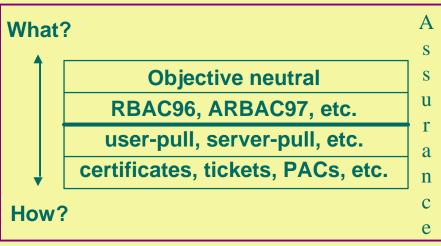
# LAYERS AND LAYERS

- Multics rings
- Layered abstractions
- Waterfall model
- Network protocol stacks
- Napolean layers
- RoFi layers
- ◆ OM-AM
- etcetera





# OM-AM AND ROLE-BASED ACCESS CONTROL (RBAC)



# DISTRIBUTED RBAC (DRBAC) CASE STUDY

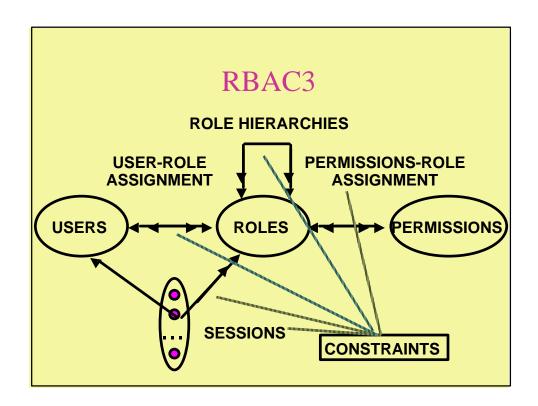
- Approximately a dozen physical sites
- Approximately 2-3 simulation models/site
- Fewer than 100 roles structured in a very shallow hierarchy
  - > A subset of roles is used in any single simulation model
- Fewer than 100 users
- A user uses only one role at a time
  - > Convenient but not critical
- Moderate rate of change

# DISTRIBUTED RBAC (DRBAC) CASE STUDY

- Permission-role assignment
  - Locally determined at each simulation model
- User-role assignment
  - A user can be assigned to a role if and only if <u>all</u> simulation models using that role agree
  - > A user is revoked from a role if and only if <u>any</u> simulation model using that role revokes the user

# DISTRIBUTED RBAC (DRBAC) CASE STUDY

- Each simulation model has a security administrator role authorized to carry out these administrative tasks
- A simulation model can assign permissions to a role X at any time
  - even if X is previously unused in that simulation model
- Consequently any simulation model can revoke any user from any role!



# MODEL CUSTOMIZATION

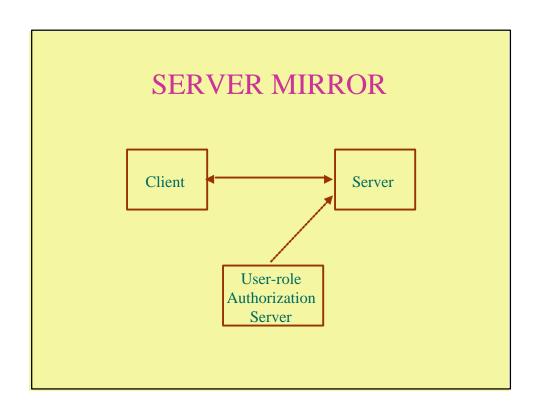
- Each session has a single role
- ❖ SM = {sm1, ..., smk}, simulation models
- ❖ OP = {op1, ..., opl}, operations
- ❖ P= SM X OP, permissions
- SMA = {sma1, ..., smk}, administrative roles
- ❖ R Ç SMA = Æ
- ❖ Admin: SM « SMA

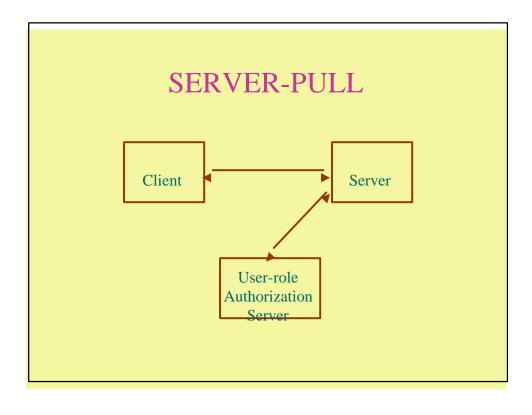
# MODEL CUSTOMIZATION

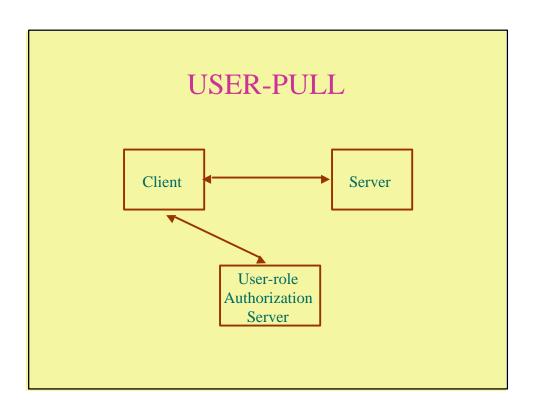
- Can formalize the administrative rules given earlier
- For each simulation model designate a unique user to be the chief security administrator who is authorized to assign and revoke users from the security administrator role for that model

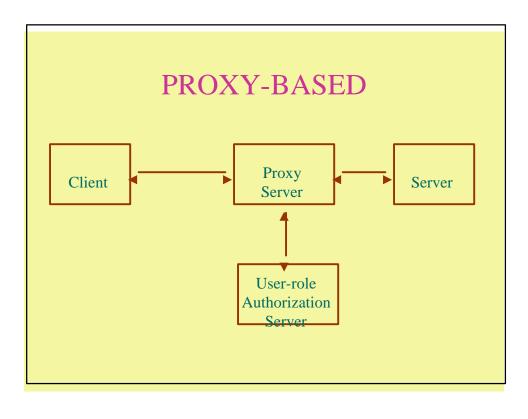
# DRBAC ARCHITECTURES

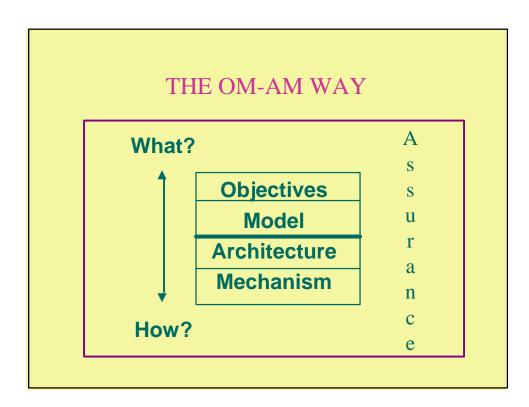
- Permission-role
  - Enforced locally at each simulation model
- Permission-role administration
  - > Enforced locally at each simulation model
  - May need to communicate to other simulation models
- User-role
  - > See following slides
- User-role administration
  - Centralized or decentralized











### Secure Attribute Services on the Web

### WWW (World Wide Web)

- widely used for electronic commerce and business
- > supports synthesis of technologies
- mostly, Web servers use identity-based access control
  - scalability problem

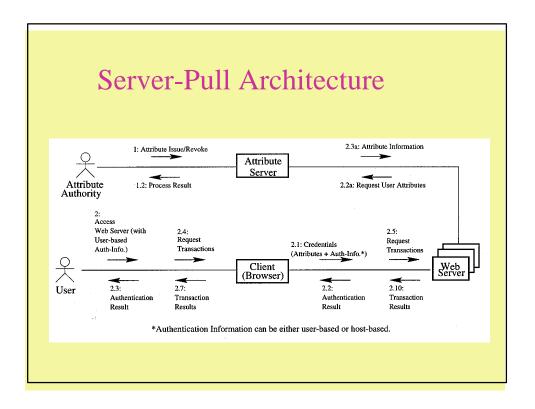
# Background

- An attribute
  - > a particular property of an entity
    - e.g., role, identity, SSN, clearance, etc.
- If attributes are provided securely,
  - Web servers can use those attributes
    - e.g., authentication, authorization, access control, electronic commerce, etc.
- A successful marriage of the Web and secure attribute services is required

#### User-Pull Architecture 1: Attribute Issue/Revoke Attribute Server 1.2: Process Result Attribute Authority Attribute Attribute 2.4: Access Information Request 2.9: 2: Attribute Web Server (with Request User-based Request 2.5: Credentials Request Auth-Info.) Transactions (Attributes + Auth-Info.\* Client (Browser) User 2.3: 2.7: 2.11: 2.6: 2.10: Attribute Validation Validation Transaction Request Result Results Result Results Result \*Authentication Information can be either user-based or host-based.

## User-Pull Architecture

- Each user
  - pulls appropriate attributes from the Attribute Server
  - presents attributes and authentication information to Web servers
- Each Web server
  - requires both identification and attributes from users
- High performance
  - > No new connections for attributes



# Related Technologies

### Cookies

- in widespread current use for maintaining state of HTTP
- becoming standard
- > not secure
- Public-Key Certificates (X.509)
  - support security on the Web based on PKI
  - > standard
  - > simply, bind users to keys
  - > have the ability to be extended

# Cookies

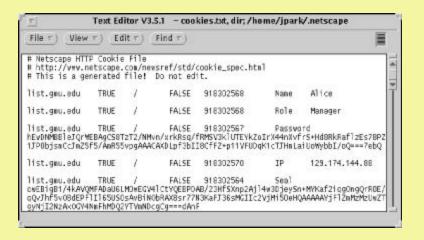
_	Domain	Flag	Path	Cookie_Name	Cookie_Value	Secure	Date
Cookie 1	acme.com	TRUE	1	Name	Alice	FALSE	12/31/99
			:		:		
Cookie n	acme.com	TRUE	/	Role	manager	FALSE	12/31/99

# Security Threats to Cookies

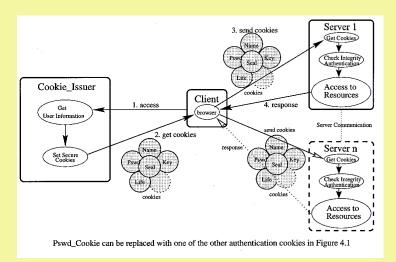
- Cookies are not secure
  - > No authentication
  - No integrity
  - > No confidentiality
- can be easily attacked by
  - > Network Security Threats
  - > End-System Threats
  - Cookie Harvesting Threats

#### Secure Cookies on the Web Domain\_ Flag Path Cookie\_Name Cookie\_Value Secure Date TRUE FALSE 12/31/99 Name\_Cookie acme.com Name\_Cookie Alice\* TRUE FALSE 12/31/99 acme.com Role\_Cookie manager\* Role\_Cookie Life Cookie acme.com TRUE Life\_Cookie 12/31/99 FALSE 12/31/99 Pswd\_Cookie TRUE Pswd\_Cookie hashed\_password FALSE 12/31/99 Key\_Cookie TRUE Key\_Cookie encryped\_key\* FALSE (Optional) Sealing Cookies Seal\_of\_Cookies\*\* FALSE Seal\_Cookie TRUE Seal\_Cookie acme.com \* Sensitive fields can be encrypted in the cookies. \*\* Seal of Cookies can be either MAC or signed message digest of cookies. Note: Pswd\_Cookie can be replaced with one of the other authentication cookies in Figure 4.1

# A Set of Secure Cookies



# How to Use Secure Cookies



# **Applications of Secure Cookies**

- User Authentication
- Electronic Transaction
- Eliminating Single-Point Failure
- Pay-per-Access
- Attribute-based Access Control

# **Authentication Cookies**

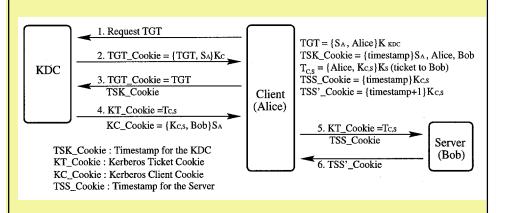
	Domain	Flag	Path	Cookie_Name	Cookie_Value	Secure	Date			
IP_Cookie	acme.com	TRUE	1	IP_Cookie	129.174.100.88	FALSE	12/31/99			
Pswd_Cookie	acme.com	TRUE	1	Pswd_Cookie	hashed_password	FALSE	12/31/99			
KT_Cookie	acme.com	TRUE	1	Kerberos_Ticket	{Alice, K c,s}Ks	FALSE	12/31/99			
Sign_Cookie	acme.com	TRUE	1	Sign_Cookie	Signature_of_Alice	FALSE	12/31/99			

# Server-Pull Architecture

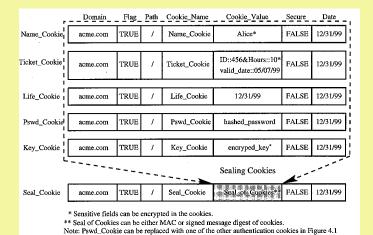
- Each user
  - presents only authentication information to Web servers
- Each Web server
  - pulls users' attributes from the Attribute Server
- Authentication information and attribute do not go together
- More convenient for users
- Less convenient for Web servers

#### Secure Cookies for **Electronic Transactions** Flag Path Cookie\_Name \_Cookie\_Value TRUE Name\_Cookie Alice\* FALSE 12/31/99 Name Cookie acme.com number::123456789\* exp\_date::Jan.2000\* Card\_Cookie TRUE Card\_Cookie 12/31/99 ID::123&off::10%\* valid\_date::05/07/99\* Coupon\_Cookie TRUE acme.com Coupon\_Cooki FALSE 12/31/99 acme.com TRUE Life\_Cookie 12/31/99 FALSE 12/31/99 TRUE Pswd\_Cookie hashed\_password FALSE 12/31/99 Key\_Cookie Key\_Cookie encryped\_key\* FALSE 12/31/99 Sealing Cookies FALSE 12/31/99 Seal\_Cookie acme.com TRUE Seal Cookie Sensitive fields can be encrypted in the cookies. \*\* Seal of Cookies can be either MAC or signed message digest of cookies. Note: Pswd\_Cookie can be replaced with one of the other authentication cookies in Figure 4.1

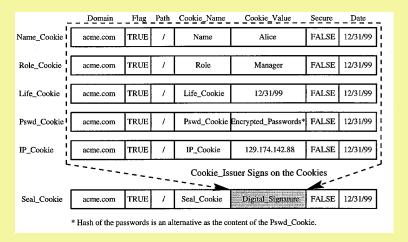
# Kerberos-Based Authentication by Secure Cookies



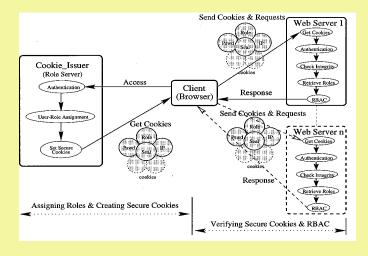
# Secure Cookies for Pay-Per-Access



# Secure Cookies for RBAC



# RBAC on the Web by Secure Cookies



## X.509 Certificate

- Digitally signed by a certificate authority
  - to confirm the information in the certificate belongs to the holder of the corresponding private key
- Contents
  - version, serial number, subject, validity period, issuer, optional fields (v2)
  - > subject's public key and algorithm info.
  - extension fields (v3)
  - digital signature of CA
- Binding users to keys
- Certificate Revocation List (CRL)

## X.509 Certificate

### **Smart Certificates**

### Short-Lived Lifetime

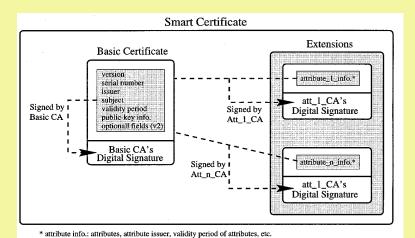
- More secure
  - typical validity period for X.509 is months (years)
  - users may leave copies of the corresponding keys behind
  - the longer-lived certificates have a higher probability of being attacked
- No Certificate Revocation List (CRL)
  - simple and less expensive PKI

## **Smart Certificates**

### Containing Attributes Securely

- Web servers can use secure attributes for their purposes
- Each authority has independent control on the corresponding information
  - basic certificate (containing identity information)
  - each attribute can be added, changed, revoked, or reissued by the appropriate authority
    - e.g., role, credit card number, clearance, etc.
- Short-lived certificate can remove CRLs

# Separate CAs in a Certificate



# **Smart Certificates**

- Postdated Certificates
  - > The certificate becomes valid at some time in the future
  - possible to make a smart certificate valid for a set of duration
  - > supports convenience
- Confidentiality
  - > Sensitive information can be
    - · encrypted in smart certificates
      - e.g. passwords, credit card numbers, etc.

## A Smart Certificate

```
Certificate Content:

Certificate:

Data:

Version: v3 (0x2)

Serial Number: 26 (0x1a)

Signature Algorithm: PKCS #1 MD5 With R8A Encryption

Islamature Algorithm: PKCS #1 MD6 With R8A Encryption

R8A Certificate With R8A Encryption

Public Rev:

Modulus

Ostarithm: PKCS #1 R8A Encryption

Public Rev:

Modulus

Ostarithm: PKCS #1 R8A Encryption

Public Rev:

Modulus

Ostarithm: PKCS #1 R8A Encryption

Public Reponent: 65537 (0x1001)

Tribbischeiaa: 36:4629:0819:09:01:31:03:11:05:77:ee:79:

17:805:30+e:aa:36:4629:0819:99:76:46:457:97:94:19:06:a7:

Public Exponent: 65537 (0x10001)

Extensions:

Identifies: Certificate Type

Certified Usage:

SSE Cilent

Identifies: Authority Key Identifier

Critical: no

Value: hBWDNMBBLOGWBBAGCS8T:TZ/NMVn/xrkRsq/fRMSV3klUTEYK20I

Identifies: Authority Key Identifier

Critical: no

Critical: no

Value: hBWDNMBBLOGWBBAGCS8T:TZ/NMVn/xrkRsq/fRMSV3klUTEYK20I

Identifies: Authority Key Identifier

Critical: no

Signature:

Algorithm: PKCS #1 MD5 With R8A Encryption

Signature:

67:44

Signature:

67:46:12:61:c7:bf:57:07:66:c5:f4:f0:c2:e1:62:27:f6:d6:ae:09:77:

46
```

# **Applications of Smart Certificates**

- On-Duty Control
- Compatible with X.509
- User Authentication
- Electronic Transaction
- Eliminating Single-Point Failure
- Pay-per-Access
- Attribute-based Access Control

# Injecting RBAC to Secure a Web-based Workflow System

Gail-Joon Ahn and Ravi Sandhu
George Mason University

Myong Kang and Joon Park Naval Research Laboratory

# WORKFLOW MANAGEMENT SYSTEMS

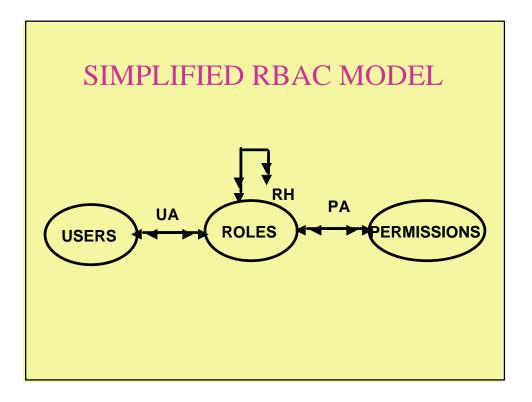
- Control and coordinate processes that may be processed by different processing entities
- □ Received much attention
- ☐ Marriage with Web technology
- Minimal security services

## **OBJECTIVE**

□ Inject role-based access control (RBAC) into an existing web-based workflow system

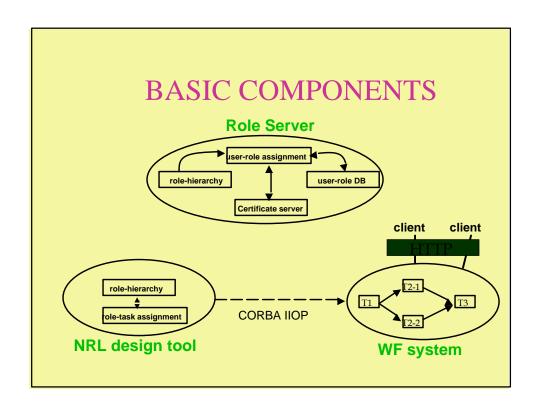
# WHY RBAC?

- □ A mechanism which allows and promotes an organization-specific access control policy based on roles
- □ Has become widely accepted as the proven technology



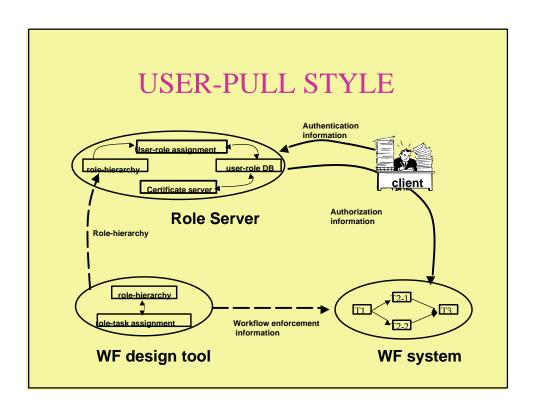
# ROLE-BASED SECURE WORKFLOW SYSTEM

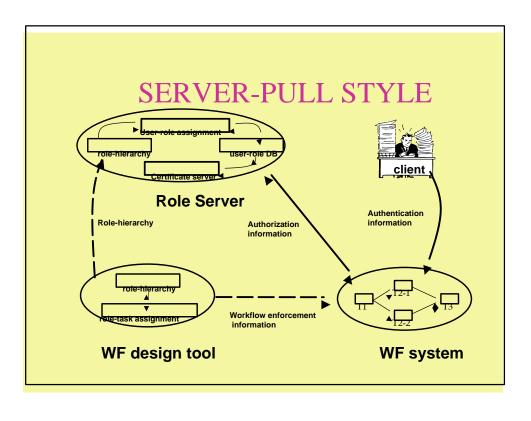
- □ Workflow Design Tool
- □ Workflow (WF) System
- □ Role Server



# **ARCHITECTURES**

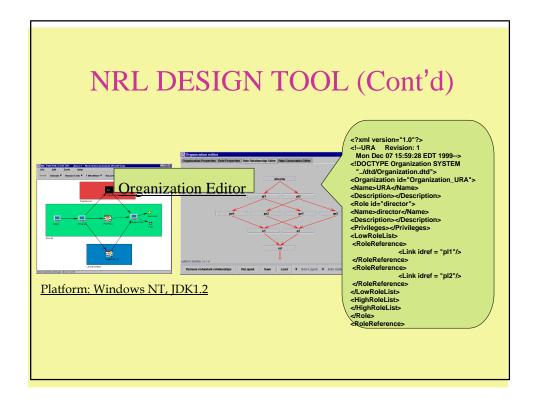
- □ USER-PULL STYLE
- □ SERVER-PULL STYLE





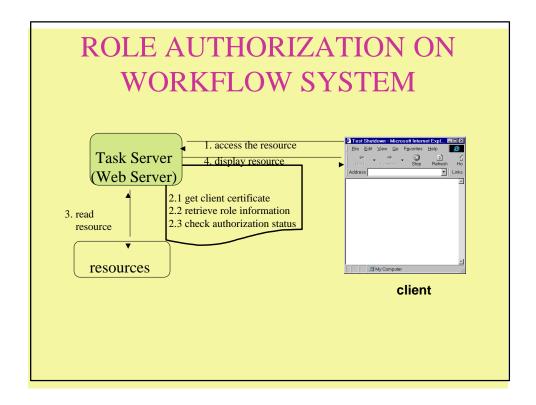
# NRL (Naval Research Lab.) DESIGN TOOL

- □ design workflow model
- □ create role and role hierarchies
- □ assign role to task
- exporting role hierarchies to role server



## **WORKFLOW SYSTEM**

- □ each task server is web server
- user should present client authentication certificate
- user's privilege is authorized by content of certificate (specially client's role information)

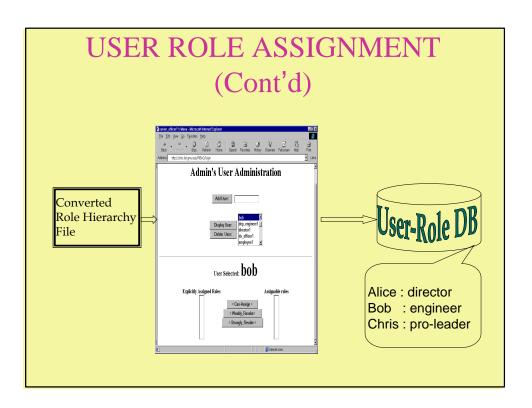


# ROLE SERVER

- □ User Role Assignment
- □ Certificate Server

# **USER ROLE ASSIGNMENT**

- □ maintain role hierarchies and user database
- □ assign users to roles
- □ generate user-role database



# **CERTIFICATE SERVER**

- □ authenticate client
- □ retrieve client's role information from user-role database
- □ issue certificate with client's role information



